

$I(g(x); Y) \text{ vs } I(x, Y)$
 $H(g(x)|x) = 0$ $H(Y|x) = H(Y|x, g(x)); \underline{H(x|x) \geq H(g(x)|Y)}$
 $I(g(x); Y) = H(g(x)) - H(g(x)|Y) \leq H(x) - \underbrace{H(g(x)|Y)}_{?}$

$I(x, g(x); Y) = I(x; Y) + \underbrace{I(g(x); Y|x)}_{=0} =$
 $= I(g(x); Y) + I(x; Y|g(x))$

$I(g(x); Y|x) = \underbrace{H(g(x)|x)}_{=0} - \underbrace{H(g(x)|x, Y)}_{=0} = 0$

$I(x; Y) = I(g(x); Y) + \underbrace{I(x; Y|g(x))}_{\geq 0} \Rightarrow$

$I(x; Y) \geq I(g(x); Y)$

DOVAZANO!!!

- Markov e platziino od DATA-LOCALIZING INEQUALITY KAD
 $I(x; Z) \leq I(x; Y)$ i.e. $I(x; X) \geq I(x; Z)$

$I(x_0 | x_{-1}) \text{ vs } I(x_0 | x_{-1}, x_1)$
 $H(x_{-1}, x_0, x_1) = H(x_{-1}) + H(x_0 | x_{-1}) + H(x_1 | x_{-1}, x_0)$
 $H(x_0 | x_{-1}) \geq H(x_0 | x_{-1}, x_1)$ CONDITIONING REDUCES ENTROPY.

$P(x||z) \geq 0$ $\sum_{x \in X} P(x|z) \frac{P(x)}{P(x)} \geq 0$

$I(x; Y) = H(x) - H(x|Y) \geq 0$ $H(x) \geq H(x|Y)$

$I(x_0 | x_1 | x_{-1}) = H(x_0 | x_{-1}) - H(x_0 | x_{-1}, x_1) \geq 0 \Rightarrow$

$H(x_0 | x_{-1}) \geq H(x_0 | x_{-1}, x_1)$

DOVAZANO!!!

$\frac{H(x, Y)}{(H(x) + H(Y))} \text{ vs } 1$

$H(x, Y) = H(x) + \underbrace{H(Y|x)}_{\text{CONDITIONING REDUCES ENTROPY}} \leq H(x) + H(Y)$

$H(x, Y) \leq H(x) + H(Y)$

$\frac{H(x, Y)}{H(x) + H(Y)} \leq 1$

DOVAZANO!!!

• HW2S SOLUTIONS: (ILITE RESULTATI NO MAJHU POZICIJENI TUKO)

(a) $Y = f(X) = 5X$, f^{-1} IS BIJECTIVE FUNCTION I.E

$$X = f^{-1}(Y) = g(Y)$$

- ZNAEME DEKA: $H(X) \geq H(f(X)) = H(5X)$

- OO DRUGA STRANA $H(Y) \geq H(g(Y)) = H(X)$

$$H(5X) \geq H(X)$$

ZNAJI EDNAŠ PODIVAJI DEKA $H(X) \geq H(5X)$ * EDNAŠ
 DEKA $H(5X) \geq H(X) \Rightarrow \boxed{H(X) = H(5X)}$

(b) $I(g(Y); Y)$ vs. $I(X; Y)$

MOZE PA ŠT FORMIRA MARKOV CHAIN:

$$Y \rightarrow X \rightarrow g(X)$$

OVA SEQRVA
 NE ZAVISI OD Y
 T.E ŠE PADOU ZA
 MARKOV LANCEL

- DATA PROCESSING INEQUALITY:

$$I(Y; X) \geq I(Y; g(X)) \Rightarrow I(X; Y) \geq I(g(X); Y)$$

$$P(X, Z | Y) = \frac{P(X, Y, Z)}{P(Y)} = \frac{P(X|Y) \cdot P(Z|X, Y)}{P(Y)} = \frac{P(X|Y) \cdot P(Z|Y)}{P(Y)}$$

2.43 MUTUAL INFORMATION OF HEADS AND TAILS

(a) CONSIDER A FAIR COIN FLIP. WHAT IS THE MUTUAL INFORMATION BETWEEN TOP AND BOTTOM SIDES OF THE COIN?

(b) A SIX-SIDE FAIR DIE IS ROLLED. WHAT IS THE MUTUAL INFORMATION BETWEEN THE TOP SIDE AND THE FRONT SIDE (THE SIDE MOST FACING YOU)?

~~(a) $X \in \{PETA, GLAVA\} = \{T, H\}$
 $P(X) = \left\{ \begin{matrix} 1/2 \\ 1/2 \end{matrix} \right\}$~~

~~$X \in \{TAIL, NON-TAIL\}$ $Y \in \{HEAD, NON-HEAD\}$~~

$X \backslash Y$	H	NH	$P(X)$
T	0	$\frac{1}{2}$	$\frac{1}{2}$
NT	$\frac{1}{2}$	0	$\frac{1}{2}$
$P(Y)$	$\frac{1}{2}$	$\frac{1}{2}$	

$I(X; Y) = H(X) - H(X|Y)$

$P(X Y)$	$X \backslash Y$	H	NH
T	0	1	
NT	1	0	

$P(X|Y) = \frac{P(X, Y)}{P(Y)}$

$$H(X) = \frac{1}{2} \log 2 + \frac{1}{2} \log 2 = \frac{1}{2} + \frac{1}{2} = 1$$

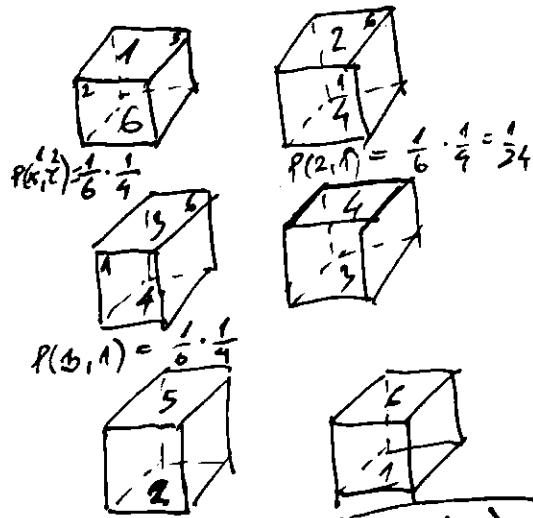
$$H(X|Y) = \sum_{x,y} p(x,y) \log \frac{1}{p(x|y)} = 0 \log 0 + \frac{1}{2} \log 1 + \frac{1}{2} \log 1 = 0$$

$$I(X;Y) = H(X) - H(X|Y) = 1 - 0 = 1 \text{ bit}$$

LOGIČNO E! AUKO 60 ZNAČEŠ Y = SIGURNO 60 ZNAČEŠ X

⑥ X = GORNJA STRANA E {1,2,3,4,5,6}

Y = PREDNA STRANA E {1,2,3,4,5,6}



X \ Y	1	2	3	4	5	6	P(X)
1	0	1/24	1/24	1/24	1/24	0	1/6
2	1/24	0	1/24	1/24	0	1/24	1/6
3	1/24	1/24	0	0	1/24	1/24	1/6
4	1/24	1/24	0	0	1/24	1/24	1/6
5	1/24	0	1/24	1/24	0	1/24	1/6
6	0	1/24	1/24	1/24	1/24	0	1/6
P(Y)	1/6	1/6	1/6	1/6	1/6	1/6	1/6

X \ Y	1	2	3	4	5	6
1	0	1/4	1/4	1/4	1/4	0
2	1/4	0	1/4	1/4	0	1/4
3	1/4	1/4	0	0	1/4	1/4
4	1/4	1/4	0	0	1/4	1/4
5	1/4	0	1/4	1/4	0	1/4
6	0	1/4	1/4	1/4	1/4	0

$$I(X;Y) = H(Y) - H(Y|X)$$

$$H(Y) = 6 \cdot \frac{1}{6} \log 6 = \log 6 = 2.58$$

$$H(Y|X) = - \sum_{x,y} p(x,y) \log p(y|x)$$

$$= 24 \cdot \frac{1}{24} \cdot \log 4 = \log 4 = 2$$

$$I(X;Y) = \log 6 - \log 4 = \log \frac{3}{2} = 0.585$$

VREDNOST

- POZNAVANJE TO NA GORNJATA STRANA ZA NAMAZANA NEKOLIKO NOSTA NA VREDNOST NA PREDNATA STRANA ZA CEI 2 BITA !!!

- ISTO 60 LEIČE VO ODAŠ SO FORMULATA HW2S SO TOA ŠTO DIREKTNO ZA TRANS-FORMACIJA
- $I(B;T) = H(B) - H(B|T) = (AUKO 60 ZNAČEŠ TO) = H(B) - 0 = 2 \cdot \frac{1}{2} \log 2 = 1$
- $I(T;F) = H(T) - H(T|F) = 6 \cdot \frac{1}{6} \log 6 - 4 \cdot \frac{1}{4} \log 4 = \log \frac{3}{2}$

AUKO PREGLEDUVA DODANA GORNJA STRANA ZA PREDNATA STRANA 3 INA 4 MOŽDAŠ SO EDNAKVI VREDI

• PRIMEROT NA REIZNETO VO HVZS TE OHRABLJENA DA NE ODISI SO CEZATA FORMULA ZA $H(Y|X)$ TUKU SO SKLADENATA T.R. $Y|X$ SO TRESTAJI KAKO ODMA SKUCATA ALOMBENIVA ZA DA NE ODISI SO FORMUZATA $-\sum_{x,y} p(x,y) \ln p(y|x)$.

- SERAK MOETO LOSIJE E MATEMATIKI POIZDABANO A REIZNETO VO HVZS E POPLANTICNO I POINZENE-BSKO.

2.44 FAIR RANDOMNES. WE WISH TO USE THREE-SIDED COIN TO GENERATE FAIR COIN TOSS. LET THE COIN HAVE PROBABILITY MASS FUNCTION

$$X = \begin{cases} A, & p_A \\ B, & p_B \\ C, & p_C \end{cases} \quad \boxed{p_C = 1 - p_A - p_B}$$

- WHERE p_A, p_B, p_C ARE UNKNOWN.
- (a) HOW WOULD YOU USE TWO INDEPENDENT FIAT X_1, X_2 TO GENERATE (IF POSSIBLE) BERNOULLI $(\frac{1}{2})$ RANDOM VARIABLE, Z ?
- (b) WHAT IS THE RESULTING MAXIMUM EXPECTED NUMBER OF FAIR FIAT GENERATED?

- (a) AA, AB, AC, BB, BA, BC, CC, CA, CB

$$H(X) = -p_A \ln p_A - p_B \ln p_B - p_C \ln p_C$$

$X \in \{A, B, C\}$ $W \in \{W_1, W_2, \dots, W_{2^n}\}$

$$\boxed{W_i = \{W_{i1}, W_{i2}\}}_{n=2}$$

$$\boxed{H(W) = n H(X)}$$

$$H(X) = -p_A \ln p_A - p_B \ln p_B - (1 - p_A - p_B) \ln p_C =$$

$$= -p_A \ln p_A - p_B \ln p_B - \ln p_C + p_A \ln p_C + p_B \ln p_C$$

$$= p_A \ln \frac{p_C}{p_A} + p_B \ln \frac{p_C}{p_B} - \ln p_C = \ln \left(\frac{p_C}{p_A} \right)^{p_A} + \ln \left(\frac{p_C}{p_B} \right)^{p_B} - \ln p_C$$

$$= \ln \left[\left(\frac{p_C}{p_A} \right)^{p_A} \cdot \left(\frac{p_C}{p_B} \right)^{p_B} \cdot \frac{1}{p_C} \right]$$

$P(W) = \{ p_A, p_A p_B, p_A p_C, p_B, p_B p_A, p_B p_C, p_C, p_C p_A, p_C p_B \}$

$\frac{p_A + p_B}{\frac{1}{2}} = \frac{1}{2} \quad p_C = \frac{1}{2} \quad \left\{ \frac{1}{8}, \frac{1}{8}, \frac{1}{8}, \frac{1}{8}, \frac{1}{8}, \frac{1}{8}, \frac{1}{8}, \frac{1}{8} \right\}$

NEMA

$$\{AA, AB, AC, BA, BC, CA, CB\}$$

$$H(x) \geq H(f(x))$$

$$\left\{ \underbrace{AA, AB, AC, BA, CA}_{\frac{1}{2}}, \underbrace{BB, BC, CB}_{\frac{1}{2}} \right\}$$

$$\left. \begin{aligned} p_A^2 + 2p_A p_B + 2p_A p_C &= \frac{1}{2} \\ p_B^2 + 2p_B p_C + p_C^2 &= \frac{1}{2} \\ p_A + p_B + p_C &= 1 \end{aligned} \right\}$$

$$\begin{aligned} x^2 + 2xy + 2xz &= \frac{1}{2} \\ y^2 + 2yz + z^2 &= \frac{1}{2} \\ \hline x + y + z &= 1 \\ z &= 1 - (x+y) \end{aligned}$$

$$x^2 + \cancel{2xy} + 2x - \cancel{2xz} - \cancel{2xy} = \frac{1}{2}$$

$$2x^2 + 4x - x^2 = 1$$

$$x^2 - 4x + 1 = 0$$

$$-2x^2 + 4x - 1 = 0$$

$$x_{1,2} = \frac{4 \pm \sqrt{16 - 8}}{4}$$

$$x_{1,2} = 1 \pm \frac{1}{4} \sqrt{8} = 1 \pm \frac{1}{4} \cdot 2\sqrt{2} = 1 \pm \frac{\sqrt{2}}{2} = \frac{2 - \sqrt{2}}{2} = \underline{\underline{0.293}}$$

$$y^2 + 2y(1-x-y) + (1-(x+y))^2 = \frac{1}{2}$$

$$= y^2 + 2y - 2xy - 2y^2 + \left(1 - \frac{2 - \sqrt{2}}{2} - y\right)^2 = \frac{1}{2}$$

$$= y^2 + 2y - 2y + \sqrt{2}y - 2y^2 + \left(\frac{2 - \sqrt{2}}{2} - y\right)^2 = -y^2 + \sqrt{2}y + \left(\frac{2 - \sqrt{2}}{2} - y\right)^2 =$$

$$= -y^2 + \sqrt{2}y + \left(\frac{2}{4} - \frac{2\sqrt{2}}{2}y + y^2\right) = \cancel{-y^2} + \sqrt{2}y + \frac{1}{2} - \sqrt{2}y + \cancel{y^2} =$$

$$\Rightarrow \boxed{\frac{1}{2} = \frac{1}{2}} \quad z = 1 - x - y = 1 - \frac{2 - \sqrt{2}}{2} - y$$

$$z = \frac{2 - 2 + \sqrt{2}}{2} - y = \frac{\sqrt{2}}{2} - y$$

$$\boxed{z = \frac{\sqrt{2}}{2} - y}$$

$$\text{IF: } y = \frac{1}{4} \Rightarrow z = \frac{\sqrt{2}}{2} - \frac{1}{4} = \underline{\underline{0.707 - 0.25 = 0.457}}$$

$$\boxed{0.457 + 0.293 + 0.25 = 0.750 + 0.250 = 1}$$

$$y^2 + 2y\left(\frac{\sqrt{2}}{2} - y\right) + \left(\frac{\sqrt{2}}{2} - y\right)^2 = \cancel{y^2} + \cancel{y\sqrt{2}} - \cancel{2y^2} + \frac{1}{2} - \sqrt{2}y + \cancel{y^2} = \frac{1}{2}$$

$$p(x) = \{p_A, p_B, p_C\} = \{0.293; 0.25; 0.457\}$$

$$H(x) = -0.293 \ln 0.293 - 0.25 \ln 0.25 - 0.457 \ln 0.457 = 1.535$$

$$H(W) = \sum_{i=1}^9 p_i(W) \log p_i(W) = -p_1^2 \log p_1^2 - p_1 p_2 \log p_1 p_2 - \dots - p_2 p_3 \log p_2 p_3$$

$$H(W) = 3.0704$$

$$H(W) = 2 \cdot H(X) = 2 \cdot 1.535 = 3.070$$

THE MAXIMUM NUMBER OF PURE RANDOM BITS THAT CAN BE PRODUCED FROM THREE SIDE COIN IS:

$$H(X)$$

Vo optima sukca:

$$p(x) = \left[\frac{2-\sqrt{2}}{2}, \gamma, \frac{\sqrt{2}}{2} - \gamma \right]$$

$$H(X) = -\frac{2-\sqrt{2}}{2} \log \frac{2-\sqrt{2}}{2} - \gamma \log \gamma - \left(\frac{\sqrt{2}}{2} - \gamma \right) \log \left(\frac{\sqrt{2}}{2} - \gamma \right)$$

$$\frac{dH(X)}{d\gamma} = 0 \quad \left[\gamma \log \gamma + \left(\frac{\sqrt{2}}{2} - \gamma \right) \log \left(\frac{\sqrt{2}}{2} - \gamma \right) \right]' = 0$$

NOTE: Multi-rop MMOCap.uu

$$\gamma = \frac{\sqrt{2}}{4}$$

$$H(X) \text{ MAXIMUM} \\ H_m(X) = 1.5795 \approx 1.580$$

⇒ BEST DISTRIBUTION FOR X IS:

$$p(x) = \{ p_1, p_2, p_3 \} = \left\{ \frac{2-\sqrt{2}}{2}, \frac{\sqrt{2}}{4}, \frac{\sqrt{2}-\sqrt{2}}{4} \right\}$$

$$p(x) = \left\{ \frac{2-\sqrt{2}}{2}, \frac{\sqrt{2}}{4}, \frac{\sqrt{2}}{4} \right\}$$

$$H(W) = 2 \cdot 1.580 = 3.160$$

MAXIMUM NUMBER OF PURE RANDOM BITS.

2.49 FINITE ENTROPY SHOW FOR A DISCRETE RANDOM VARIABLE $X \in \{1, 2, \dots\}$ IF $E[X] < \infty$ THEN $H(X) < \infty$.

JENSEN'S INEQUALITY: IF $f(t)$ IS CONVEX FUNCTION $E[f(X)] \geq f(E[X])$

$$H(X) = -\sum_{x \in X} p(x) \log p(x) = -E[\log p(X)] \geq -\log E[p(X)]$$

$$E[\log p(X)] \leq \log E[p(X)]$$

$$E[p(X)] = \sum_{x \in X} p^2(x)$$

$$-E[\log p(X)] \geq -\log E[p(X)]$$

$$\log(x) \Rightarrow \text{CONCAVE} \\ 2^{-H(X)} = 2^{E[\log p(X)]} \leq E[2^{\log p(X)}] \\ = E[p(X)] = \sum_{x \in X} p^2(x) \\ 2^{-H(X)} \leq \sum_{x \in X} p^2(x)$$

$$H(x) \geq -\log E[\gamma(x)] = -\log \left[\sum_{x \in X} \gamma(x) \right] =$$

$$x \in \{1, 2, \dots, \dots\}$$

$$\gamma(x) = \{p_1, p_2, \dots, \dots\}$$

$$E[x] = \sum_{x \in X} x \cdot \gamma(x) = \sum_{i=1}^{\infty} i \cdot p_i \quad \sum_{i=1}^{\infty} i \cdot x^i = \frac{x}{(1-x)^2}$$

$$E[\log x] = \sum_{x \in X} \gamma(x) \log(x) \leq \log \left[\sum_{x \in X} x \cdot \gamma(x) \right] =$$

$$E[\log x] \leq \log E[x]$$

$$\textcircled{+} \Rightarrow -H(x) \leq \log \sum_{x \in X} \gamma^2(x) \quad H(x) \geq \log \sum_{x \in X} \gamma(x)$$

• ТРЕТА ОА ПОКАЗАНА ДЕНА $H(x) \leq \underline{E[\log x]}$

$$\sum_{x \in X} \gamma(x) \log(x) = \sum_{x \in X} \log x^{\gamma(x)} = \log \left(\prod_{x \in X} x^{\gamma(x)} \right)$$

$$x = \log 2^x \quad \sum_{x \in X} x \gamma(x) = \sum_{x \in X} \log 2^x \gamma(x) = E[\log 2^x]$$

$$\leq \log 2^{E[x]} = E[x]$$

$$H(x) = -E[\log \gamma(x)] \geq -\log E[\gamma(x)] = -\log \sum_{x \in X} \gamma^2(x)$$

$$-H(x) \leq \log \sum_{x \in X} \gamma^2(x) \quad \boxed{2^{-H(x)} \leq \sum_{x \in X} \gamma^2(x)}$$

$$D(\gamma || \gamma) = \sum_{i=0}^{\infty} \gamma_i \log \frac{\gamma_i}{q_i} \geq 0$$

$$\sum_{i=0}^{\infty} \gamma_i \log \gamma_i - \sum_{i=0}^{\infty} \gamma_i \log q_i \geq 0$$

$$\sum_{i=0}^{\infty} \gamma_i \log \gamma_i \geq \sum_{i=0}^{\infty} \gamma_i \log q_i$$

$$-\sum_{i=0}^{\infty} \gamma_i \log \gamma_i \leq -\sum_{i=0}^{\infty} \gamma_i \log q_i$$

$$\boxed{q_i = 2^{-i}}$$

$$-\sum_{i=0}^{\infty} \gamma_i \log \gamma_i \leq -\sum_{i=0}^{\infty} \gamma_i \log 2^{-i} = + \sum_{i=0}^{\infty} i \gamma_i$$

$$\boxed{H(x) \leq +E[x]}$$

ИЗВЕСТНО Е КОЛКОТО ВО ПРЕД. СЛ.

$$\boxed{\log x \leq x - 1}$$

$$\boxed{\log x \leq \frac{1}{\ln 2} (x - 1)}$$

$$f(x) = f(x_0) + \frac{f'(x_0)}{1!} (x - x_0) + \frac{f''(x_0)}{2!} (x - x_0)^2 + \dots$$

$\ln x \leq \frac{1}{\ln 2} (x-1)$ $\ln(x) \leq x-1$ TARTOG
 $E[\ln(x)] \leq E[x] \leq \ln[E(x)+1] \leq E[x]$
 $E[\ln(x)] \leq \frac{1}{\ln 2} E[x-1] = \frac{1}{\ln 2} [E[x] - 1] \geq \frac{1}{\ln 2} [H(x) - 1]$
 $\ln x \leq \ln(x+1) \leq x$

ZNAČI:
 1° $E[x] \geq E[\ln(x)]$
 2° $E[x] \geq H(x)$
 $E[\ln(x)] < \infty$ NE KAKIČIŠA DENA $E(x) < \infty$

$H(x) = -\sum_{x \in X} p(x) \ln p(x) \geq -\sum_{x \in X} p(x) \ln(p(x)+1) \geq$
 $\geq -\sum_{x \in X} p(x) \cdot p(x) = -\sum_{x \in X} p^2(x)$
 $-H(x) \leq \sum_{x \in X} p^2(x)$

$\ln x = \frac{\ln x}{\ln 2}$
 $\ln x = \ln 2 \cdot \ln x$
 $\ln x \geq \frac{1}{\ln 2} - \frac{1}{\ln 2(x)}$

- ISTEKO SE DOĐNA ISO JENSEN'S INEQUALITY
 $H(x) = -\sum_{x \in X} p(x) \ln p(x) = -E[\ln p(x)] \geq -\ln\{E[p(x)]\} \geq$
 $\geq -\ln\{E[p(x)]+1\} \geq -\frac{E[p(x)]}{E[p(x)]+1} = -\sum_{x \in X} p^2(x)$

$E[\ln(x)] < \infty$ $E[\ln(x)] \leq \ln\{E[x]\}$

 $\ln x \leq x-1$ $\gamma = \frac{1}{x}$ $x = \frac{1}{\gamma}$ $-\ln \gamma \leq \frac{1}{\gamma} - 1$

$\ln \gamma \geq 1 - \frac{1}{\gamma}$

$\gamma > 0$

 $\gamma \in \{1, 2, 3, \dots\}$
 $p(\gamma) = \{p_1, p_2, \dots\}$
 $E[\ln \gamma] \geq \sum_{\gamma \in T} (1 - \frac{1}{\gamma}) p(\gamma) = 1 - \sum_{\gamma \in T} \frac{p(\gamma)}{\gamma} = 1 - E[\frac{1}{\gamma}]$

$$H(X) = - \sum_{Y \in \mathcal{Y}} p(Y) \log p(Y)$$

$$\log p(Y) \leq p(Y) - 1$$

$$-\log p(Y) \geq 1 - p(Y) \quad p \in [0, 1]$$

$$H(X) = - \sum_{Y \in \mathcal{Y}} p(Y) \log p(Y) \geq \sum_{Y \in \mathcal{Y}} p(Y) - \sum_{Y \in \mathcal{Y}} p^2(Y) = 1 - \sum_{Y \in \mathcal{Y}} p^2(Y)$$

$$E[\ln X] \geq 1 - E\left[\frac{1}{X}\right]$$

$$H(X) \geq 1 - \sum_{Y \in \mathcal{Y}} p^2(Y) \geq 1 - 2^{-H(X)}$$

$$\log H(X) - H(X) \geq 0 ?$$

$$H(X) + 2^{-H(X)} \geq 1 \quad / \log$$

$$\log H(X) \geq H(X) ?$$

$$H(X) \geq 1 - \sum_{Y \in \mathcal{Y}} p^2(Y) \quad E[\ln X] \geq 1 - E\left[\frac{1}{X}\right]$$

$$E[\log X] \leq \log E[X]$$

$$H(X) = -p_1 \log p_1 - p_2 \log p_2 - p_3 \log p_3 - \dots$$

$$E[\log X] = p_1 \log x_1 + p_2 \log x_2 + p_3 \log x_3 + \dots$$

$$X \in \{1, 2, 3, \dots\}$$

$$E[\log X] = p_2 \log 2 + p_3 \log 3 + p_4 \log 4 + \dots$$

$$E[\log X] = \sum_{i=1}^{\infty} p_i \log i = \sum_{i=1}^{\infty} p_i \log \frac{1}{\left(\frac{1}{i}\right)}$$

$$H(X) = \sum_{i=1}^{\infty} p_i \log \frac{1}{p_i} = \sum_{i=1}^{\infty} p_i \log \frac{1}{p_i} = \sum_{i=1}^{\infty} p_i \log \frac{1}{p_i} = \sum_{i=1}^{\infty} p_i \log \frac{1}{p_i}$$

$$E[\log X] \geq \frac{1 - E\left[\frac{1}{X}\right]}{\ln 2} = \left[1 - \sum_{i=1}^{\infty} \frac{p_i}{X}\right] \cdot \frac{1}{\ln 2}$$

$$\log p_i = \left| \log \frac{1}{2^i} \right| = \log 2^{-i} = -i$$

$$H(X) = - \sum_{i=1}^{\infty} p_i \log p_i \quad E[\log X] \geq \left[1 - \sum_{i=1}^{\infty} \frac{p_i}{X}\right] \cdot \frac{1}{\ln 2}$$

$$H(X) \geq 1 - \sum_{i=1}^{\infty} p_i^2(X)$$

што е познато?

$$\sum_{i=1}^{\infty} \frac{p_i}{X} \cdot i < 1$$

$$\sum_{i=1}^{\infty} p_i^2(X) ?$$

$$\sum_{\gamma \in \Sigma} \frac{p(\gamma)}{\gamma} \geq \sum_{\gamma \in \Sigma} p^2(\gamma)$$

EMANVI SE 37
 $\gamma=1.$

V.1.11 MAKE

Multilog MIMOC

3.27.15

$$\frac{1}{\ln 2} \left[1 - \sum_{\gamma \in \Sigma} \frac{p(\gamma)}{\gamma} \right] \leq 1 - \sum_{\gamma \in \Sigma} p^2(\gamma)$$

FINAZ!!!

$[0, \dots, 1]$

$[0 \dots 1]$

$$H(x) \geq E[\log(x)]$$

JAS MISLAM DEVA E ODRADO TE DEVA AVO
 $H(x) < \infty$ TOGAIS $E[\log(x)] \leq \infty$

$$E[\log(x)] < H(x) < \infty \Rightarrow \underline{E[\log(x)] < \infty}$$

1 REARITE
 TEST VIKROSDP
 ZA $H(x)$
 $E[\log(x)]$ GO KONCU-
 VART OVA. MAKE.

$$p = \alpha \cdot \beta^i$$

VO GENERACEN
 $x \in \{i\}$ SLUCIT:
 $i = 1, 2, \dots$

$$\frac{1}{\ln 2} - \frac{3\alpha\beta}{2\ln 2} \leq 1 - 2\alpha^2\beta^2$$

$$\frac{1}{\ln 2} - \frac{\alpha}{\ln 2} \leq 1 - \alpha^2$$

$$1 - \alpha \leq \ln 2 - \ln 2 \alpha^2$$

$$\ln 2 \alpha^2 - \alpha + 1 - \ln 2 \leq 0$$

OVA E SO
 POMBI NA MOXEM
 (2.19) Ed.2

$$2 - 3\alpha\beta \leq 2\ln 2 - 4\ln 2 \alpha^2 \beta^2$$

$$4\ln 2 \alpha^2 \beta^2 - 3\alpha\beta + 2 - 2\ln 2 \leq 0$$

$$p = \alpha \cdot \beta^i$$

$$H(x) = \sum_{i=1}^{\infty} \alpha \cdot \beta^i \log(\alpha \cdot \beta^i) =$$

$$= \sum_{i=1}^{\infty} \alpha \cdot \beta^i \log \alpha + \sum_{i=1}^{\infty} \alpha \cdot \beta^i \cdot i \log \beta = \log \alpha \cdot \sum_{i=1}^{\infty} \alpha \cdot \beta^i + \log \beta \cdot \sum_{i=1}^{\infty} \alpha \cdot \beta^i \cdot i$$

$$= \log \alpha + \log \beta \cdot E[X]$$

$$E[\log x] \leq \log E[x] \leq E[x]$$

$$H(x) = \alpha \alpha^i + \beta \beta^i \cdot E[X] \Rightarrow \alpha \alpha^i + \beta \beta^i E[E[X]]$$

$$E[\alpha \alpha^i] \leq \frac{H(x) - \alpha \alpha^i}{\beta \beta^i}$$

$$H(x) \geq -\alpha \alpha^i + \beta \beta^i E[E(x)]$$

UZIKNO OD REZULTATOS VO P.R. 50

OVA KAZIVA DEKA VO GENERALIZOVANU SLUCIA ZA $p = \alpha \cdot \beta^i$ ZA $\alpha, \beta \in [0..1]$ $H(x) \geq E[\alpha \alpha^i]$

Ako izrazimo se bazira na modelu (2.30) TOGAJ :

$$E[\alpha \alpha^i] \leq \alpha \{E[X]\} \leq E[X]$$

zasto $x \in [1, 2, \dots]$

$$E[X] = \sum_{i=1}^{\infty} i \cdot p(i)$$

Ako $E[\alpha \alpha^i] < \infty$ NE MOGA DA ZNAČI DEKA $E[X] < \infty$!!

NO SREĆAK AKO $E[\alpha \alpha^i] = \text{const} \leq \infty$ TOGAJ

$$D(\gamma | q) \geq 0 \Rightarrow \sum_{i=0}^{\infty} \gamma_i \ln \frac{p_i}{q_i} \geq 0$$

$$+ \sum_{i=0}^{\infty} p_i \ln p_i - \sum_{i=0}^{\infty} p_i \ln q_i \geq 0 \Rightarrow - \sum_{i=0}^{\infty} \gamma_i \ln p_i \leq - \sum_{i=0}^{\infty} \gamma_i \ln q_i =$$

$$= \sum_{i=0}^{\infty} \gamma_i \ln \frac{q_i}{p_i} = \sum_{i=0}^{\infty} \gamma_i \ln \frac{1}{\alpha \cdot \beta^i} = - \sum_{i=0}^{\infty} \gamma_i \ln \alpha - \sum_{i=0}^{\infty} \gamma_i \ln \beta^i =$$

$$= - \ln \alpha - \ln \beta \sum_{i=0}^{\infty} i \gamma_i = - \ln \alpha - \ln \beta E[X]$$

$$\alpha, \beta \in [0..1] \Rightarrow H(x) \leq - \ln \alpha - \ln \beta E[X]$$

$$H(x) \leq \ln \alpha^{-1} + \ln \beta^{-1} E[X] = \begin{cases} \alpha=1 \\ \beta=1/2 \end{cases} = 0 + E[X]$$

$H(x) \leq E[X]$ PAK SE PORADA NA ISTO !!

$$\sum_{i=0}^{\infty} \alpha \cdot \beta^i = \alpha \frac{1}{1-\beta} = 1 \quad \alpha \sum_{i=0}^{\infty} i \beta^i = \frac{\alpha \cdot \beta}{(1-\beta)^2} = A$$

$$\frac{\alpha}{1-\beta} = \frac{1}{1-\beta} \Rightarrow \alpha = 1-\beta$$

$$\frac{\alpha}{1-\beta} = A \Rightarrow \alpha = A(1-\beta)$$

$$1-\alpha = \alpha \cdot A \Rightarrow \alpha(A+1) = 1 \Rightarrow \alpha = \frac{1}{A+1}$$

$$\beta = 1 - \frac{1}{A+1} = \frac{A}{A+1}$$

$$p = \frac{1}{A+1} \left(\frac{A}{A+1}\right)^i$$

• Ako tražimo so od nasen $E[X] = \sum_{i=0}^{\infty} i p(i) = A$

$p(i) = \frac{1}{A+1} \left(\frac{A}{A+1}\right)^i$ (1)

$E[\log(x)] \leq \log E[X] \rightarrow E[\log(x)] < \infty$

$E[X] = A < \infty$

$-H(x) = E[\log(p(x))] \leq \log E[p(x)] = \log \sum_{i=0}^{\infty} p(i) \leq E[p(x)] - 1$

$E[0 \dots 1]$ MACLOREN

(11): $-H(x) = E[\log(p(x))] \leq \log \{E[p(x)]\} \leq \log \{E[p(x)] + 1\} \leq E[p(x)]$

$E[0 \dots 1]$

U13.12 $\Rightarrow H(x) \leq -\log \alpha - \log \beta \cdot E[X] =$

$= -\log \frac{1}{1+A} - \log \frac{A}{1+A} E[X] = \log(1+A) + \log\left(\frac{1+A}{A}\right) E[X]$

$H(x) \leq \log(1+A) + \log(1+A) \cdot E[X] - \log(A) E[X]$

$= \log(1+A) [1 + E[X]] - \log A \cdot E[X] = \frac{(1+A) \log(1+A) - A \log A}{A}$

$\leq (1+A) \cdot \left(1 - \frac{1}{A}\right) - A \left(1 - \frac{1}{A}\right) = (1+A) \frac{A-1}{A} - \frac{A-1}{A} \cdot A$

$= \frac{A-1}{A}$

$H(x) \leq \frac{A-1}{A}$

• Vo ovom slučaju za bilo kakvo $A = p(x)$ geometrija e konvergentna funkcija (vidi nabra MULTIMODALNOCITOSTI MW (3.27.25))

$H(x) = E[\log p(x)] \leq \log E[p(x)] \leq \log [E[p(x)] + 1] \leq E[p(x)]$

$\leq p(E[X])$

$H(x) = \sum_{i=0}^{\infty} p \log \frac{1}{p} = - \sum_{i=0}^{\infty} \frac{1}{A+1} \left(\frac{A}{A+1}\right)^i \log \frac{1}{A+1} \left(\frac{A}{A+1}\right)^i$

for $A \rightarrow 1$ $p(x) = \left(\frac{1}{2}\right)^x$

$H(x) = -E[\log p(x)] = -E\left[\log\left(\frac{1}{2}\right)^x\right] = E[\log 2^x] = E[X]$

$$H(x) \leq \frac{(1+A) \log(1+A) - A \log A}{\text{MAXIMUM ENTROPY PR. 2.30}} = \log(1+A) + \underbrace{A \log(A/A) - A \log A}_{\text{FRIDLIEN ZA GOLEMO A (VEIC OD A \ge 4)}}$$

$$H(x) \leq \log(1+A) \leq A$$

$$\log(1+x) \leq x \quad \left. \vphantom{\log(1+x)} \right\} \text{MACLORENOV RED}$$

ZNAČI KONEČNO, AKO SE ZEME PDF-OT DA NAJDE KAKO $\text{E}[X] < \infty$ TOGA:

$$H(x) \leq A = E[X] \quad (\text{ZA } A \ge 4)$$

BIDEŽI: $E[\log(x)] \leq \log E[x] \leq E[x]$
 USLOVOT $E[\log(x)] < \infty$ ZA DA $H(x) < \infty$
 S SIGURNO IŠTAJET ŽOŽTO UŠTE ROVEČE $E[x] < \infty$.

→ OVA SEMOGAŠ VAIŽI AVO ZEME $f(x) = 2^{-x}$
 POHAJOT S SO RELATIVA ENTROPJA ($P(q||z) \ge 0$) KAKO NA PF7. SUŠTINSKIOT POHAJOT E A.9 .

2.46 AXIOMATIC DEFINITION OF ENTROPY. IF WE ASSUME CERTAIN AXIOMS FOR OUR MEASURE OF INFORMATION WE WILL BE FORCED TO USE A LOGARITMIC MEASURE SUCH AS ENTROPY. SHANNON USED THIS TO JUSTIFY HIS INITIAL DEFINITION OF ENTROPY. IN THIS BOOK WE RECAT MORE ON THE OTHER PROPERTIES OF ENTROPY RATHER THAN ITS AXIOMATIC DERIVATION TO JUSTIFY ITS USE.

IF THE SEQUENCE OF SYMBERIC FUNCTIONS $H_n(p_1, p_2, \dots, p_n)$ SATISFIES THE FOLLOWING PROPERTIES:

- NORMALIZATION: $H_2\left(\frac{1}{2}, \frac{1}{2}\right) = 1$
- CONTINUITY: $H_2(p, 1-p)$ IS CONTINUOUS FUNCTION OF "P"
- GROUPING: $H_n(p_1, p_2, \dots, p_n) = H_{n-1}(p_1 + p_2, p_3, \dots, p_n) + (p_1 + p_2) H_2\left(\frac{p_1}{p_1 + p_2}, \frac{p_2}{p_1 + p_2}\right)$

PROVE THAT H_n MUST BE OF THE FORM:

$$H_n(p_1, p_2, \dots, p_n) = -\sum_{i=1}^n p_i \log p_i \quad \text{or } = 2^3, \dots$$

GROUPING PROPERTY i.e. AXIOM:

$$H_m(p_1, p_2, \dots, p_m) = H_{m-1}(p_1+p_2, p_3, p_4, \dots, p_m) + (p_1+p_2) H_2\left(\frac{p_1}{p_1+p_2}, \frac{p_2}{p_1+p_2}\right)$$

(*)

$$H_m = \sum_{i=1}^m p_i \log \frac{1}{p_i} = p_1 \log \frac{1}{p_1} + p_2 \log \frac{1}{p_2} + \dots + p_m \log \frac{1}{p_m}$$

$$+ (p_1+p_2) \log \frac{1}{p_1+p_2} - (p_1+p_2) \log \frac{1}{p_1+p_2}$$

$$H_{m-1}(p_1+p_2, p_3, p_4, \dots, p_m) = (p_1+p_2) \log \frac{1}{p_1+p_2} + p_3 \log \frac{1}{p_3} + \dots$$

$$= (p_1+p_2) \left[\frac{p_1}{p_1+p_2} \log \frac{p_1+p_2}{p_1} + \frac{p_2}{p_1+p_2} \log \frac{p_1+p_2}{p_2} \right] =$$

$$= p_1 \log \left(\frac{p_1+p_2}{p_1} \right) + p_2 \log \left(\frac{p_1+p_2}{p_2} \right)$$

$$= p_1 \log \frac{1}{p_1} + p_2 \log \frac{1}{p_2} - p_1 \log \frac{1}{p_1+p_2} - p_2 \log \frac{1}{p_1+p_2} =$$

$$= p_1 \log \frac{p_1+p_2}{p_1} + p_2 \log \frac{p_1+p_2}{p_2} = (p_1+p_2) H_2\left(\frac{p_1}{p_1+p_2}, \frac{p_2}{p_1+p_2}\right)$$

अतः

दुहरावो!!!

$$H_m = (p_1+p_2) H_2\left(\frac{p_1}{p_1+p_2}, \frac{p_2}{p_1+p_2}\right) + H_{m-1}(p_1+p_2, p_3, p_4, \dots, p_m)$$

SOLUTION 1 SOLUTION:

- EXTEND THE GROUPING AXIOM BY INDUCTION AND PROVE THAT:

$$H_m(p_1, p_2, \dots, p_m) = H_{m-k}(p_1+p_2+\dots+p_k, p_{k+1}, \dots, p_m) + (p_1+p_2+\dots+p_k) H_k\left(\frac{p_1}{p_1+p_2+\dots+p_k}, \dots, \frac{p_k}{p_1+p_2+\dots+p_k}\right)$$

LET $f(m)$ BE ENDSIDE OF UNIFORM DISTRIBUTION OF m SYM BOLS i.e.:

$$f(m) = H_m\left(\frac{1}{m}, \frac{1}{m}, \dots, \frac{1}{m}\right)$$

- For ANY TWO INTEGERS r AND s :

$$f(r \cdot s) = f(r) + f(s)$$

- For RATIONALS: $y = \frac{r}{s}$

$$H_2(y, 1-y) = -y \log y - (1-y) \log(1-y)$$

$$f(n) = H_n\left(\frac{1}{n}, \frac{1}{n}, \dots, \frac{1}{n}\right) = -\sum_{i=1}^n p_i \log p_i =$$

$$= \sum_{i=1}^n \frac{1}{n} \log \frac{1}{n} = n \cdot \frac{1}{n} \log \frac{1}{n} = \log \frac{1}{n}$$

DA HVA AKO
DA ZNACS
DEFINICATA NA
ENTROPIJA NA
ZAPISIVA CIMA
NE ZA ZNACS.

$$\boxed{n = r \cdot s} \Rightarrow f(r \cdot s) = \log(r \cdot s) = \log r + \log s$$

$$H_2\left(\frac{r}{s}, 1 - \frac{r}{s}\right) = \frac{r}{s} \log \frac{s}{r} + \left(1 - \frac{r}{s}\right) \log \frac{1}{s-r}$$

$$= \frac{r}{s} \log \frac{s}{r} + \log \frac{s}{s-r} - \frac{r}{s} \log \frac{s}{s-r}$$

$$= \frac{r}{s} \log \frac{s}{r} + \log \frac{s}{s-r} = \frac{r}{s} \log \left(\frac{s-r}{s}\right) + \log \frac{s}{s-r}$$

$$= \frac{r}{s} \cdot \log \frac{s}{r} \cdot \frac{s-r}{s} + \log \frac{s}{s-r} = \frac{s-r}{s} \log \frac{s}{s-r}$$

$$H\left(\frac{r_2}{s_2}\right) = H\left(1 - \frac{r_2}{s_2}, \frac{r_2}{s_2}\right) = H\left(\frac{q_1 + r_2 - r_2}{s_2}, \frac{r_2}{s_2}\right) = H\left(\frac{q_1}{s_2}, \frac{r_2}{s_2}\right)$$

$$\boxed{S_k = \sum_{i=1}^k p_i}$$

$$\boxed{H(2) = H_2\left(\frac{1}{2}, 1 - \frac{1}{2}\right)}$$

$$\boxed{S_3 = S_2 + p_3}$$

$$H_n(p_1, \dots, p_n) = H_{n-1}(S_2, p_2, \dots, p_n) + S_2 H\left(\frac{p_1}{S_2}\right)$$

$$= H_{n-2}(S_3, p_3, \dots, p_n) + S_3 H\left(\frac{p_2}{S_3}\right) + S_2 H\left(\frac{p_1}{S_2}\right)$$

$$\vdots = H_{n-(k-1)}(S_k, p_{k+1}, \dots, p_n) + \sum_{i=2}^k S_i H\left(\frac{p_i}{S_i}\right) \quad \textcircled{\$}$$

- AKA SE UOVEDI REKURZIVNO GORNJE AKSIOMATA NA:

$$H_k\left(\frac{p_1}{S_k}, \frac{p_2}{S_k}, \dots, \frac{p_k}{S_k}\right) = H_2\left(\frac{S_{k-1}}{S_k}, \frac{p_k}{S_k}\right) + \sum_{i=2}^{k-1} \frac{S_i}{S_k} H\left(\frac{p_i/S_k}{S_i/S_k}\right)$$

$$= H_2\left(\frac{S_k}{S_k}\right) + \sum_{i=2}^k \frac{S_i}{S_k} H\left(\frac{p_i}{S_i}\right) = \frac{1}{S_k} \sum_{i=2}^k H\left(\frac{p_i}{S_i}\right)$$

ZNAČI: $H_k \left(\frac{p_1}{s_k}, \frac{p_2}{s_k}, \dots, \frac{p_k}{s_k} \right) = \frac{1}{s_k} \sum_{i=1}^k s_i \ln \left(\frac{p_i}{s_i} \right)$

Ako se zameni uo (8) se dobiva

$$H(p_1, p_2, \dots, p_n) = H_{n-k+1} \left(s_k, p_{k+1}, \dots, p_n \right) + s_k H_k \left(\frac{p_1}{s_k}, \frac{p_2}{s_k}, \dots, \frac{p_k}{s_k} \right)$$

- So ova je dokazana proširena forma na slučajna na garantu.

• CONSIDER: $f(n, y) = H_{n+1} \left(\frac{1}{ny}, \frac{1}{ny}, \dots, \frac{1}{ny} \right) =$

$$= H_{n+1-y+1} \left(n \frac{1}{ny}, \frac{1}{ny}, \dots, \frac{1}{ny} \right) + \frac{1}{ny} H_n \left(\frac{1}{y}, \frac{1}{y}, \dots, \frac{1}{y} \right)$$

$$= H_{n+1-y+1} \left(\frac{1}{n}, \frac{1}{ny}, \dots, \frac{1}{ny} \right) + \frac{1}{n} H_n \left(\frac{1}{y}, \frac{1}{y}, \dots, \frac{1}{y} \right)$$

$$= H_{n+1-2y+1} \left(\frac{1}{n}, \frac{1}{n}, \frac{1}{ny}, \dots, \frac{1}{ny} \right) + \frac{2}{n} H_n \left(\frac{1}{y}, \frac{1}{y}, \dots, \frac{1}{y} \right)$$

$$\dots$$

$$= H_n \left(\frac{1}{n}, \frac{1}{n}, \dots, \frac{1}{n} \right) + H_n \left(\frac{1}{y}, \frac{1}{y}, \dots, \frac{1}{y} \right)$$

$$= f(n) + f(y) \quad \text{DOKAZANO !!!}$$

- Otvor automatski sledi: $f(n^k) = k \cdot f(n)$

UŠTE EDNO INTERESNO SVOJSTVO KOJE ODMA SE GLEDA ADO ZNAČI: DETA ENTROPIJA I DETA KANA COGAMINAMKI:

$$H(n^k) = \sum_{i=1}^{n^k} \frac{1}{n^k} \ln n^k = n^k \cdot \frac{1}{n^k} \ln n^k = k \frac{\ln n}{H(n)}$$

• $H_2(1, 0) = H(1) = 0$ $(p_1 + p_2 = 1)$

- PROOF:

$$H_3(p_1, p_2, 0) = H_2(p_1, p_2) + p_2 H_2(1, 0) =$$

$$= H_2(p_1 + p_2, 0) + (p_1 + p_2) H_2 \left(\frac{p_1}{p_1 + p_2}, \frac{p_2}{p_1 + p_2} \right) =$$

$$= H_2(1, 0) + H_2(p_1, p_2)$$

$$\Rightarrow p_2 H_2(1, 0) = H_2(1, 0) \quad \text{za } \forall p_2 \Rightarrow H_2(1, 0) = 0$$

① $f(n+1) - f(n) \rightarrow 0$ AS $n \rightarrow \infty$ (DA SE DOKAZE !!!)

$$f(n+1) = H\left(\frac{1}{n+1}, \frac{1}{n+1}, \dots, \frac{1}{n+1}\right) = H_{\underbrace{n+1-n}_{2}}\left(\frac{1}{n+1}, \frac{1}{n+1}\right) + \frac{1}{n+1} H_n\left(\frac{1}{n+1}, \frac{1}{n+1}, \dots, \frac{1}{n+1}\right) = H_2\left(\frac{1}{n+1}, \frac{1}{n+1}\right) + \frac{1}{n+1} H_n\left(\frac{1}{n+1}, \frac{1}{n+1}, \dots, \frac{1}{n+1}\right) = H_2\left(\frac{1}{n+1}, \frac{1}{n+1}\right) + \frac{1}{n+1} f(n)$$

$$f(n+1) = h\left(\frac{1}{n+1}\right) + \frac{1}{n+1} f(n)$$

$$f(n+1) - f(n) = h\left(\frac{1}{n+1}\right) - \frac{1}{n+1} f(n)$$

$$f(n+1) - \frac{1}{n+1} f(n) = h\left(\frac{1}{n+1}\right)$$

$$\lim_{n \rightarrow \infty} \frac{(n+1)f(n+1) - nf(n)}{n+1} = \lim_{n \rightarrow \infty} h\left(\frac{1}{n+1}\right)$$

$$\lim_{n \rightarrow \infty} \frac{(n+1)f(n+1) - nf(n)}{n+1} = \frac{f(n) - f(n)}{h\left(\frac{1}{n+1}\right)} = 0 \Rightarrow h\left(\frac{1}{n+1}\right) = 0$$

Let us define:

$$a_{n+1} = f(n+1) - f(n)$$

$$b_n = h\left(\frac{1}{n}\right)$$

$$a_{n+1} = h\left(\frac{1}{n+1}\right) + \frac{1}{n+1} f(n) - f(n) = h\left(\frac{1}{n+1}\right) + \frac{nf(n) - (n+1)f(n)}{n+1} = h\left(\frac{1}{n+1}\right) - \frac{f(n)}{n+1}$$

$$a_{n+1} = -\frac{1}{n+1} f(n) + b_{n+1}$$

$$a_{n+1} = -\frac{1}{n+1} \sum_{i=2}^n a_i + b_{n+1}$$

$$a_2 = f(2) - f(1) \quad a_3 = f(3) - f(2) \quad a_4 = f(4) - f(3)$$

$$a_n = f(n) - f(n-1) \quad a_2 + a_3 + a_4 + \dots + a_n = f(2) - f(1) + f(3) - f(2) + f(4) - f(3) + \dots + f(n) - f(n-1) = -f(1) + f(n) = f(n) - f(1)$$

$$(n+1) \cdot a_{n+1} = - \sum_{i=2}^n a_i + (b_{n+1}) \cdot (n+1) \quad (n+1) b_{n+1} = (n+1) a_{n+1} + \sum_{i=2}^n a_i$$

• Summing over n

$$\sum_{n=2}^N n b_n = \sum_{n=2}^N n a_n + \sum_{n=2}^N \sum_{i=2}^n a_i =$$

$$= \sum_{n=2}^N (n a_n + a_{n-1} + a_{n-2} + \dots + a_2)$$

$$\boxed{n+1 = m} \quad m b_m = m a_m + \sum_{i=2}^{m-1} a_i \quad \left| \sum_{m=2}^N \right.$$

$$\sum_{m=2}^N m b_m = \sum_{m=2}^N m a_m + \sum_{m=2}^N \sum_{i=2}^{m-1} a_i =$$

$$= \sum_{m=2}^N m a_m + a_{m-1} + a_{m-2} + \dots + a_2 \Rightarrow$$

$$\boxed{\sum_{n=2}^N n b_n = \sum_{n=2}^N (n a_n + a_{n-1} + a_{n-2} + \dots + a_2)}$$

$$\sum_{n=2}^N n a_n + \sum_{n=2}^N a_{n-1} + \sum_{n=2}^N a_{n-2} + \dots + \sum_{n=2}^N a_2$$

$\underbrace{\sum_{n=2}^N a_2}_{(N-1) a_2}$

$$\sum_{n=2}^N a_{n-2} = \left| \begin{array}{l} n=4-2 \quad n=4-2 \\ n=2 \quad n=0 \\ n=N \quad n=N-2 \end{array} \right| = \sum_{n=0}^{N-2} a_n$$

$$\sum_{n=2}^N a_{n-1} = \left| \begin{array}{l} n=4-1 \\ n=2 \quad n=1 \\ n=N \quad n=N-1 \end{array} \right| = \sum_{n=1}^{N-1} a_n$$

$$\sum_{n=2}^4 (n a_n + a_{n-1} + a_{n-2} + \dots + a_2) = 2a_2 + 3a_3 + a_2 + 4a_4 + a_3 + a_2 = 4a_2 + 4a_3 + 4a_4$$

$$\sum_{n=2}^5 (n a_n + a_{n-1} + a_{n-2} + \dots + a_2) = 2a_2 + 3a_3 + a_2 + 4a_4 + a_3 + a_2 + 5a_5 + a_4 + a_3 + a_2 = 5a_2 + 5a_3 + 5a_4 + 5a_5$$

$$\sum_{n=2}^N (4a_n + a_{n-1} + a_{n-2} + \dots + a_2) = N \cdot \sum_{i=2}^N a_i \quad / \quad \sum_{n=1}^N n = \frac{N(N+1)}{2}$$

$$\begin{array}{l} S = 1 + 2 + \dots + (n-1) + n \\ S = n + (n-1) + \dots + 2 + 1 \end{array} \quad | \quad +$$

$$2S = (n+1) + (n+1) + \dots + (n+1) \quad \boxed{S = \frac{n \cdot (n+1)}{2}}$$

$$\sum_{n=2}^N n b_n = N \sum_{i=2}^N a_i \quad / \quad \sum_{n=1}^N n = \frac{(N+1)N}{2}$$

$$\begin{array}{l} S = 2 + 3 + \dots + n \\ S = n + (n-1) + \dots + 2 \end{array} \quad | \quad +$$

$$2S = (n+2) + (n+2) + \dots + (n+2) \quad \boxed{S = \frac{(n-1) \cdot (n+2)}{2}}$$

$$\sum_{n=2}^N n = \frac{(N-1)(N+2)}{2}$$

$$\frac{\sum_{n=2}^N n b_n}{\sum_{n=1}^N n} = \frac{2}{N+1} \sum_{i=2}^N a_i$$

$$\boxed{b_n = n \left(\frac{1}{n} \right)}$$

$$n \left(\frac{1}{n} \right) = \left(1 - \frac{1}{n} \right) \ln \left(1 - \frac{1}{n} \right) + \frac{1}{n} \ln \left(\frac{1}{n} \right)^{-1} =$$

$$= \frac{n-1}{n} \ln \left(\frac{n-1}{n} \right) + \frac{1}{n} \ln n = \frac{n-1}{n} \ln \frac{n}{n-1} + \frac{1}{n} \ln n$$

$$= \frac{n-1}{n} \ln n - \frac{n-1}{n} \ln(n-1) + \frac{1}{n} \ln n = \left(\frac{n-1}{n} + \frac{1}{n} \right) \ln n - \frac{n-1}{n} \ln(n-1)$$

$$= \ln n - \frac{n-1}{n} \ln(n-1) = 0 \quad \log_1 x > 0 \in \mathbb{R} \quad \text{MADE VIKTA TAKA!}$$

$$\Rightarrow b_n \rightarrow 0 \quad \text{IF} \quad n \rightarrow \infty \quad \Rightarrow \quad \sum_{n=2}^N n b_n \rightarrow 0$$

$$\Rightarrow a_{n+1} = - \frac{1}{n+1} \sum_{i=2}^n a_i + b_{n+1} \quad \Rightarrow a_{n+1} \rightarrow 0 \quad \Rightarrow a_{n+1} = f(n+1) - f(n) = 0 \quad \text{DOMEN!!!}$$

LEMA 2.0.1 LET THE FUNCTION $f(n)$ SATISFY THE FOLLOWING ASSUMPTIONS:

- $f(mn) = f(m) + f(n)$
- $\lim_{n \rightarrow \infty} [f(n+1) - f(n)] = 0$
- $f(2) = 1$

THEN THE FUNCTION $f(n) = \log_2 n$.

PROOF: LET p BE ARBITRARY PRIME NUMBER AND
LET:

$$g(n) = f(n) - \frac{f(p) \cdot \log n}{\log p}$$

$$g(mn) = f(mn) - \frac{f(p) \cdot \log(mn)}{\log p} =$$

$$= f(mn) - \frac{f(p) \log n}{\log p} - \frac{f(p) \log(m)}{\log p} =$$

$$f(mn) = f(m) + f(n)$$

$$= f(m) - \frac{f(p) \log n}{\log p} + f(n) - \frac{f(p) \log(m)}{\log p} =$$

$$= g(m) + g(n)$$

ZNAČI $g(n)$ JA UPOKUVATA ISTOVA PREPOSTAVKA
OO $g(p)$.

$$g(p) = f(p) - \frac{f(p) \cdot \log p}{\log p} = f(p) - f(p) = 0$$

$$a_n = g(n+1) - g(n) = f(n+1) - \frac{f(p) \cdot \log(n+1)}{\log p} - f(n) + \frac{f(p) \log n}{\log p}$$

$$= f(n+1) - f(n) - \frac{f(p)}{\log p} \log \frac{n+1}{n}$$

$$\lim_{n \rightarrow \infty} [f(n+1) - f(n)] - \lim_{n \rightarrow \infty} \left[\frac{f(p)}{\log p} \log \frac{n+1}{n} \right] = 0$$

$$\Rightarrow \boxed{\lim_{n \rightarrow \infty} a_n = 0}$$

$$= \lim_{n \rightarrow \infty} [c \cdot x] = c \lim_{n \rightarrow \infty} \frac{n+1}{n} = 0$$

FOR AN INTEGER n , DEFINE

$$n^{(k)} = \left\lfloor \frac{n}{p^k} \right\rfloor$$

$$n^{(k)} < \frac{n}{p^k}$$

$$n = p \cdot n^{(1)} + l$$

$$0 \leq l < p$$

$$g(p) = 0 \quad g(p \cdot n^{(k)}) = g(n^{(k)})$$

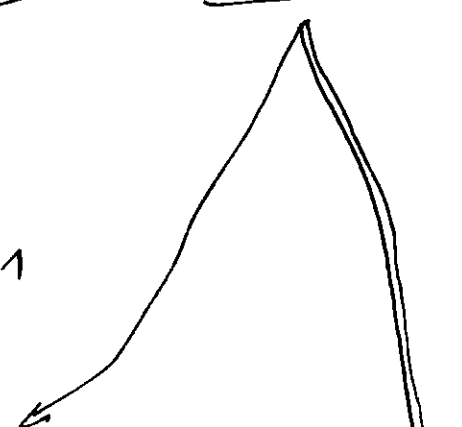
$$g(p \cdot n^{(k)}) = g(p) + g(n^{(k)}) = g(n^{(k)})$$

$$g(n) = g(n^{(k)}) - g(p \cdot n^{(k)}) + \sum_{i=1}^{n^{(k)}} g(i) = g(n^{(k)}) + \sum_{i=p \cdot n^{(k)}}^{n-1} a_i$$

$$g(n) = g(n^k) + \sum_{i=1}^k \left(\sum_{i=p^i n^{(k)}}^{n^{(k)}} a_i \right)$$

$$n^{(k)} \leq \frac{n}{p^k}$$

$$k = \left\lfloor \frac{\log n}{\log p} \right\rfloor + 1$$



$$a_i = g(i+1) - g(i)$$

$$\sum_{i=p \cdot n^{(k)}}^{n-1} a_i = a_{p \cdot n^{(k)}} + \dots + a_{n-1} = g(p \cdot n^{(k)} + 1) - g(p \cdot n^{(k)}) + g(p \cdot n^{(k)} + 2) - g(p \cdot n^{(k)} + 1) + \dots + g(n) - g(n-1) = g(p \cdot n^{(k)} + 1) - g(n-1)$$

• AKO DEFINICIJAMA ZA $a_i = \epsilon : (kMO NA P. 17)$

$$a_i = g(i) - g(i-1)$$

$$\sum_{i=p \cdot n^{(k)}}^{n-1} a_i = g(p \cdot n^{(k)}) - g(p \cdot n^{(k)} - 1) + g(p \cdot n^{(k)} + 1) - g(p \cdot n^{(k)}) + g(n-1) - g(n-2) = g(p \cdot n^{(k)}) - g(n-1)$$

• SVEOK TUDA SO OVAJ DEFINICIJA

$$\sum_{i=p \cdot n^{(k)}}^{n-1} a_i = g(p \cdot n^{(k)} + 1) - g(p \cdot n^{(k)}) + g(p \cdot n^{(k)} + 2) - g(p \cdot n^{(k)} + 1) + g(p \cdot n^{(k)} + 3) - g(p \cdot n^{(k)} + 2) + \dots + g(n) - g(n-1)$$

$$\sum_{i=p \cdot n^{(k)}}^{n-1} a_i = g(n) - g(p \cdot n^{(k)})$$

DAKLENO!!!

- Kano isto do definicijama $h^{(k)}$ od h se definišu

$$g(h^{(1)}) = g(h^{(2)}) + \sum_{i=p_{h^{(2)}}}^{h^{(1)}-1} a_i$$

$$g(h) = g(h^{(1)}) + \sum_{i=p_{h^{(1)}}}^{h^{(1)}-1} a_i = g(h^{(2)}) + \sum_{i=p_{h^{(2)}}}^{h^{(1)}-1} a_i + \sum_{i=p_{h^{(1)}}}^{h^{(1)}-1} a_i$$

$$g(h) = g(h^{(k)}) + \sum_{i=p_{h^{(k)}}}^{h^{(k)}-1} a_i + \dots + \sum_{i=p_{h^{(2)}}}^{h^{(1)}-1} a_i + \sum_{i=p_{h^{(1)}}}^{h^{(1)}-1} a_i =$$

$$= g(h^{(k)}) + \sum_{j=1}^k \sum_{i=p_{h^{(j)}}}^{h^{(j)}-1} a_i$$

$$h^{(k)} = \left\lfloor \frac{h}{p^k} \right\rfloor$$

$$h^{(k)} \leq \frac{h}{p^k}$$

$$h = p^k \cdot h^{(k)} + c$$

$$0 \leq c < p^k$$

$$k = \left\lfloor \frac{\log h}{\log p} \right\rfloor + 1$$

$$\left\lfloor \frac{\log h}{\log p} \right\rfloor = k-1$$

$$k-1 = \left\lfloor \frac{\log h}{\log p} \right\rfloor \leq \frac{\log h}{\log p}$$

$$\frac{\log h}{\log p} \geq k-1$$

$$\log h \geq \log p^{k-1}$$

$$h \geq p^{k-1}$$

$$h^{(k)} = \left\lfloor \frac{p^{k-1}}{p^k} \right\rfloor = \left\lfloor \frac{1}{p} \right\rfloor = 0 \quad (p > 1)$$

$$h^{(k)} = 0 \Rightarrow g(h^{(k)}) = g(0) = 0$$

$$g(h) = \sum_{i=1}^{b_h} a_i \quad b_h \leq p \cdot \left\lfloor \frac{\log h}{\log p} + 1 \right\rfloor$$

Since $a_n \rightarrow 0 \Rightarrow \frac{g(h)}{\log h} \rightarrow 0$ $g(h)$ HAS AT MOST

$O(\log_2 h)$ TERMS

$$\lim_{h \rightarrow \infty} \frac{g(h)}{\log h} = \frac{f(p)}{\log p}$$

$$g(h) = f(h) - \frac{f(p) \log h}{\log p}$$

$$\frac{g(h)}{\log h} = \frac{f(h)}{\log h} - \frac{f(p)}{\log p}$$

$$\lim_{h \rightarrow \infty} \frac{f(h)}{\log h} = \frac{f(1)}{\log p}$$

THIRD AXIOM

$$\frac{f(p)}{\log p} = c \quad \lim_{h \rightarrow \infty} \frac{f(h)}{\log h} = 1 \Rightarrow$$

$$\boxed{f(p) = \log p}$$

$$N = p_1 p_2 \dots p_k$$

$$f(N) = \sum f(p_i) = \sum \log p_i = \log N$$

SAHA DA WARE DUNA KECHU PRIME NUMBERS NOTES
 DA ONSEI DICO 400 INTEGER POSITIVE BROT.
 (FUNDAMENTAL THEOREM OF ARITHMETICS)

- SO OVA LEMATA E DOWAZANA
- DOWAZ AND SE PREKOSTANI DEKA $f(n)$ E MONOTONA.
- ONLY FUNCTION $f(n)$ SUCH THAT $f(mn) = f(m) + f(n)$ FOR ALL INTEGERS m, n IS OF THE FORM $f(n) = \log_a(n)$ FOR SOME BASE "a".

LET $c = f(2)$. $f(4) = f(2) + f(2) = f(2 \cdot 2) = \underline{\underline{2c}}$

$$f(2^k) = k \cdot f(2) = k \cdot c = c \cdot \log 2^k$$

- EXTENSION FOR INTEGERS THAT ARE NOT POWER OF 2.

FOR ANY INTEGER m , LET $r > 0$, BE ANOTHER INTEGER AND LET $2^k \leq m < 2^{k+1}$

$$kc \leq r f(m) < (k+1)c \quad \underline{\underline{c \frac{k}{r} \leq f(m) < c \frac{k+1}{r}}}$$

$$\log 2^k \leq \log m < \log 2^{k+1} \quad \underline{\underline{\frac{k}{r} \leq \log m < \frac{k+1}{r}}}$$

$$\frac{ck}{r} \leq c \log m \leq c \frac{k+1}{r}$$

$$\frac{k}{r} \leq \frac{f(m)}{c} \leq \frac{k+1}{r}$$

$$\frac{k}{r \cdot c} \leq \frac{\log m}{c} \leq \frac{k+1}{r \cdot c}$$

$$|x| = \begin{cases} x & x > 0 \\ -x & x < 0 \end{cases}$$

$$\left| \frac{f(m) - \log m}{c} \right| = \begin{cases} \frac{f(m) - \log m}{c} & f(m) - \log m > 0 \\ \frac{\log m - f(m)}{c} & f(m) - \log m < 0 \end{cases}$$

$$\frac{k}{r} - \frac{k}{rc} \leq f(m) - \frac{\log m}{c} \leq \frac{k+1}{r} - \frac{k+1}{rc}$$

$$L = \frac{kc - k}{rc} \quad R = \frac{(k+1)c - (k+1)}{rc} = \frac{c(k+1) - (k+1)}{rc}$$

$$= \frac{rc - (k+1)}{rc}$$

$$\left| \frac{f(m) - \log_2 m}{c} \right| < \frac{1}{r} \Rightarrow \text{MOKA } f(m) = \frac{\log_2 m}{c}$$

$c = 1$ $\exists \text{ } r > 0$ $f(2) = \frac{\log_2(2)}{c} = 1$ $c = \log_2(2) = 1$

- POUZITIE DOKAZA:

$$f(m) = H_m \left(\frac{1}{m}, \frac{1}{m}, \dots, \frac{1}{m} \right) = \ln m$$

- DA POUZIJEME DOKAZ:

$$H_2(q, 1-q) = -q \ln q - (1-q) \ln(1-q) \quad \text{(*)}$$

- BT EXTENDED GROUPING NOTION!

$$H(\gamma_1, \gamma_2, \dots, \gamma_n) = H_{m-k+n}(s_k, \gamma_{k+1}, \dots, \gamma_n) + s_k H_k \left(\frac{\gamma_1}{s_k}, \frac{\gamma_2}{s_k}, \dots, \frac{\gamma_k}{s_k} \right)$$

$$f(s) = H_s \left(\frac{1}{s}, \dots, \frac{1}{s} \right) = H \left(\frac{1}{s}, \frac{1}{s}, \dots, \frac{1}{s}, \frac{s-v}{s} \right) + \frac{s-v}{s} f(s-v) \quad \text{(**)}$$

$$s_k = (s-v) \cdot \frac{1}{s}$$

$$f(s-v) = H \left(\frac{1}{\frac{s-v}{s}}, \frac{1}{\frac{s-v}{s}}, \dots, \frac{1}{\frac{s-v}{s}} \right) = H_{s-v} \left(\frac{1}{s-v}, \frac{1}{s-v}, \dots, \frac{1}{s-v} \right)$$

$$\begin{aligned} \text{(**)} &= H_2 \left(\frac{v}{s}, \frac{s-v}{s} \right) + \frac{v}{s} H_v \left(\frac{1}{\frac{v}{s}}, \frac{1}{\frac{v}{s}}, \dots, \frac{1}{\frac{v}{s}} \right) + \frac{s-v}{s} f(s-v) \\ &= H_2 \left(\frac{v}{s}, \frac{s-v}{s} \right) + \frac{v}{s} H_v \left(\frac{1}{\frac{v}{s}}, \frac{1}{\frac{v}{s}}, \dots, \frac{1}{\frac{v}{s}} \right) + \frac{s-v}{s} f(s-v) \end{aligned}$$

$$= H_2 \left(\frac{v}{s}, \frac{s-v}{s} \right) + \frac{v}{s} f(v) + \frac{s-v}{s} f(s-v)$$

$$f(s) = \ln s \Rightarrow H_2 \left(\frac{v}{s}, \frac{s-v}{s} \right) = \ln s - \frac{v}{s} f(v) - \frac{s-v}{s} f(s-v) \quad \text{(***)}$$

$$= \ln s - \frac{v}{s} \ln v - \frac{s-v}{s} \ln(s-v) = \ln s - \frac{v}{s} \ln v - \ln(s-v) + \frac{v}{s} \ln(s-v)$$

~~$$\begin{aligned} \text{(***)} &= H_2 \left(\frac{v}{s}, \frac{s-v}{s} \right) + \frac{v}{s} f(v) + \frac{s-v}{s} f(s-v) \\ \ln s - \frac{v}{s} \ln v - \frac{s-v}{s} \ln(s-v) &= f(s) + \frac{v}{s} f(v) + \frac{s-v}{s} f(s-v) \\ \ln s - \frac{v}{s} \ln v - \frac{s-v}{s} \ln(s-v) &= \left(1 - \frac{v}{s}\right) \ln s - \left(1 - \frac{v}{s}\right) \ln(s-v) \end{aligned}$$~~

$$= \ln s - \frac{v}{s} \ln v + \frac{v}{s} \ln s - \frac{v}{s} \ln v - \left(1 - \frac{v}{s}\right) \ln(s-v) =$$

$$= \ln s - \frac{2v}{s} \ln v - \left(1 - \frac{v}{s}\right) \ln(s-v) =$$

$$= -\frac{v}{s} \ln \frac{v}{s} + \left(1 - \frac{v}{s}\right) \ln s - \left(1 - \frac{v}{s}\right) \ln(s-v) =$$

NE HOJE KOEFICIENT DA GO POTICAS!!!

TUKA E GREŠKA !!

$$= -\frac{v}{s} \ln \frac{v}{s} - \left(1 - \frac{s}{v}\right) \ln \frac{1}{s} - \left(1 - \frac{v}{s}\right) \ln(s-v) = \textcircled{1}$$

$$-\left(1 - \frac{s}{v}\right) = \left(\frac{s}{v} - 1\right) = \frac{s-v}{v} \quad - \left(1 - \frac{v}{s}\right) = \left(\frac{v}{s} - 1\right) = \frac{v-s}{s}$$

$$\textcircled{1} = -\frac{v}{s} \ln \frac{v}{s} + \frac{s-v}{v} \ln \frac{1}{s} + \frac{v-s}{s} \ln(s-v) =$$

$$= -\frac{v}{s} \ln \frac{v}{s} + \frac{v-s}{v} \ln s + \frac{v-s}{s} \ln(s-v) =$$

$$= -\frac{v}{s} \ln \frac{v}{s} + \left(1 - \frac{s}{v}\right) \ln s + \left(\frac{v}{s} - 1\right) \ln(s-v)$$

$$p = \frac{v}{s} \quad \left(1 - \frac{1}{p}\right) \ln s + (p-1) \ln(s-v)$$

IMA GRESKA
VO IZVEDU-
VANJE TO !!!
000

~~$$\left(1 - \frac{1}{p}\right) \ln s + (p-1) \ln(s-v)$$~~

~~$$h_2\left(\frac{v}{s} \mid \frac{s-v}{s}\right) = h_2\left(\frac{v}{s} \mid \frac{s-v}{s}\right) + \frac{v}{s} h_2\left(\frac{v}{s} \mid \frac{v}{s}\right) + \frac{s-v}{s} h_2\left(\frac{s-v}{s} \mid \frac{s-v}{s}\right)$$~~

~~$$h_2\left(\frac{v}{s} \mid \frac{s-v}{s}\right) = \left(1 - \frac{s}{v}\right) h_2\left(\frac{v}{s} \mid \frac{v}{s}\right) + \left(1 - \frac{v}{s}\right) h_2\left(\frac{s-v}{s} \mid \frac{s-v}{s}\right)$$~~

~~$$= \left(1 - \frac{s}{v}\right) \ln \frac{v}{s} + \left(1 - \frac{v}{s}\right) \ln \frac{s-v}{s} =$$~~

~~$$\ln \frac{v}{s} \ln \frac{v}{s} + \frac{v}{s} \ln \frac{v}{s} + \ln s - \frac{s}{v} \ln s + \left(1 - \frac{v}{s}\right) \ln \frac{s-v}{s}$$~~

~~$$= -\frac{v}{s} \ln \frac{v}{s} + \frac{v}{s} \ln v - \frac{v}{s} \ln s + \ln s - \frac{s}{v} \ln s + \left(1 - \frac{v}{s}\right) \ln \frac{s-v}{s}$$~~

~~$$\textcircled{1} = \ln s \left(1 - \frac{v}{s}\right) - \frac{v}{s} \ln \frac{v}{s} = \ln s \left(\frac{s-v}{s}\right) - \frac{v}{s} \ln \frac{v}{s}$$~~

$$h_2\left(\frac{v}{s} \mid \frac{s-v}{s}\right) = \ln s - \frac{v}{s} \ln v - \frac{s-v}{s} \ln(s-v) = \textcircled{2}$$

$$= \ln s - \frac{v}{s} \ln v + \frac{v}{s} \ln s - \frac{v}{s} \ln s - \frac{s-v}{s} \ln(s-v) =$$

$$= \ln s - \frac{v}{s} \ln v - \frac{v}{s} \ln \left(\frac{v}{s}\right) - \frac{v}{s} \ln s - \frac{s-v}{s} \ln(s-v) =$$

$$= \ln s \left(1 - \frac{v}{s}\right) - \frac{v}{s} \ln \frac{v}{s} - \frac{s-v}{s} \ln(s-v) =$$

$$= -\frac{v}{s} \ln \frac{v}{s} - \left(1 - \frac{v}{s}\right) \ln \frac{1}{s} - \left(1 - \frac{v}{s}\right) \ln(s-v) =$$

$$= -\frac{v}{s} \ln \frac{v}{s} - \left(1 - \frac{v}{s}\right) \ln \frac{s-v}{s} = -\frac{v}{s} \ln \frac{v}{s} - \left(1 - \frac{v}{s}\right) \ln \left(1 - \frac{v}{s}\right)$$

POKREMO

• ASSUME THAT NZL. (4) HOLDS FOR $n = n-1$

$$\begin{aligned}
 H_n(r_1, r_2, \dots, r_n) &= \frac{H_{n-1}(r_1+r_2, r_3, \dots, r_n) + (r_1+r_2) \log_2 \left(\frac{r_1}{r_1+r_2}, \frac{r_2}{r_1+r_2} \right)}{1} \\
 &= - \sum_{i=3}^n r_i \log_2 r_i + (r_1+r_2) \log_2 \left(\frac{1}{r_1+r_2} \right) + (r_1+r_2) \cdot \frac{r_1}{r_1+r_2} \log_2 \frac{r_1+r_2}{r_1} \\
 &\quad + (r_1+r_2) \cdot \frac{r_2}{r_1+r_2} \log_2 \frac{r_1+r_2}{r_2} = - \sum_{i=2}^n r_i \log_2 r_i + (r_1+r_2) \log_2 \frac{1}{r_1+r_2} \\
 &\quad + \frac{r_1 \log_2 (r_1+r_2)}{1} - \frac{r_1 \log_2 r_1}{1} + \frac{r_2 \log_2 (r_1+r_2)}{1} - \frac{r_2 \log_2 r_2}{2} \\
 &= - \sum_{i=3}^n r_i \log_2 r_i - (r_1+r_2) \log_2 (r_1+r_2) + (r_1+r_2) \log_2 (r_1+r_2) + \sum_{i=1}^2 r_i \log_2 \frac{1}{r_i} \\
 &= \sum_{i=1}^n r_i \log_2 r_i \quad \text{POKAZANO SO INDUKCIJA!!!}
 \end{aligned}$$

DVAJA ČENI OD INDUKCIJA SE ISKORISTI
 1.) POKAZANO VLEČOŠT $n=2$
 2.) ANO VLEČI ZA $n=n-1$ VLEČI 1 ZA $n=n$.

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WORDY OF MISORTED FLE. A DECK OF n^2

CARDS IN ORDER $1, 2, \dots, n$ IS PROVIDED. ONE CARD IS REMOVED AT RANDOM, AND THEN REPLACED AT RANDOM. WHAT IS ENTROPY OF THE RESULTING DECK.

$$H(n) = \sum_{i=1}^n \frac{1}{n} \log_2 n = n \frac{1}{n} \log_2 n = \log_2 n$$

$(n-1)$ -PROVI
 .1.2.3.4.

PLAVIČI
 ČIŽI NA NZT

1 2 3 4			
2 1 3 4*	2 1 3 4*	3 2 1 4	4 2 3 1
3 2 1 4	1 3 2 4*	1 3 2 4*	1 4 3 2
4 2 3 1	1 4 3 2	1 2 4 3*	1 2 4 3*

(1 2 3) 2 1 3* 3 2 1
 2 1 3* 1 3 2* 1 3 2*
 3 2 1
 VUKRO IMA

6 MISORTED KOMBINACII
 KOI SE POKREKNO VELOZAVANI

4.3 = 12 MISORTED KOMBINACII.
 ADŽEŠENT $2(n-1) = 6$

$n \cdot (n-1)$ BILU KOD OD n -TE BLOEVI MOZE DA GO STAVI NA $(n-1)$ MESTA

$2(n-1) \Rightarrow$ OD KAKOOTE $n(n-1)$ MA $(n-1)$ SOZEMTI PAKOVI I SGLUOJ SOZEMEN KESS MA 2 DUKLIKA TI.

$$P(x) = \begin{cases} \frac{1}{4^2} & \text{CASE I } x(i)=j; x(j)=i; x(k)=k \quad |i-j| > 1 \\ \frac{2}{4^2} & \text{CASE II } x(i)=j; x(i)=i; x(k)=k \quad |i-j| = 2 \\ \frac{1}{4} & \text{CASE III } x(k)=k \text{ for all "k"} \end{cases}$$

$n(n-1) - 2(n-1) = n^2 - n - 2n + 1 = n^2 - 3n + 1 =$

$\rightarrow = (n-2)(n-1)$ INSTANCES OF CASE I =

$(n-1)$ INSTANCES OF CASE II
 "n" DUKLIKA
 (1) INSTANCE OF CASE III (ALL NUMBERS ARE AT ITS OWN RIGHT PLACE I.E. NO MISORDER FILES)
 (VADIS 1 = STAVI TA 1, VADIS 2)
 - OPRIPATA III "e" VOZMOZNA ZOITO KAKOVA MORE DA ZA VRATIS VO NEZEMOTO INKIPAZO MESTO
 - NE SE DOVM PRIMERITE NA P.26 (THE SE AND IZVLEKIS 2 KATI I M GI ZAKREVIS MESTOVA)

1 2 3		
2 1 3*	2 1 3*	3 1 2
2 3 1	1 2 2*	1 3 2*

1 2 3 4			
* 2 1 3 4	* 2 1 3 4	3 1 2 4	4 1 2 3
2 3 1 4	1 3 2 4*	1 3 2 4*	1 4 2 3
2 3 4 1	1 3 4 2	1 2 4 3*	1 2 4 3*

(*) ADJESCENT CASES

$(n-2)(n-1)$

$$H(x) = \sum_{i=1}^n \frac{1}{4^2} \log 4^2 + \sum_{k=1}^{(n-1)} \frac{2}{4^2} \log \frac{4^2}{2} + \frac{1}{4} \log 4$$

$$= \frac{(n-2)(n-1)}{4^2} \log 4^2 + \frac{2(n-1)}{4^2} \log \frac{4^2}{2}$$

~~$(n-1) + (n-2)(n-1) + 1 = n - 1 + 1 + n^2 - 3n + 1 = n^2 - 2n + 1 = (n-1)^2$~~

~~$P(x=x_1) = \frac{(n-2)(n-1)}{(n-1)^2} = \frac{n-2}{n-1}$~~

~~$P(x=x_2) = \frac{n-1}{(n-1)^2} = \frac{1}{n-1}$~~

$$P(x=x_1) = \frac{n-2}{n-1}$$

$$P(x=x_2) = \frac{1}{n-1}$$

$$P(x=x_3) = \frac{1}{(n-1)^2}$$

VELOZATNOST NA SE SVUČI MISSORTING CASE I

VELOZATNOST NA SE SVUČI MISSORTING CASE II

VELOZATNOST NA SE SVUČI MISSORTING CASE III
T.R. SOLUTIONS

• SAMO ZA EDEN MISSORTING OD DAVEN VTT

$$y_1 = \frac{P(x=x_1)}{(n-2)(n-1)} = \frac{1}{(n-1)^2}$$

$$y_2 = \frac{1(x=x_2)}{(n-1)^2}$$

VKURNO NASTAVNI:

VADIS 1 STAVAS NA 1, VADIS 2 STAVAS NA 2, ...

$$n(n-1) + n = \frac{n^2 - n + n}{1} = n^2$$

$$P(x \in X_1) = \frac{(n-2)(n-1)}{n^2}$$

$$P(x \in X_2) = \frac{2(n-1)}{n^2}$$

$$P(x \in X_3) = \frac{n}{n^2}$$

$$(n-2)(n-1) + 2(n-1)$$

• SAMO ZA EDEN MISSORTING OD DAVEN TIP

$$y_1 = P(x=x_1) = \frac{P(x \in X_1)}{(n-2)(n-1)} = \frac{1}{n^2}$$

$$y_2 = P(x=x_2) = \frac{P(x \in X_2)}{(n-1)}$$

$$y_3 = \frac{P(x \in X_3)}{1} = \frac{1}{n}$$

DUPLIKATIFE ZA EZOLEKUVAT VELOZAT-OKTA

$$y_1 = \frac{1}{n^2} ; y_2 = \frac{2}{n^2} ; y_3 = \frac{1}{n}$$

$$H(x) = (n-1)(n-2) \cdot \frac{1}{n^2} \ln n^2 + \underbrace{(n-1) \cdot \frac{2}{n^2}}_{\text{BEZ DUKLIFATI}} \cdot \ln \frac{n^2}{2} + \frac{1}{n} \ln n$$

$$= \frac{2(n-1)(n-2)}{n^2} \ln n + \frac{4(n-1)}{n^2} \ln n - \frac{2(n-1)}{n^2} \ln 2 + \frac{1}{n} \ln n =$$

$$= \ln(n) \left[\frac{2(n-1)(n-2) + 4(n-1) + n}{n^2} \right] - \frac{n-1}{n^2} \ln 2$$

$$2(4-1)(4-2) + 4(4-1) + 4 = 2(4^2 - 2 \cdot 4 - 4 + 2) + 4 \cdot 4 - 4 + 4 =$$

$$= 2 \cdot 4^2 - 6 \cdot 4 + 4 + 5 \cdot 4 - 4 = 2 \cdot 4^2 - 4 = 4(2 \cdot 4 - 1)$$

$$H(x) = \frac{4(2 \cdot 4 - 1)}{4^2} \log_2(4) - \frac{4-1}{4^2} \log_2 4$$

$$H(x) = \frac{2 \cdot 4 - 1}{4} \log_2 4 - \frac{4-1}{4^2} \log_2 4$$

МОДЕЛЬ ЗАПИСИ ДАННЫХ НА СЕ ВИКА: ENTROPY
OF TAKING OUT AND RETURNING BACK THE FILE
FROM ARCHIVE.

2.48 SEQUENCE LENGTH; HOW MUCH INFORMATION
DOES THE LENGTH OF THE SEQUENCE GIVE
ABOUT THE CONTENT OF THE SEQUENCE? SUPPOSE

WE CONSIDER A BERNOULLI $(\frac{1}{2})$ PROCESS $\{X_i\}$.
STOP THE PROCESS WHEN THE FIRST 1 APPEARS.
LET "N" DESIGNATE THIS STOPPING TIME. THUS,
 X^N IS AN ELEMENT OF THE SET OF ALL FINITE-
LENGTH BINARY SEQUENCES $\{0,1\}^* = \{0, 1, 00, 01, 10, 11,$
 $000, \dots\}$.

(a) FIND $I(N; X^N)$. (b) FIND $H(X^N|N)$ (c) FIND $H(X^N)$

LET'S NOW CONSIDER A DIFFERENT STOPPING TIME. FOR
THIS PART, AGAIN ASSUME THAT $X_i \sim \text{BERNOULLI}(\frac{1}{2})$
BUT STOP AT TIME $N=6$ WITH PROBABILITY $\frac{1}{3}$
AND STOP AT TIME $N=12$ WITH PROBABILITY $\frac{2}{3}$.
LET THIS STOPPING TIME BE INDEPENDENT OF THE
SEQUENCE X_1, X_2, \dots, X_{12} .

(d) FIND $I(N; X^N)$. (e) FIND $H(X^N|N)$. (f) FIND $H(X^N)$.

$$(a) P(N) = \frac{1}{2} \cdot \frac{1}{2} \cdot \dots \cdot \frac{1}{2} \left(1 - \frac{1}{2}\right) = \left(\frac{1}{2}\right)^N$$

$$H(N) = \sum_{N=1}^{\infty} \left(\frac{1}{2}\right)^N \log_2 2^N = \log_2 \sum_{N=1}^{\infty} N \left(\frac{1}{2}\right)^N = \log_2 \frac{\frac{1}{2}}{\left(1 - \frac{1}{2}\right)^2}$$

$$H(N) = \log_2 \frac{\frac{1}{2}}{\frac{1}{4}} = (\log_2 2) \cdot 2 = \underline{\underline{2}}$$

$$I(N; X^N) = H(N) - H(N|X^N) = H(N) = 2 \cdot \log_2 2 = 2$$

⊙ NO GO ZONES X^N GO
ZONES "N"

(6) $H(X^N|N) = ?$ $I(N; X^N) = H(N) - H(N|X^N) = H(N)$

$H(X^N|N) = H(X_1|N) + H(X_2|N, X_1) + \dots + H(X_N|N, X_1, \dots, X_{N-1})$

$X^N \in \{x_1, x_2, \dots, x_N\}$

$H(X^N|N) = \sum_{i=1}^N P(i) \cdot H(X^N|N=i)$

$I(N; X^N) = H(X^N) - H(X^N|N) = N - H(X^N|N)$

$H(X^N) = N \cdot H(X_i) = N \cdot [-p \log p - (1-p) \log (1-p)] = N \cdot \dots$

$H(X^N|N) = N - I(N; X^N) = N - 2 \log 2$

$X^N \in \{0, 1, 00, 01, 10, 11, 000, 001, 010, 011, 100, 101, 110, 111\}$

$N=1 \Rightarrow X^N \in \{1, 10, 11, 100, 101, 110, 111\}$

$X^N \in \{0, 1\}$

$N=2 \quad X^N \in \{0, 1, 00, 01, 10, 11\}$

AVO PAVATA "1" E TOXE N=2
A X^N E OGIANICHO NA POZINA
2. NETA NEBEVOST.

$H(X^N|N=2) = \emptyset$

$N=3 \quad X^N \in \{0, 1, 00, 01, 10, 11, 000, 001, 010, 011, 100, 101, 110, 111\}$

$H(X^N|N=2) = P(00) \log \frac{1}{P(00)} + P(01) \log \frac{1}{P(01)} + P(10) \log \frac{1}{P(10)} + P(11) \log \frac{1}{P(11)} = \left(\frac{1}{4} \cdot \log 4\right) \cdot 4 = \frac{4}{2} \cdot 4 = 2$

$H(X^N|N=4) = \left(\frac{1}{2^4} \log 2^4\right) \cdot 2^4 = 4 \log 2 = 4$

$H(X^N|N) = \sum_{n=1}^N P(n) H(X^N|N=n) = \sum_{n=1}^N \frac{1}{2^n} n$

$S = \sum_{n=1}^N n x^n \quad \int x^n dx = \frac{x^{n+1}}{n+1} \quad \int x^{n-1} dx = \frac{x^n}{n}$

$\frac{S}{x} = \sum_{n=1}^N n x^{n-1} \quad \int \Rightarrow \int \frac{S}{x} dx = \sum_{n=1}^N n \frac{x^n}{x}$

$$\int \frac{S dt}{t} = \sum_{k=1}^N t^k$$

$$S = x + x^2 + \dots + x^N$$

$$\frac{S}{x} = 1 + x + \dots + x^{N-1}$$

$$S - \frac{S}{x} = x^N - 1$$

$$S \left(1 - \frac{1}{x}\right) = x^N - 1$$

$$S = \frac{x^N - 1}{\frac{x-1}{x}} = \frac{x^{N+1} - x}{x-1}$$

$$\int \frac{S dt}{t} = \frac{x^{N+1} - x}{x-1} \quad \frac{S}{x} = \left(\frac{x^{N+1} - x}{x-1} \right)' = \frac{[(N+1)x^N - 1](x-1)}{(x-1)^2} - \frac{x^{N+1} + x}{(x-1)^2}$$

$$S = x \cdot \frac{(N+1)x^N(x-1) - x^{N+1} - x}{(x-1)^2} = x \frac{(N+1)x^{N+1} - (N+1)x^N - x^{N+1} + x}{(x-1)^2}$$

$$S = x \frac{Nx^{N+1} + x^{N+1} - Nx^N - x^N + 1 - x^{N+1}}{(x-1)^2} = x \frac{Nx^{N+1} - (N+1)x^N + 1}{(x-1)^2}$$

$$S = x \cdot \frac{x^N(Nx - N - 1) + 1}{(x-1)^2}$$

$$x = \frac{1}{2}$$

$$S = \frac{1}{2} \frac{\frac{1}{2^N} \left(\frac{N}{2} - N - 1 \right) + 1}{\frac{1}{4}} = 2 \left[\frac{1}{2^N} \left(-\frac{N}{2} - 1 \right) + 1 \right]$$

$$S = 2 - \frac{1}{2^{N+1}} \left(1 + \frac{N}{2} \right)$$

$$S = 2 - \frac{N+2}{2^N}$$

$$S = 2 - N \cdot 2^{-N} - 2^{1-N}$$

$$H(x^N | N) = 2 - \frac{N+2}{2^N}$$

$$N \rightarrow \infty \quad H(x^N | N) = 2$$

$$(c) \quad I(N; x^N) = H(x^N) - H(x^N | N) = 2 \Rightarrow$$

$$H(x^N) = 2 + H(x^N | N) = 4 - \frac{N+2}{2^N}$$

$$P(N=1) \quad H(x^N | 1) = 2 - \frac{3}{2} = 0.5$$

$$H(x^1 | 1) = P(N=1) \cdot H(x^N | N=1) = \frac{1}{2} \cdot \left[\frac{1}{2} \log 2 + \frac{1}{2} \log 2 \right] = \frac{1}{2}$$

$$H(x^N) = 4 - \frac{N+2}{2^N}$$

$$(d) P(N=6) = \frac{1}{3} \quad P(N=12) = \frac{2}{3}$$

$$H(N) = P(N=6) \log \frac{1}{P(N=6)} + P(N=12) \log \frac{1}{P(N=12)} =$$

$$= \frac{1}{3} \log 3 + \frac{2}{3} \log \frac{3}{2} = \frac{1}{3} \log 3 + \frac{2}{3} \log 3 - \frac{2}{3} \log 2$$

$$= \log 3 - \frac{2}{3} \log 2 \quad H(N) = \log 3 - \frac{2}{3} \log 2 = \log 3 - \frac{2}{3}$$

$$I(N; X^N) = H(N) - H(N|X^N) = H(N) = \log 3 - \frac{2}{3} = 0,92$$

$$(e) H(X^N|N) = ? \quad I(N; X^N) = H(X^N) - H(X^N|N)$$

$$H(X^N|N=6) = -P(000000) \log P(000000) - P(000001) \log P(000001) -$$

$$- P(000010) \log P(000010) - P(000011) \log P(000011) - \dots - P(111111) \log P(111111)$$

$$H(X^N|N=6) = \left(\frac{1}{2^6} \log 2^6 \right) \cdot 2^6 = \log 2^6 = 6$$

$$H(X^N|N=12) = \left(\frac{1}{2^{12}} \log 2^{12} \right) \cdot 2^{12} = 12$$

$$H(X^N|N) = P(N=6) \cdot H(X^N|N=6) + P(N=12) \cdot H(X^N|N=12) =$$

$$= \frac{1}{3} \cdot 6 + \frac{2}{3} \cdot 12 = 2 + 8 = \underline{\underline{10}}$$

$$H(X^N|N) = 10$$

$$(d) H(X^N) = ? \quad H(X^N) = I(N; X^N) + H(X^N|N) \Rightarrow$$

$$H(X^N) = \log 3 - \frac{2}{3} + 10 = \log 3 + \frac{30-2}{3} = \log 3 + \frac{28}{3}$$

$$\boxed{H(X^N) = 10,92}$$

CHAPTER 3 ASYMPTOTIC EQUIVALENCE

$$\frac{1}{n} \log \frac{1}{P(x_1, x_2, \dots, x_n)} \rightarrow H \quad x_1, x_2, \dots, x_n$$

$$\left(P(x_1, x_2, \dots, x_n) \right)^{\frac{1}{n}} \rightarrow 2^{-H} \quad \left[P(x_1, x_2, \dots, x_n) \right]^{\frac{1}{n}} = 2^{-H} \quad ()^n$$

$$\boxed{P(x_1, x_2, \dots, x_n) \rightarrow 2^{-nH}}$$

SE NAKAZANI NA PROBLEMI 2.48

$$H(X^N) = 4 - \frac{N+2}{2^N}$$

$$P(0,0000000000000001) = \frac{1}{2^{12}} \quad 2^{-4H} = 2^{-4N + \frac{N+2}{2^N}}$$

• TYPICAL SET WHERE SAMPLE ENTROPY IS CLOSE TO THE TRUE ENTROPY.

• EXAMPLE. LET RANDOM VARIABLE $X \in \{0,1\}$ HAVE PROBABILITY MASS FUNCTION $P(X=1) = p$ $P(X=0) = q$
 IF X_1, X_2, \dots, X_n ARE I.I.D. PROBABILITY OF X_1, X_2, \dots, X_n IS $\prod_{i=1}^n P(X_i)$

e.g: $(1,0,1,1,0,1) \rightarrow P(X) = p^4 q^2 = 2^{-4H} = 2^{-4 \cdot 0.92} = 2^{-3.68} \approx 2^{-4 + 0.32}$

NOT ALL SEQUENCES OF LENGTH n HAS SAME PROBABILITY.

$$P\{X_1, X_2, \dots, X_n : P(X_1, X_2, \dots, X_n) = 2^{-n(H \pm \epsilon)}\} \approx 1$$

$$P(X_1, X_2, \dots, X_n) = p^{\sum x_i} q^{n - \sum x_i}$$

NUMBER OF 1'S IN SEQUENCE IS CLOSE TO $n \cdot p$ AND ALL SUCH SEQUENCES HAVE RUGHLY THE SAME PROBABILITY $2^{-nH(p)}$

ALL SUCH SEQUENCES HAVE RUGHLY THE SAME PROBABILITY $2^{-nH(p)}$

$$X_1 X_2 X_3 X_4 X_5$$

$$p = 0.6 \quad q = 0.4$$

$$n = 5$$

$$4 \cdot p = 0.6 \cdot 5 = 3$$

$$4 \cdot q = 5 \cdot 0.4 = 2$$

$$p^3 q^2 = (0.6)^3 \cdot (0.4)^2 = 0.03456$$

$$2^{-5 \cdot H(p)} = 0.03575$$

DEFINITION: (CONVERGENCE OF RANDOM VARIABLES) GIVEN A SEQUENCE OF RANDOM VARIABLES X_1, X_2, \dots , WE SAY THAT SEQUENCE X_1, X_2, \dots CONVERGES TO RANDOM VARIABLE X :

1. In probability IF EVERY $\epsilon > 0$, $P\{|X_n - X| > \epsilon\} \rightarrow 0$
2. In mean square IF $E[(X_n - X)^2] \rightarrow 0$
3. With probability 1 (also called almost surely) IF $P\{\lim_{n \rightarrow \infty} X_n = X\} = 1$

3.1 ASYMPTOTIC EQUIPARTITION PROPERTY THEORY

THEOREM 3.1.1 (AEP) If x_1, x_2, \dots are i.i.d $\sim p(x)$

THEN:

$$-\frac{1}{n} \log p(x_1, x_2, \dots, x_n) \xrightarrow{i.p.} H(X)$$

PROBABILITY:

OBNEZNO GLAVNO BI IZVEDUJANJA VO MAREZNI
MAGLEJNINO SPOJITNO

$$-\frac{1}{n} \log p(x_1, x_2, \dots, x_n) = -\frac{1}{n} [\log p(x_1) + \log p(x_2) + \dots + \log p(x_n)] =$$

$$-\frac{1}{n} \sum_{i=1}^n \log p(x_i) \rightarrow -E[\log p(x_i)] \text{ IN PROBABILITY}$$

SAMPLE AVERAGE

EXPECTED VALUE OF $\log p(x_i)$

LAW OF LARGE NUMBERS

$$\Rightarrow \log [p(x_1, x_2, \dots, x_n)] \stackrel{-1}{\sim} H(X) \quad \log p(x_1, x_2, \dots, x_n) = -nH(X)$$

$$p(x_1, x_2, \dots, x_n) = 2^{-n \cdot H(X)} \quad (*)$$

EDROKOVOSTA VAEI ZA SEKVENCI SO $n \rightarrow \infty$
BLOD NA EDKICI

DEFINITION: THE TYPICAL SET $A_\epsilon^{(n)}$ WITH RESPECT TO $p(x)$ IS THE SET OF SEQUENCES $(x_1, x_2, \dots, x_n) \in X^n$ WITH THE PROPERTY

$$2^{-n(H(X)+\epsilon)} \leq p(x_1, x_2, \dots, x_n) \leq 2^{-n(H(X)-\epsilon)} \quad (\#)$$

ALTERNATIVA FORMA: $|\frac{1}{n} \log p(x_1, x_2, \dots, x_n) - H(X)| \leq \epsilon$

TYPICAL SET IS SET WHERE SAMPLE ENTROPY IS CLOSE TO REAL ENTROPY.

THEOREM 3.1.2 AS A CONSEQUENCE OF AEP WE CAN SHOW THAT THE SET $A_\epsilon^{(n)}$ HAS THE FOLLOWING PROPERTIES:

- (1) If $(x_1, x_2, \dots, x_n) \in A_\epsilon^{(n)}$, THEN $H(X) - \epsilon \leq -\frac{1}{n} \log p(x_1, x_2, \dots, x_n) \leq H(X) + \epsilon$
- (2) $\Pr\{A_\epsilon^{(n)}\} \geq 1 - \epsilon$ FOR n SUFFICIENTLY LARGE
- (3) $|A_\epsilon^{(n)}| \leq 2^{n(H(X)+\epsilon)}$ WHERE $|A|$ DENOTES THE NUMBER OF ELEMENTS IN THE SET A .
- (4) $|A_\epsilon^{(n)}| \geq (1 - \epsilon) 2^{n(H(X)-\epsilon)}$

THUS THE TYPICAL SET HAS PROBABILITY NEARLY 1, AND ALL ELEMENTS OF TYPICAL SET ARE NEARLY EQUIPROBABLE, AND

(34) NUMBER OF ELEMENT IS NEARLY 2^{nH}

PROOF: (1) $- \ln(H(x) + \epsilon) \leq \ln \gamma(x_1, x_2, \dots, x_n) \leq - \ln(H(x) - \epsilon)$

$\ln(H(x) + \epsilon) \geq - \ln \gamma(x_1, x_2, \dots, x_n) \geq \ln(H(x) - \epsilon) \Rightarrow$

$$H(x) - \epsilon \leq - \frac{1}{n} \ln \gamma(x_1, x_2, \dots, x_n) \leq H(x) + \epsilon$$

(2) $\Rightarrow \gamma(x_1, x_2, \dots, x_n) = 2^{-n H(x)}$

$\lim_{n \rightarrow \infty} \gamma(x_1, x_2, \dots, x_n) = \lim_{n \rightarrow \infty} 2^{-n H(x)} = \frac{1}{2^{n H(x)}} \rightarrow 0$

FOR ANY $\delta > 0$ THERE EXIST n_0 SUCH THAT $n \geq n_0$ WE HAVE:

$P\{ \left| - \frac{1}{n} \ln \gamma(x_1, x_2, \dots, x_n) - H(x) \right| < \epsilon \} > 1 - \delta$

000, 001, 010, 011, 100, 101, 110, 111
 a b c 2 3

$$\begin{cases} p = 0.7 \\ q = 0.3 \end{cases}$$

$(0.7)^2 (0.3) = 0.147$

$(0.3)^2 (0.7) = 0.063$

(3) $1 = \sum_{x \in X^n} \gamma(x) \geq \sum_{x \in A_\epsilon^{(n)}} \gamma(x) \geq \sum_{x \in A_\epsilon^{(n)}} 2^{-n(H(x) + \epsilon)} =$

$= |A_\epsilon^{(n)}| \cdot 2^{-n(H(x) + \epsilon)}$

$\Rightarrow 2^{n(H(x) + \epsilon)} \geq |A_\epsilon^{(n)}|$
 i.e. $|A_\epsilon^{(n)}| \leq 2^{n(H(x) + \epsilon)}$

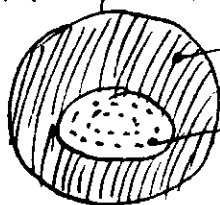
\sum ELEMENTS OF $A_\epsilon^{(n)}$

(4) $(1 - \epsilon) \leq P\{A_\epsilon^{(n)}\} \leq \sum_{x \in A_\epsilon^{(n)}} 2^{-n(H(x) - \epsilon)} = |A_\epsilon^{(n)}| \cdot 2^{-n(H(x) - \epsilon)}$

$|A_\epsilon^{(n)}| \geq (1 - \epsilon) 2^{n(H(x) - \epsilon)}$ DOKAZANO!!!

3.2 CONSEQUENCES OF THE AEP: DATA COMPRESSION

In $A_\epsilon^{(n)}$ THERE ARE $\leq 2^{n(H(x) + \epsilon)}$ SEQUENCES WHICH CAN BE INDEXED BY $n(H(x) + \epsilon) + 1$ BITS. PREFIX \emptyset FOR EACH SEQUENCE \Rightarrow LENGTH OF INDEXING NUMBER $\leq n(H(x) + \epsilon) + 2$ BITS



NON-TYPICAL SET: $n \log |A| + 2$ bits
 TYPICAL SET: $n(H(x) + \epsilon) + 2$ bits

x^n DENOTES A SEQUENCE x_1, x_2, \dots, x_n
 $L(x^n)$ - LENGTH OF THE CODEWORD CORRESPONDING TO x^n
 - IF n IS SUFFICIENTLY LARGE SO THAT $P\{A_E^{(n)}\} > 1 - \epsilon$
 THE EXPECTED LENGTH OF CODEWORD IS:

$$E[L(x^n)] = \sum_{x^n} p(x^n) \cdot L(x^n)$$

$$\begin{aligned}
 E[L(x^n)] &= \sum_{x^n} p(x^n) L(x^n) = \sum_{x^n \in A_E^{(n)}} p(x^n) L(x^n) + \sum_{x^n \in A_E^{(n)c}} p(x^n) L(x^n) \\
 &\leq \sum_{x^n \in A_E^{(n)}} p(x^n) (n(H+\epsilon) + 2) + \sum_{x^n \in A_E^{(n)c}} p(x^n) (n \log 2 + 2) \\
 &= P\{A_E^{(n)}\} (n(H+\epsilon) + 2) + P\{A_E^{(n)c}\} (n \log 2 + 2) \\
 &= n(H+\epsilon) + \epsilon n \log 2 + 2 = n(H+\epsilon')
 \end{aligned}$$

$\epsilon' = \epsilon \log 2 + \frac{2}{n} + \epsilon$

THEOREM 3.2.1 LET x^n BE i.i.d $\sim p(x)$. LET $\epsilon > 0$
 THEN THERE EXISTS A CODE THAT MAPS x^n OF
 LENGTH n INTO BINARY STRINGS SUCH THAT
 THE MAPPING IS ONE-TO-ONE (AND THEREFORE INVE-
 RTIBLE) AND:

$$E\left[\frac{1}{n} L(x^n)\right] \leq H(x) + \epsilon$$

HENCE WE CAN REPRESENT SEQUENCE x^n USING
 $n(H(x) + \epsilon)$ BITS ON AVERAGE.

3.3 TIGHT PROBABILITY SETS AND TYPICAL SET

FROM THE DEFINITION OF $A_E^{(n)}$ IT IS CLEAR THAT
 $A_E^{(n)}$ IS FAIRLY SMALL SET THAT CONTAINS MOST OF
 THE PROBABILITY. HOWEVER IT IS THE SMALLEST
 SET? WE WILL PROVE THAT THE TYPICAL SET HAS
 ESSENTIALLY THE SAME NUMBER OF ELEMENTS
 AS THE SMALLEST SET, TO FIRST ORDER IN ELEMENT.

DEFINITION: FOR EACH $n = 1, 2, \dots$ LET $B_n^{(n)} \subset X^n$
 BE THE SMALLEST SET WITH:

$$P\{B_n^{(n)}\} \geq 1 - \delta$$

WE ARGUE THAT $B_n^{(n)}$ MUST HAVE SIGNIFICANT INTER-
 SECTION WITH $A_E^{(n)}$ AND THEREFORE MUST HAVE ABOUT

AS MANY ELEMENTS. IN PROBLEM 3.11 WE OBTAIN THE PROOF OF FOLLOWING THEOREM:

THEOREM 3.3.1. LET x_1, x_2, \dots, x_n BE I.I.D.N $\gamma(t)$, FOR $\delta < \frac{1}{2}$ AND $\delta' > 0$, IF $\Pr\{B_{\delta}^{(n)}\} > 1 - \delta$ THEN

$$\frac{1}{n} \log |B_{\delta}^{(n)}| > H - \delta' \quad \text{FOR } n \text{ SUFFICIENTLY LARGE.}$$

THUS, $B_{\delta}^{(n)}$ MUST HAVE AT LEAST $2^{n(H - \delta')}$ ELEMENTS TO FIRST ORDER IN THE EXPONENT. $A_{\epsilon}^{(n)}$ HAS $2^{n(H(\epsilon) \pm \epsilon)}$ ELEMENTS.

DEFINITION: THE NOTATION $a_n \doteq b_n$ MEANS

$$\lim_{n \rightarrow \infty} \left[\frac{1}{n} \log \frac{a_n}{b_n} \right] = 0$$

THIS IMPLIES THAT a_n & b_n ARE EQUIV TO THE FIRST ORDER IN THE EXPONENT.

• THEREFORE, IF $\delta_n \rightarrow 0$ & $\epsilon_n \rightarrow 0$ THEN

MMV $|B_{\delta_n}^{(n)}| \doteq |A_{\epsilon_n}^{(n)}| \doteq 2^{nH}$ [071309700]

3.1 MARKOV INEQUALITY & CHEBYSHEV INEQUALITY.

(a) (Markov's inequality) FOR ANY NONNEGATIVE RANDOM VARIABLE X AND ANY $t > 0$ SHOW THAT

$$\Pr\{X \geq t\} \leq \frac{E[X]}{t}$$

EXIST RANDOM VARIABLE THAT ACHIEVES THIS INEQUALITY WITH EQUALITY

(b) (Chebyshev's inequality) LET γ BE A RANDOM VARIABLE WITH MEAN μ AND VARIANCE σ^2 . BY LETTING $X = (\gamma - \mu)^2$ SHOW THAT FOR ANY $\epsilon > 0$

$$(a) \Pr\{X \geq t\} = 1 - \Pr\{X \leq t\} \quad \Pr\{|\gamma - \mu| > \epsilon\} \leq \frac{\sigma^2}{\epsilon^2}$$

$$E[X] = \sum_{i=1}^n p(x_i) \cdot x_i$$

$$X \in \{1, 2, \dots, n\}$$

$$p(x) = \{p_1, p_2, \dots, p_n\}$$

$$Pr(X > t) = 1 - Pr\{X \leq t\} = 1 - \sum_{i=1}^t p_i$$

$$\frac{E[X]}{t} = \frac{1}{t} \sum_{i=1}^n x_i p_i \geq 0 \quad \boxed{\frac{E[X]}{t} \geq Pr(X > t)}$$

$$\frac{1}{t} \sum_{i=1}^n x_i p_i \geq 1 - \sum_{i=1}^t x_i p_i = \sum_{i=t+1}^n x_i p_i$$

$$\sum_{i=1}^n x_i p_i \geq t \sum_{i=t+1}^n x_i p_i \quad \sum_{i=t+1}^n x_i p_i - \sum_{i=t+1}^n t x_i p_i + \sum_{i=1}^t x_i p_i$$

$$\sum_{i=t+1}^n (1-t) x_i p_i + \sum_{i=1}^t x_i p_i \geq 0 \quad ?$$

$$\sum_{i=1}^n x_i p_i \geq t + \sum_{i=1}^t x_i p_i \quad \sum_{i=1}^t x_i p_i + \sum_{i=t+1}^n x_i p_i + t \sum_{i=1}^t x_i p_i \geq t$$

$$(1+t) \sum_{i=1}^t x_i p_i + \sum_{i=t+1}^n x_i p_i \geq t$$

$$\frac{1}{t} \sum_{i=1}^n x_i p_i = \frac{1}{t} \left[\sum_{i=1}^t x_i p_i + \sum_{i=t+1}^n x_i p_i \right]$$

$X \in [1, 2, \dots, t, \dots, n]$

$P(X) = \{p_1, p_2, \dots, p_t, \dots, p_n\}$

$$E[X] = \sum_{i=1}^n i \cdot p_i = p_1 + 2p_2 + 3p_3 + \dots + t p_t + \dots + n p_n$$

$$\frac{E[X]}{t} = \frac{p_1}{t} + \frac{2p_2}{t} + \dots + p_t + \dots + \frac{n}{t} p_n \geq$$

$$p_1 + p_2 + \dots + p_t + \dots + \frac{n}{t} p_n$$

$$1 - \sum_{i=1}^t p_i = 1 - p_1 - p_2 - \dots - p_t$$

$$Pr(X > t) = \sum_{i=t+1}^n p_i = p_{t+1} + p_{t+2} + \dots + p_n$$

$$\frac{E[X]}{t} = \frac{1}{t} + \frac{2p_2}{t} + \dots + p_t + \frac{(t+1)p_{t+1}}{t} + \dots + \frac{n}{t} p_n$$

$$\leq \frac{p_1}{t} + \frac{2p_2}{t} + \dots + p_t + \underbrace{(t+1)p_{t+1} + \dots + (n-1)p_{n-1}}_{Pr(X > t)}$$

$$x \in [t, 2t, 3t, \dots, nt]$$

$$\begin{aligned} S &= 1+2t+\dots+nt \\ S &= n+2t+\dots+nt \\ 2S &= (n+1)n \\ S &= \frac{(n+1)n}{2} \end{aligned}$$

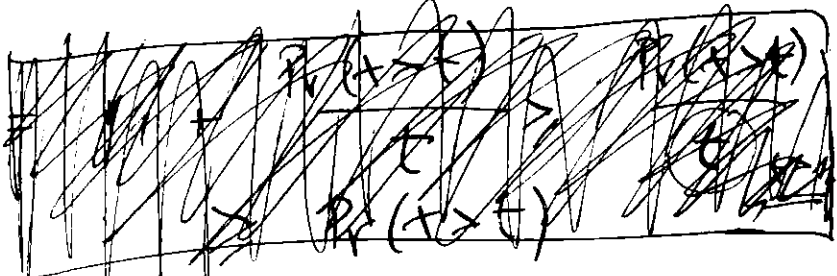
$$\frac{E[X]}{t} = \sum_{i=1}^n i p_i$$

$$Pr(X > t) = \sum_{i=2}^n i p_i$$

$$\frac{E[X]}{t} = \sum_{i=1}^n i p_i$$

$$Pr(X > t) = t \sum_{i=2}^n p_i$$

$$\frac{E[X]}{t} = p_1 + \sum_{i=2}^n i p_i$$



$$\frac{E[X]}{t}$$

$$= \frac{1}{t} + \frac{2p_2}{t} + \dots + p_t + \frac{(t+1)p_{t+1}}{t} + \dots + \frac{n-1}{t} p_{n-1}$$

$$\geq \underbrace{p_1 + p_2 + \dots + p_t}_{1 - Pr(X > t)} + \frac{(t+1)p_{t+1}}{t} + \dots + \frac{n-1}{t} p_{n-1}$$

$$1 - Pr(X > t)$$

$$E[X] = \frac{1}{n} \sum_{i=1}^n i = \frac{1}{n} \frac{n \cdot (n+1)}{2} = \frac{n+1}{2}$$

$$\frac{E[X]}{t} = \frac{1}{t} \cdot \frac{n+1}{2}$$

$$x \in [1, 2, \dots, t, \dots, n] \\ Pr(X > t) = \sum_{i=t+1}^n p_i$$

$$Pr(X > t) = p_{t+1} + p_{t+2} + \dots + p_n$$

$$\frac{E[X]}{t} = \frac{1}{t} + \frac{2p_2}{t} + \dots + \frac{(t+1)p_{t+1}}{t} + \dots + \frac{n}{t} p_n \leq \frac{p_1}{t} + \frac{2p_2}{t} + \dots + p_{t+1} + \dots + p_n$$

• РИМКА ЗА УНИФОРМА ПАРЦЕЛЕНА:

$$x \in \{1, 2, 3, \dots, 4\}$$

$$p(x) = \left\{ \frac{1}{4}, \frac{1}{4}, \dots, \frac{1}{4} \right\}$$

$$Pr\{x > t\} = (4-t) \cdot \frac{1}{4}$$

$$E[X] = \frac{1}{4} \sum_{i=1}^4 i = \frac{1}{4} \frac{4(4+1)}{2}$$

$$E[X] = \frac{4+1}{2}$$

$$\frac{E[X]}{t} = \frac{4+1}{2t} \quad (\text{?})$$

$$\frac{4+1}{2t} \geq \frac{(4-t)}{4}$$

$$4^2 + 4 \geq 24t - 2t^2$$

$$4^2 + 4 + 2t^2 - 24t \geq 0$$

$$f(t) = 2t^2 - 24t + 4(4+1) \quad \text{JPSDDP}$$

VO OVO
SUCAS NE
VAZI
NEPRAVNO.

• РИМКА СО ГЕОМЕТРИСКА ПАРЦЕЛЕНА

$$x \in \{1, 2, 3, \dots, 4\}$$

$$p(x) = \left\{ \frac{1}{2}, \frac{1}{2^2}, \dots, \frac{1}{2^{4+1}} \right\}$$

$$S = 2^1 + 2^2 + \dots + 2^4$$

$$2S = 2^2 + 2^3 + \dots + 2^{4+1}$$

$$S(1-2) = 2 - 2^{4+1}$$

$$S = \frac{2(1 - \frac{1}{2^4})}{\frac{1}{2}}$$

$$S = \left(1 - \frac{1}{2^4}\right)$$

$$S = \left(1 - \frac{1}{2^{4-1}}\right) + \frac{1}{2^{4-1}} = 1$$

$$Pr(x > t) = 1 - \left(1 - \frac{1}{2^t}\right) = \frac{1}{2^t}$$

МАКЕТ!!!

$$E[X] = \sum_{i=1}^{4-1} i \cdot p(i) = \sum_{i=1}^{4-1} i \cdot \frac{1}{2^i} + \frac{4}{2^{4-1}} = \frac{4}{2^{4-1}} - \frac{1}{2^{4-1}} + 2 + \frac{4}{2^{4-1}}$$

$$E[X] = 2 - \frac{1}{2^{4-1}}$$

$$\frac{4}{2^{4-1}} - \frac{1}{2^{4-1}} + 2 + \frac{4}{2^{4-1}}$$

~~$E[X] = \sum_{i=1}^n (i-1) \cdot \frac{1}{2^{i-1}}$~~ ~~$E[X] = \sum_{i=1}^n (i-1) \cdot \frac{1}{2^{i-1}}$~~ ~~$E[X] = \sum_{i=1}^n (i-1) \cdot \frac{1}{2^{i-1}}$~~

$$\frac{E[X]}{t} = \frac{2}{t} - \frac{1}{t \cdot 2^{n-1}}$$

$$Pr(X > t) = \frac{1}{2} t$$

$$\frac{2}{t} - \frac{1}{t \cdot 2^{n-1}} \geq \frac{1}{2} t$$

$$\frac{2 \cdot 2^t}{t} - \frac{2^t}{t \cdot 2^{n-1}} \geq 1$$

$$2 \cdot 2^t - 2^{t-n+1} \geq t$$

$$2^{t-1} - 2^{t-n+1} - t \geq 0$$

$$\frac{2^t}{t} \left[2 - \frac{1}{2^{n-1}} \right] \geq 1$$

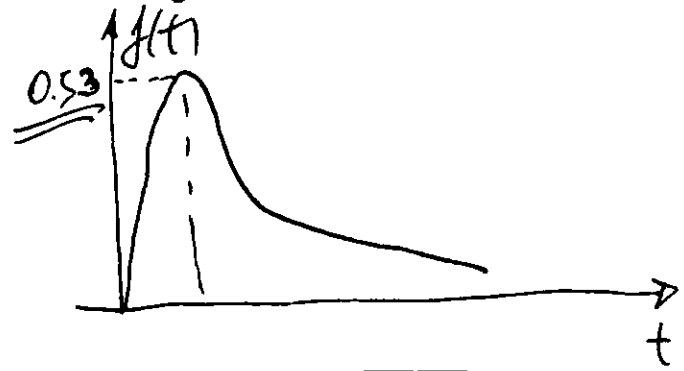
$$\left(2 - \frac{1}{2^{n-1}} \right) \geq \frac{t}{2^t}$$

$$\left(2 - \frac{1}{2^{n-1}} \right) \geq \underbrace{t \cdot 2^{-t}}_{f(t)}$$

$$f(0) = 0$$

$$f(1) = \frac{1}{2}$$

$$f(2) = \frac{2}{2^2} = \frac{1}{4}$$



$$\left. \begin{aligned} g(1) &= g(1) = 1 \\ g(2) &= 2 - \frac{1}{2} = \frac{3}{2} \\ g(3) &= 2 - \frac{1}{4} = \frac{7}{4} \end{aligned} \right\}$$

VARI NEMENSTVO !!!

• SE KAKVA ŽAM NA UIFORMNATA Ⓢ 113.40

$$Pr\{X \geq t\} = \frac{(n-t+1)}{n}$$

$$\frac{n+1}{2t} \geq \frac{n-t+1}{n}$$

$$n+1 \geq 2t \cdot n - 2t^2 + 2t$$

$$2t^2 - 2(n+1)t + n(n+1) \geq 0$$

ISTOTO SE DONJA KAKO I ZA $Pr\{X > t\}$

• SEGA DA KODARIT EA:

$$X = \left\{ \frac{1}{2^{n-i}} : i=1, \dots, n \right\}$$

$$P = \left\{ \frac{1}{2^i} : i=1, \dots, n-1 \right\}, \frac{1}{2^{n+1}}$$

$$+ \frac{1}{2^{n+1}} \cdot 1$$

$$E[X] = \sum_{i=1}^n x_i \cdot p(i) = \sum_{i=1}^{n-1} \frac{1}{2^{n-i}} \cdot \frac{1}{2^i} + \frac{1}{2^{n+1}} \cdot 1$$

$$E[X] = \frac{n+1}{2^n}$$

$$= \sum_{i=1}^{n-1} \frac{1}{2^n} = \frac{n-1}{2^n} + \frac{1}{2^{n+1}} = \frac{n+1}{2^{n+1}}$$

$$\frac{E[X]}{t} = \frac{n+1}{t \cdot 2^n}$$

$$Pr[X > t] = ? \quad Pr[X > t] = \frac{1}{2^m}$$

$$x = \left[\frac{1}{2^1}, \frac{1}{2^{n-2}}, \dots, \frac{1}{2^2}, \frac{1}{2^1}, 1 \right]$$

$$f(x) = \left[\frac{1}{2}, \frac{1}{2^2}, \dots, \frac{1}{2^n}, \dots, \frac{1}{2^{n-2}}, \frac{1}{2^{n-1}}, \frac{1}{2^n} \right]$$

$$t = \frac{1}{2^{n-m}} \quad t = \frac{1}{2^m} \cdot 2^n \quad \boxed{2^m = t \cdot 2^n}$$

$$\boxed{Pr[X > t] = \frac{1}{t \cdot 2^n}}$$

$$2^m = \frac{1}{256} \cdot 2^{10} = 4$$

$$m = \log_2 4 = 2$$

$$\frac{n+1}{t \cdot 2^n} \geq \frac{1}{t \cdot 2^m} \quad n+1 \geq 1 \quad \boxed{n \geq 0}$$

1 во орд \leq орд \rightarrow \forall $t \in \mathbb{R}^+$ \rightarrow \forall $n \in \mathbb{N}$ \rightarrow \forall $n \geq 0$.

$$X = [1, 2, \dots, t, \dots, n]$$

$$\frac{E[X]}{t} = \frac{1}{t} \frac{n(n+1)}{2n}$$

$$\boxed{\frac{E[X]}{t} = \frac{n+1}{2t}}$$

$$\left. \frac{1}{t}, \frac{1}{t}, \dots, \frac{1}{t} \right\} \text{OPT. 1}$$

$$f = [f_1, f_2, \dots, f_t, \dots, f_n]$$

$$\sum_{i=1}^n f_i = \underbrace{f_t + f_{t+1} + \dots + f_n}_{\frac{n+1}{2} \text{ OPT. 1}} = \frac{n+1}{2 \cdot t}$$

$$f = \left[\frac{1}{n}, \frac{1}{n-1}, \dots, \frac{1}{t}, \dots, \frac{1}{2} \right]$$

$$\boxed{Pr[X \geq t] \leq \frac{E[X]}{t}}$$

$$\sum_{i=t}^n f_i \leq \frac{1}{t} \sum_{i=1}^n f_i \cdot x_i$$

$$t \cdot \sum_{i=t}^n \gamma_i \leq \sum_{i=1}^{t-1} \gamma_i x_i + \sum_{i=t}^n \gamma_i x_i$$

$$\sum_{i=t}^n (\textcircled{t} - \gamma_i) x_i \leq \sum_{i=1}^{t-1} \gamma_i x_i$$

$$\gamma_1 x_1 + \gamma_2 x_2 + \dots + \gamma_{t-1} x_{t-1} + (t - \gamma_t) x_t + (t - \gamma_{t+1}) x_{t+1} + \dots + (t - \gamma_n) x_n \geq 0$$

$$X \in [x_1, x_2, \dots, x_t, \dots, x_n] = [x_1, x_2, \dots, \textcircled{\gamma_t}, \dots, \textcircled{\gamma_t}]$$

$$\gamma \in [\gamma_1, \gamma_2, \dots, \gamma_t, \dots, \gamma_n] = [\gamma_1, \gamma_2, \dots, \gamma_t, \dots, \gamma_n]$$

$$\frac{E[X]}{t} = \sum_{i=1}^{t-1} \gamma_i x_i + \sum_{i=t}^n \textcircled{\gamma_i} x_i = \sum_{i=1}^{t-1} \gamma_i x_i + \gamma_t \sum_{i=t}^n x_i$$

EXERCISE 1 SOLUTION

(a) IF X HAS DISTRIBUTION $F_{\infty}(x)$

$$E[X] = \int_0^{\infty} x dF = \int_0^{\delta} x dF + \int_{\delta}^{\infty} x dF \geq \int_{\delta}^{\infty} x dF$$

$$\geq \int_{\delta}^{\infty} \delta dF = \delta \cdot \Pr[X \geq \delta]$$

$\Pr[X \geq \delta] \leq \frac{E[X]}{\delta}$
 PROOFED!!!

$$E[X] = E[X | X \leq \delta] \cdot \Pr[X \leq \delta] + E[X | X > \delta] \cdot \Pr[X > \delta]$$

$$\textcircled{*} E[X] = E[X | X \geq \delta] \cdot \Pr\{X \geq \delta\} + E[X | X < \delta] \cdot \Pr\{X < \delta\}$$

$$E[X | Y=7] = \sum_{x \in X} x \cdot P(X=x | Y=7) = \sum_{x \in X} x \frac{P(X=x, Y=7)}{P(Y=7)}$$

$$E[X | X \leq \delta] = \sum_{x \in X} x \cdot \frac{P(X=x | X \leq \delta)}{P(X \leq \delta)} = \sum_{x \in X} x \cdot \frac{P(X=x, X \leq \delta)}{P(X \leq \delta)}$$

$$E[X] = \sum_{x \in X} x \cdot \gamma(x) = \sum_{x \in X} x \cdot \sum_n \gamma(x|n)$$

$$\rightarrow E[X] \geq E[X | X \geq \delta] \cdot \Pr\{X \geq \delta\} \geq \delta \cdot \Pr\{X \geq \delta\} \quad \text{JAKO!!!} \quad (43)$$

CONDITIONAL EXPECTATION AND TOTAL EXPECTATION

- IF Z IS RANDOM VARIABLE AND " E " IS EVENT THEN CONDITIONAL EXPECTATION IS DEFINED BY:

$$E_x(Z|E) = \sum_{\omega \in E} Z(\omega) \cdot P_r(\omega|E)$$

- FOR EXAMPLE, LET Z BE THE NUMBER THAT COMES UP ON A ROLL OF A FAIR DIE, AND LET " E " BE THE EVENT THAT NUMBER IS 'EVEN'. LET'S COMPUTE $E_x(Z|E)$, THE EXPECTED VALUE OF A DIE ROLL, GIVEN THAT THE RESULT IS EVEN.

$$E_x(Z|E) = \sum_{\omega \in E} Z(\omega) \cdot P_r(\omega|E) = 1 \cdot 0 + 2 \cdot \frac{1}{3} + 3 \cdot 0 + 4 \cdot \frac{1}{3} + 5 \cdot 0 + 6 \cdot \frac{1}{3} =$$

$$= \frac{2}{3} + \frac{4}{3} + \frac{6}{3} = \frac{12}{3} = 4$$

THEOREM (TOTAL EXPECTATION)

$$E_x(Z) = E_x(Z|E_1) \cdot P_r(E_1) + E_x(Z|E_2) \cdot P_r(E_2) + \dots + E_x(Z|E_n) \cdot P_r(E_n)$$

ANALOGOUS TO TOTAL PROBABILITY THEOREM

$$P(B) = P(B|A_1) \cdot P(A_1) + P(B|A_2) \cdot P(A_2) + \dots + P(B|A_n) \cdot P(A_n)$$

- FOR EXAMPLE, LET Z BE THE NUMBER THAT COMES UP ON FAIR DIE, AND $E =$ THE EVENT THAT THE RESULT IS EVEN. THEN THE \bar{E} IS THE EVENT THAT THE RESULT IS ODD. SO TOTAL EXPECTATION

THEORY SAYS:

$$E_x(Z) = \underbrace{E_x(Z|E)}_4 \cdot \underbrace{P_r(E)}_{1/2} + \underbrace{E_x(Z|\bar{E})}_? \cdot \underbrace{P_r(\bar{E})}_{1/2}$$

171
341
512

$$E_x(Z|\bar{E}) = 1 \cdot \frac{1}{3} + 2 \cdot 0 + 3 \cdot \frac{1}{3} + 4 \cdot 0 + 5 \cdot \frac{1}{3} + 6 \cdot 0$$

$$E_x(Z|\bar{E}) = \frac{1+3+5}{3} = \frac{9}{3} = 3$$

$$E_x(Z) = 4 \cdot \frac{1}{2} + 3 \cdot \frac{1}{2} = 2 + \frac{3}{2} = \frac{4+3}{2} = \frac{7}{2}$$

$E \setminus R$	1	2	3	4	5	6	7	8	9	10	$P(E)$	$P(Z E)$
E	0	1/4	0	1/16	0	1/4	0	1/16	0	1/16	171/512	
\bar{E}	1/2	0	1/8	0	1/32	0	1/128	0	1/512	0	341/512	
$P(R)$	1/2	1/4	1/8	1/16	1/32	1/64	1/128	1/256	1/512	1/512	1	

$P(R|E)$

MOTB
6.001

77.28.101.27

$E \backslash F$	1	2	3	4	5	6	7	8	9	10
E	0	$\frac{512 \cdot 1}{171 \cdot 4}$	0	$\frac{512 \cdot 1}{171 \cdot 16}$	0	$\frac{512 \cdot 1}{171 \cdot 64}$	0	$\frac{512 \cdot 1}{171 \cdot 256}$	0	$\frac{512 \cdot 1}{171 \cdot 1024}$
\bar{E}	$\frac{512 \cdot 1}{341 \cdot 2}$	0	$\frac{512 \cdot 1}{341 \cdot 8}$	0	$\frac{512 \cdot 1}{341 \cdot 32}$	0	$\frac{512 \cdot 1}{341 \cdot 128}$	0	$\frac{512 \cdot 1}{341 \cdot 512}$	0

$$E_x(R) = \underbrace{E_x(R|E)}_{!} \cdot \underbrace{P_R(E)}_{\frac{171}{512}} + \underbrace{E_x(R|\bar{E})}_{?} \cdot \underbrace{P_R(\bar{E})}_{\frac{341}{512}}$$

CONTINUE FROM PP. 43

- For given δ , THE DISTRIBUTION ACHIEVING:
 $P_r\{X \geq \delta\} = \frac{E[X]}{\delta}$ IS: $X = \begin{cases} \delta & \text{WITH PROB. } \frac{\mu}{\delta} \\ 0 & \text{WITH PROB. } 1 - \frac{\mu}{\delta} \end{cases}$
 WHERE $\mu \leq \delta$

$$E[X] = \delta \cdot \frac{\mu}{\delta} + 0 \cdot (1 - \frac{\mu}{\delta}) \quad \boxed{E[X] = \mu}$$

$$P_r\{X \geq \delta\} = \frac{\mu}{\delta} \quad \boxed{\frac{\mu}{\delta} = \frac{\mu}{\delta}} \quad \checkmark$$

- IZGLEBA VARN ZA UTE. POEDJOSTAVNA DISTRO:

ϵ : $X = \begin{cases} \delta & \text{WITH PROBABILITY } \mu \\ 0 & \text{WITH PROBABILITY } 1 - \mu \end{cases}$

$$E[X] = \delta \cdot \mu + 0 \cdot (1 - \mu) = \delta \mu$$

$$P_r\{X \geq \delta\} = \mu \Rightarrow \mu = \frac{\delta \mu}{\delta} \quad \boxed{\mu = \mu} \quad \checkmark$$

(b) Chebyshev INEQUALITY: Let X BE A RANDOM VARIABLE WITH MEAN μ AND VARIANCE σ^2 . BY LEVY-NGI $Z = (X - \mu)^2$, SHOW THAT FOR ANY $\epsilon > 0$

$$P_r\{|X - \mu| > \epsilon\} \leq \frac{\sigma^2}{\epsilon^2}$$

$$\boxed{P_r\{X \geq \delta\} \leq \frac{E[X]}{\delta}}$$

MARKOV INEQUALITY

$$\rightarrow P_r\{|X - \mu|^2 > \epsilon^2\} \leq P_r\{(X - \mu)^2 \geq \epsilon^2\} \leq \frac{E[(X - \mu)^2]}{\epsilon^2}$$

$$E[(X - \mu)^2] = \sigma^2 \Rightarrow P_r\{|X - \mu|^2 > \epsilon^2\} \leq \frac{\sigma^2}{\epsilon^2}$$

$$P_r\{|X - \mu|^2 > \epsilon^2\} = P_r\{(X - \mu) > \epsilon\} \Rightarrow P_r\{(X - \mu) > \epsilon\} \leq \frac{\sigma^2}{\epsilon^2}$$

PROVED !!!

(c) (SPECIAL IN EDITION 1) PROVE THE WEAK LAW OF LARGE NUMBERS. LET Z_1, Z_2, \dots, Z_n BE A SEQUENCE OF I.I.D. RANDOM VARIABLES WITH MEAN μ AND VARIANCE σ^2 . LET $\bar{Z}_n = \frac{1}{n} \sum_{i=1}^n Z_i$ BE A SAMPLE MEAN.

SHOW THAT:

$$P\left\{|\bar{Z}_n - \mu| > \epsilon\right\} \leq \frac{\sigma^2}{n\epsilon^2}$$

THUS $P\left\{|\bar{Z}_n - \mu| > \epsilon\right\} \rightarrow 0$ AS $n \rightarrow \infty$. THIS IS KNOWN AS THE WEAK LAW OF LARGE NUMBERS.

$$E[\bar{Z}_n] = \frac{1}{n} E\left[\sum_{i=1}^n Z_i\right] = \frac{1}{n} \sum_{i=1}^n E[Z_i] = \frac{1}{n} \cdot n \cdot \mu = \mu$$

$$\begin{aligned} \text{Var}[\bar{Z}_n] &= E[(\bar{Z}_n - \mu)^2] = E[\bar{Z}_n^2 - 2\mu\bar{Z}_n + \mu^2] \\ &= E[\bar{Z}_n^2] - 2\mu E[\bar{Z}_n] + \mu^2 = E[\bar{Z}_n^2] - \mu^2 \end{aligned}$$

$$\begin{aligned} E[\bar{Z}^2] &= \frac{1}{n^2} E[(Z_1 + Z_2) \cdot (Z_1 + Z_2)] = \frac{1}{n^2} E[Z_1^2 + 2Z_1Z_2 + Z_2^2] \\ &= \frac{1}{n^2} \left\{ E[Z_1^2] + 2E[Z_1Z_2] + E[Z_2^2] \right\} = \frac{2}{n^2} E[Z_1^2] = \frac{1}{n^2} (\sigma^2 + \mu^2) \end{aligned}$$

FOR $n=1$ $E[Z^2] = \frac{1}{n^2} (\sigma^2 + \mu^2) = \frac{1}{1} (\sigma^2 + \mu^2)$

NE E POSRO VARI 90 AKOUSTAN $\cdot \frac{1}{n}$

$$\sigma^2 = E[(Z_1 - \mu)^2] = E[Z_1^2] - 2\mu E[Z_1] + \mu^2 = E[Z_1^2] - \mu^2$$

$$\text{Var}[\bar{Z}_n] = E\left[\left(\frac{1}{n} \sum_{i=1}^n Z_i - \mu\right)^2\right] = \frac{1}{n^2} E\left[\sum_{i=1}^n Z_i^2\right] - 2\mu E\left[\frac{1}{n} \sum_{i=1}^n Z_i\right] + \mu^2$$

$$\boxed{\text{Var}[\bar{Z}_n] = \frac{1}{n^2} E\left[\sum_{i=1}^n Z_i^2\right] - \mu^2}$$

VARI E POSRO!!!

$$E\left[\left(\sum_{i=1}^n Z_i\right)^2\right] = E[(Z_1 + Z_2)(Z_1 + Z_2)] = E[Z_1^2] + 2E[Z_1Z_2] + E[Z_2^2]$$

$$E\left[\left(\sum_{i=1}^n Z_i\right)^2\right] = 2 \cdot (\sigma^2 + \mu^2) \quad \text{ANALOGNO: } E\left[\left(\sum_{i=1}^n Z_i\right)^2\right] = n \cdot (\sigma^2 + \mu^2)$$

$$\text{Var}[\bar{Z}_n] = \frac{1}{n^2} n \cdot (\sigma^2 + \mu^2) - \mu^2 = \frac{\sigma^2}{n} + \frac{\mu^2}{n} - \mu^2$$

$$\bar{z}_n = \frac{1}{n} \sum_{i=1}^n z_i$$

$$E[\bar{z}_n] = \mu \quad E[(z_i - \mu)^2] = \sigma^2 \quad (*)$$

$$E[\bar{z}_n] = \frac{1}{n} \sum_{i=1}^n E[z_i] = \frac{1}{n} \cdot n \cdot \mu = \mu$$

~~$$E[(\bar{z}_n - \mu)^2] = E[\bar{z}_n^2] - 2\mu E[\bar{z}_n] + \mu^2 = E[\bar{z}_n^2] - \mu^2$$~~

~~$$E[(\bar{z}_n - \mu)^2] = E\left[\frac{1}{n^2} \left(\sum_{i=1}^n z_i\right)^2\right] = \frac{1}{n^2} \cdot n \cdot E[z_i^2]$$~~

$$E[(\bar{z}_n - \mu)^2] = E\left[\left(\frac{1}{n} \sum_{i=1}^n z_i - \frac{1}{n} \sum_{i=1}^n \mu\right)^2\right] =$$

$$= E\left[\frac{1}{n} \sum_{i=1}^n (z_i - \mu)\right]^2 = \frac{1}{n^2} \sum_{i=1}^n E[(z_i - \mu)^2] = \frac{1}{n} \cdot \sigma^2$$

STATISTI. INDEP.

OVA E DOPO VERIFICANZE!!! MMV

$$[(z_1 - \mu) + (z_2 - \mu)] [(z_1 - \mu) + (z_2 - \mu)] =$$

$$(z_1 - \mu)^2 + 2(z_1 - \mu) \cdot (z_2 - \mu) + (z_2 - \mu)^2$$

$$E[(z_1 - \mu) \cdot (z_2 - \mu)] = E[z_1 z_2] - \mu \cdot E[z_1] - \mu \cdot E[z_2] + \mu^2 = 0$$

- OVA FOLE KACUVA NEJEDNO DRUGO NO MSLAJI OČETA
 POJAVO E TOA ŠTO SU NA STATISTIČKI NEZAVISNI
 (KOMENČIVI ZVUKIŠI PORAVANJE ICI OZEMENJE SA KOMI-
 KITA NE OZEMENJE STATISTIČKI NEZAVISNI TA
 ZAVOJA: $E[(z_1 - \mu) \cdot (z_2 - \mu)] = 0$

$$z = x + y \quad E[x] = \mu_x \quad E[y] = \mu_y \quad E[(x - \mu_x)^2] = \sigma_x^2$$

$$E[(y - \mu_y)^2] = \sigma_y^2 \quad \sigma_x^2 = E[x^2] - \mu_x^2$$

$$E[z] = E[x + y] = E[x] + E[y] = \mu_x + \mu_y$$

$$E[(z - \mu_z)^2] = E[z^2] - 2\mu_z \cdot E[z] + \mu_z^2 = E[z^2] - \mu_z^2$$

$$E[z^2] = E[x^2 + 2xy + y^2] = E[x^2] + E[y^2] = \sigma_x^2 + \mu_x^2 + \sigma_y^2 + \mu_y^2$$

$$E[(z - \mu_z)^2] = \sigma_x^2 + \mu_x^2 + \sigma_y^2 + \mu_y^2 - (\mu_x + \mu_y)^2 =$$

$$\sigma_z^2 = \sigma_x^2 + \sigma_y^2 - 2\mu_x \mu_y$$

1 VO
 TRONKA NA
 INFO.M.
 ZA NOLNAR A DUTRI
 DUCIA IMA DOBA
 DOVA OČETA $\sigma_z^2 = \sigma_x^2 + \sigma_y^2$
 JO KAZIVANJE ICI
 PUKIČI (UPI WIKIČI) 47

CHEBYSHEV INEQUALITY:

$$P\{|Y - \mu| > \varepsilon\} \leq \frac{\sigma^2}{\varepsilon^2} \quad \bar{z}_n = \frac{1}{n} \sum_{i=1}^n z_i$$

$$E[\bar{z}_n] = \mu \quad \text{Var}\{\bar{z}_n\} = \frac{\sigma^2}{n}$$

$$P\{| \bar{z}_n - \mu | > \varepsilon\} \leq \frac{\sigma^2}{n \cdot \varepsilon^2}$$

BY LETTING $n \rightarrow \infty \rightarrow P\{| \bar{z}_n - \mu | > \varepsilon\} \rightarrow 0$
PROVED!!!

5.2 ACI AND MUTUAL INFORMATION. LET (X_i, Z_i)

BE I.I.D $\sim p(x, y)$. WE FORM THE LOG LIKELIHOOD RATIO OF THE HYPOTHESIS THAT X AND Y ARE INDEPENDENT VS. HYPOTHESIS THAT X AND Y ARE DEPENDENT. WHAT IS THE LIMIT OF:

$$I = \frac{1}{n} \log \frac{p(x^n) p(y^n)}{p(x^n, y^n)}$$

$$-\frac{1}{n} \log p(x_1, x_2, \dots, x_n) = -\frac{1}{n} \sum_{i=1}^n \log p(x_i) \rightarrow -E[\log p(x_i)] = H(X)$$

SAME NOISE

$$p(x_1, x_2, \dots, x_n) \rightarrow 2^{-nH(X)}$$

$$I = \frac{1}{n} \log p(x^n) + \frac{1}{n} \log p(y^n) - \frac{1}{n} \log p(x^n, y^n) \xrightarrow{n \rightarrow \infty}$$

$$\rightarrow -H(X) - H(Y) + H(X, Y) = -H(X) - H(Y) + (H(X) + H(X|Y))$$

$$= -H(X) + H(X|Y) = \underline{\underline{-I(X, Y)}}$$

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log \frac{p(x^n) \cdot p(y^n)}{p(x^n, y^n)} = I(X, Y)$$

CONCEPTUAL UNIVERSAL SECTION n

$$\frac{1}{n} \log \frac{p(x^n) \cdot p(y^n)}{p(x^n, y^n)} = \frac{1}{n} \log \prod_{i=1}^n \frac{p(x_i) \cdot p(y_i)}{p(x_i, y_i)} =$$

$$= \frac{1}{4} \sum_{i=1}^4 \lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{f(x_i) \cdot g(y_i)}{f(x_i, y_i)} \cdot \lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{f(x_i) \cdot g(y_i)}{f(x_i, y_i)} =$$

$$= \sum_{i=1}^4 \lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{f(x_i) \cdot g(y_i)}{f(x_i, y_i)} = \sum_{i=1}^4 I(x_i, y_i) \quad \text{VATKRO}$$

$$\frac{1}{4} \sum_{i=1}^4 \ln \frac{f(x_i) \cdot g(y_i)}{f(x_i, y_i)} \rightarrow E \left[\ln \frac{f(x_i) \cdot g(y_i)}{f(x_i, y_i)} \right] = -I(x, y) \quad \text{TROVED}$$

3.3 Piece of Cake. A cake is sliced roughly in half, the other pieces discarded. We will assume that a random cut creates pieces of proportion

$$p = \begin{cases} \left(\frac{2}{3}, \frac{1}{3}\right) & \text{WITH PROBABILITY } \left(\frac{3}{4}\right) = p \\ \left(\frac{2}{5}, \frac{3}{5}\right) & \text{WITH PROBABILITY } \left(\frac{1}{4}\right) = 1-p \end{cases}$$

THUS FOR EXAMPLE, THE FIRST CUT (AND CHANCE OF LARGEST PIECE) MAY RESULT IN PIECE SIZE $\frac{3}{5}$. CUTTING AND CHOOSING FROM THIS PIECE MIGHT REDUCE IT TO SIZE $\left(\frac{2}{5}\right) \cdot \left(\frac{2}{3}\right)$ AT THAT 2, AND SO ON. HOW LARGE TO FIRST OFFER IN EXHONENT, IS THE PIECE OF CAKE AFTER "n" CUTS.

$$1) \frac{1}{4} \cdot \frac{1}{2} = P\left(\frac{3}{5}\right) = \frac{1}{8} \quad P\left(\frac{2}{5}\right) = \frac{1}{4} \cdot \frac{1}{2} = \frac{1}{8}$$

$$2) P\left(\frac{3}{5} \cdot \frac{2}{3}\right) = P\left(\frac{2}{5}\right) \cdot P\left(\frac{2}{3}\right) = \frac{1}{8} \cdot \frac{3}{4} \cdot \frac{1}{2} = \frac{1}{8} \cdot \frac{3}{8} = \frac{3}{64}$$

$$P(2) = P\left[\frac{1}{2} P\left(\frac{2}{3}\right) + \frac{1}{2} P\left(\frac{1}{3}\right)\right]$$

$$P\left(\frac{3}{15}\right) = P\left(\frac{2}{5} \cdot \frac{1}{3}\right) = P\left(\frac{2}{5}\right) \cdot P\left(\frac{1}{3}\right) = \frac{1}{8} \cdot \frac{3}{4} \cdot \frac{1}{2} = \frac{3}{64} = P\left(\frac{2}{5} \cdot \frac{2}{3}\right) = P\left(\frac{4}{15}\right)$$

$$P\left(\frac{4}{15}\right) = P\left(\frac{2}{5} \cdot \frac{2}{3}\right) = P\left(\frac{2}{5}\right) \cdot P\left(\frac{2}{3}\right) = \frac{1}{8} \cdot \frac{3}{4} \cdot \frac{1}{2} = \frac{3}{64} = P\left(\frac{2}{5} \cdot \frac{1}{3}\right) = P\left(\frac{2}{15}\right)$$

- AND SEE HOW TO DO IT

$$P\left(\frac{2}{3}\right) = \frac{3}{4} \cdot \frac{1}{2} = \frac{3}{8} \quad P\left(\frac{1}{3}\right) = \frac{3}{4} \cdot \frac{1}{2} = \frac{3}{8}$$

$$P\left(\frac{4}{15}\right) = P\left(\frac{2}{3} \cdot \frac{2}{5}\right) = P\left(\frac{2}{3}\right) \cdot P\left(\frac{2}{5}\right) = \frac{3}{8} \cdot \frac{1}{8} = \frac{3}{64} = P\left(\frac{1}{3} \cdot \frac{2}{5}\right) = P\left(\frac{2}{15}\right)$$

$$P\left(\frac{2}{15}\right) = P\left(\frac{2}{3} \cdot \frac{2}{5}\right) = P\left(\frac{2}{3}\right) \cdot P\left(\frac{2}{5}\right) = \frac{3}{64} = P\left(\frac{1}{3} \cdot \frac{2}{5}\right) = P\left(\frac{2}{15}\right)$$

• АНО СЛУЧАЈИ ГО ИЗДИКАЈИ ПОСОЛЕМОТО ПАКЕ ПРОЦЕНУО СЕ ПОСМОТРЕВА:

$$P = \begin{cases} \left(\frac{2}{3} | \frac{1}{3}\right) & \text{WITH PROBABILITY } \left(\frac{3}{4}\right) \cdot P'' \\ \left(\frac{2}{5} | \frac{3}{5}\right) & \text{WITH } \dots \left(\frac{1}{4}\right) \cdot Z'' \end{cases}$$

• АНО СЕ ТРЕЊЕ ОД $\left(\frac{2}{3}\right)$

1) $P\left(\frac{2}{3}\right) = P\left(\frac{3}{4}\right) \cdot \frac{1}{2} = \frac{3}{4} \cdot \frac{1}{2} = \frac{3}{8} = \frac{1}{2}$ } СЕКОЈ ГО ДИКАЈИ ПОСОЛЕМОТО Т.Е. $\frac{2}{3}$

$P\left(\frac{1}{3}\right) = \left(\frac{1}{4}\right) \cdot \frac{1}{2} = \frac{1}{8} = \frac{1}{2}$

2) $P\left(\frac{2}{3} \cdot \frac{3}{5}\right) = P\left(\frac{2}{3}\right) \cdot P\left(\frac{3}{5}\right) = \frac{3}{8} \cdot \frac{1}{4} \cdot \frac{1}{2} = \frac{3}{8} \cdot \frac{1}{8} = \frac{3}{64}$ } ГО ДИКАЈИ ПОСОЛЕМОТО $\frac{3}{5}$

$P\left(\frac{2}{3} \cdot \frac{2}{5}\right) = P\left(\frac{2}{3}\right) \cdot P\left(\frac{2}{5}\right) = \frac{3}{8} \cdot \frac{1}{8} = \frac{3}{64}$

$P\left(\frac{2}{3} \cdot \frac{2}{3}\right) = P\left(\frac{2}{3}\right) \cdot P\left(\frac{2}{3}\right) = \frac{3}{8} \cdot \frac{3}{8} = \frac{9}{64}$

$P\left(\frac{2}{3} \cdot \frac{1}{3}\right) = P\left(\frac{2}{3}\right) \cdot P\left(\frac{1}{3}\right) = \frac{3}{8} \cdot \frac{1}{8} = \frac{3}{64}$ } ГО ДИКАЈИ ПОСОЛЕМОТО $\frac{2}{3}$

$P\left(\frac{2}{3} \cdot \frac{3}{5}\right) = P\left(\frac{6}{15}\right) = \frac{1}{2} \cdot \frac{2}{2} = \frac{1 \cdot 2}{2^2}$ } ВО ВТОРОТ СЕКОЈ ГО ДИКАЈИ ИЛИ $P\left(\frac{6}{15}\right)$ ИЛИ $P\left(\frac{4}{9}\right)$

• АНО СЕ ТРЕЊЕ ОД $\left(\frac{3}{5}\right)$ ВО ВТОРА СЕКОЈ ОДИТЕ СЕ:

$P\left(\frac{3}{5} \cdot \frac{2}{3}\right) = \frac{1 \cdot 2}{2^2}$ } $P\left(\frac{2}{5} \cdot \frac{3}{5}\right) = \frac{2}{2^2}$

БУТРОИПАТА ЗА ВТОРОТ СЕКОЈ Е:

$$H = - \sum_{k=0}^2 \binom{4}{k} p^k \cdot 2^{4-k} \log \binom{4}{k} p^k \cdot 2^{4-k} =$$

$$= - \left[\binom{4}{0} p^0 \cdot 2^4 \log \binom{4}{0} p^0 \cdot 2^4 + \binom{4}{1} p^1 \cdot 2^3 \log \binom{4}{1} p^1 \cdot 2^3 + \binom{4}{2} p^2 \cdot 2^2 \log \binom{4}{2} p^2 \cdot 2^2 \right]$$

ВО ОБЕДИНЕН СЛУЧАЈ БУТРОИПАТА СЕ ВИДЕ:

$$H(S) = - \sum_{k=0}^4 \binom{4}{k} p^k \cdot 2^{4-k} \log \binom{4}{k} p^k \cdot 2^{4-k}$$

$S = \{x_1, x_2, \dots, x_n\}$
 БУТРОИПА НА НИСТА ОД СЛУЧАЈИ ПЛОТ.

МОЗЕТИ НА ТРЕТА ПА МОЗЕТИ СО $\frac{1}{2}$ БОРТО ПРОВИКО Е ОН ГО ДИКАЈИ ПОСОЛЕМОТО ПАКЕ !!!

30) $P\left(\frac{2}{3} \cdot \frac{2}{3} \cdot \frac{2}{3}\right) = \left| \begin{array}{l} \text{OPAM ISTE} \\ \text{PEKENTE SO} \\ \text{2 ZALADJANJE} \\ \text{NA ROZLOZENJE VALICINA} \end{array} \right| = P \cdot 2 \cdot 1 = \underline{1}$

$P\left(\frac{2}{3} \cdot \frac{2}{3} \cdot \frac{2}{3}\right) = 1^3$; $P\left(\frac{2}{3} \cdot \frac{2}{3} \cdot \frac{2}{3}\right) = \underline{1 \cdot 2^2}$;

$P\left(\frac{2}{3} \cdot \frac{2}{3} \cdot \frac{2}{3}\right) = \underline{1^2 \cdot 2}$; $P\left(\frac{2}{3}, \frac{2}{3}, \frac{2}{3}\right) = \underline{2 \cdot 1^2}$;

$P\left(\frac{2}{3} \cdot \frac{2}{3} \cdot \frac{2}{3}\right) = \underline{2^2 \cdot 1}$; $P\left(\frac{2}{3} \cdot \frac{2}{3} \cdot \frac{2}{3}\right) = \underline{2^2 \cdot 1}$; $P\left(\frac{2}{3}, \frac{2}{3}, \frac{2}{3}\right) = \underline{2^3}$

IMAŠ VUKRO: $1+1^3$, $3 \times 1^2 \cdot 2$, $3 \times 1 \cdot 2^2$, $1+2^3$
 KOENI ISTORI.

$P\left(\frac{9}{64}\right)$; $3 \cdot P\left(\frac{9}{16} \cdot \frac{1}{4}\right) = 3 \cdot P\left(\frac{3}{64}\right)$; $3 \cdot P\left(\frac{3}{4} \cdot \frac{1}{16}\right) = 3 \cdot P\left(\frac{3}{64}\right)$;

$P\left(\left(\frac{1}{4}\right)^2\right) = P\left(\frac{1}{64}\right)$

$X_i \in \begin{cases} \frac{2}{3} & \text{PROBABILITY } 1 \\ \frac{2}{5} & \text{PROBABILITY } 2 = 1 - 1 \end{cases}$

$P(x_1, x_2, \dots, x_i, \dots, x_n) \rightarrow 2^{-4+H(x)} = 2^{-4 \cdot (1 \cdot \log \frac{1}{1} + 2 \cdot \log \frac{1}{2})}$

$H(x) = 1 \cdot \log \frac{1}{1} + 2 \cdot \log \frac{1}{2} = \frac{2}{4} \log \frac{4}{3} + \frac{1}{4} \log 4 = \underline{0.8113}$

• TIPICNIOT SET JE IMA

- $4 \cdot \frac{3}{4}$ $\frac{2}{3}$ -KI OTSEČOCI
(DVOLETINSKI)

- $4 \cdot \frac{1}{4}$ $\frac{3}{5}$ -KI (TRILETINSKI) OSEČOCI

$\begin{aligned} L_n x &= x \\ L_n x &= L_n x \end{aligned}$

ZA DOKAZO SAKERU "4" $P_V \{A \in \mathcal{A}\} \geq 1 - \epsilon$

PRIMER: $n=4$ \Rightarrow NAJVEROJATNIJA GOLEMINA NA
 VIKOTO JE DINE

$\left(\frac{2}{3}\right)^{\frac{3 \cdot 4}{4}} \cdot \left(\frac{3}{5}\right)^{\frac{4}{4}} = \left(\frac{2}{3}\right)^3 \cdot \left(\frac{3}{5}\right)^1 = \frac{8}{27} \cdot \frac{3}{5} = \frac{24}{135} = \frac{8}{45}$

$P\left(\left(\frac{2}{3}\right)^3 \cdot \left(\frac{3}{5}\right)^1\right) = 2^{-4H(x)} = 2^{-4 \cdot 0.8113}$

POKAZ DEKA
 $P^n \cdot 2^{-n \cdot 2} = 2^{-4 \cdot H(x)}$

$2^{-4 \left(\frac{3}{4} \log \frac{4}{3} + \frac{1}{4} \log 4\right)} = 2^{-4 \left(\frac{3}{4} \left(\log 4 + \frac{1}{3}\right)\right)} = 2^{-3 \log \frac{4}{3} - 2 \cdot 2 - 3 \log \frac{4}{3}}$
 $= \frac{1}{4} \cdot 2 \log \left(\frac{4}{3}\right)^3 = 2^{-1} \cdot \left(\frac{4}{3}\right)^3 = 2$

• Showed mean and variance of binomial distribution
 "n" = number of trials

$$L(n) = \left(\frac{2}{3}\right)^{\frac{3}{4}n} \cdot \left(\frac{3}{5}\right)^{\frac{1}{4}n}$$

$$L(4) = \frac{8}{45} = 0.178 \quad L(10) = 0.0133$$

$$2^{n \cdot H} = |x=4| = 2^{4 \cdot 0.8} = 9.48$$

$H(4) = 0.8113$
 MAXIMUM
 PROBABILITY
 IN BINOMIAL DISTRIBUTION

$$\binom{4}{1} = \frac{4!}{3!} = \frac{4 \cdot 3 \cdot 2 \cdot 1}{3 \cdot 2 \cdot 1} = 4$$

3.4

AEP. Let x_i be iid $\sim \gamma(x)$, $x \in \{1, 2, \dots, 4\}$
 Let $\mu = E[X]$ AND $H = -\sum \gamma(x) \log \gamma(x)$

Let $A^n = \{x^n \in X^n : \left| -\frac{1}{n} \log \gamma(x^n) - H \right| \leq \epsilon\}$

Let $B^n = \left\{ x^n \in X^n : \left| \frac{1}{n} \sum_{i=1}^n x_i - \mu \right| \leq \epsilon \right\}$ • LAW OF LARGE NUMBERS

- (a) Does $\Pr\{x^n \in A^n\} \rightarrow 1$ Yes
 (b) Does $\Pr\{x^n \in A^n \cap B^n\} \rightarrow 1$ Yes

(a) A^n is typical macrostate and hence $\Pr\{A^n\} \leq 1 - \epsilon$ Theorem 3.2.2

(b) B^n - smallest set $\Pr\{B^n\} \geq 1 - \delta$

(c) Show that: $|A^n \cap B^n| \leq 2^{n(H+\epsilon)}$ FOR $\forall n$

(d) Show that: $|A^n \cap B^n| \geq \left(\frac{1}{2}\right) \cdot 2^{n(H-\epsilon)}$ FOR n SUFFICIENTLY LARGE

(c) $|A^n| \leq 2^{n(H+\epsilon)}$ A^n is typical set
 $|A^n \cap B^n| \leq |A^n|$ $|A^n| \geq |A^n \cap B^n|$
 $2^{n(H+\epsilon)} \geq |A^n| \geq |A^n \cap B^n|$ $|A^n \cap B^n| \leq 2^{n(H+\epsilon)}$

(d) $|A^n| \geq (1-\epsilon) 2^{n(H-\epsilon)}$
 $1 = \sum_{x^n \in X^n} \gamma(x^n) \geq \sum_{x^n \in A^n} \gamma(x^n) = \sum_{x^n \in A^n} \gamma(x^n) \geq 2^{n(H-\epsilon)}$ $\gamma(x^n) \geq 2^{n(H-\epsilon)}$ \Rightarrow $\gamma(x^n) \leq 2^{-n(H-\epsilon)}$

$$\textcircled{1} = \sum_{x^y \in A^y} \gamma(x^y) \geq \sum_{x^y \in A^y} 2^{-y(n+\epsilon)} = |A^y| 2^{-y(n+\epsilon)}$$

$$|A^y| \leq \sum_{x^y \in X^y} \gamma(x^y) = 1$$

$$1 = \sum_{x^y \in X^y} \gamma(x^y) \leq \sum_{x^y \in X^y} 2^{-y(n-\epsilon)} \leq |X^y| \cdot 2^{-y(n-\epsilon)}$$

$$|X^y| \geq 2^{y(n-\epsilon)}$$

$$|A^y \cap B^y| \leq |A^y|$$

$$|A^y| \geq 2^{y(n-\epsilon)} - \epsilon 2^{y(n-\epsilon)}$$

$$\frac{2^{y(n-\epsilon)}}{2^{y(n-\epsilon)}} \leq \frac{|A^y| + \epsilon 2^{y(n-\epsilon)}}{2^{y(n-\epsilon)}} \Rightarrow \frac{1}{2^{y(n-\epsilon)}} \geq \frac{1}{|A^y| + \epsilon 2^{y(n-\epsilon)}}$$

$$\gamma(x^y) \geq \frac{1}{|A^y| + \epsilon 2^{y(n-\epsilon)}} \quad 2^{-y(n-\epsilon)} \geq \frac{1}{|A^y| + \epsilon 2^{y(n-\epsilon)}}$$

$$|A^y| 2^{-y(n-\epsilon)} + \epsilon \geq 1 \quad |A^y| \geq (1-\epsilon) 2^{y(n-\epsilon)}$$

$$(1+\epsilon) \geq \Pr\{X^{(n)}\} \geq \Pr\{A^{(n)}\} = \sum_{x^y \in A^y} \gamma(x^y) \geq \sum_{x^y \in A^y} 2^{-y(n+\epsilon)}$$

$$= |A^y| 2^{-y(n+\epsilon)} \quad (1+\epsilon) \geq |A^y| 2^{-y(n+\epsilon)}$$

$$|A^y| \leq (1+\epsilon) 2^{y(n+\epsilon)}$$

$$|A^y \cap B^y| \leq |A^y| \leq (1+\epsilon) 2^{y(n+\epsilon)}$$

$$(1+\epsilon) 2^{y(n+\epsilon)} \geq (1-\epsilon) 2^{y(n-\epsilon)} \quad (1+\epsilon) 2^{4\epsilon} \geq (1-\epsilon) 2^{-4\epsilon}$$

$$(1+\epsilon) 2^{24\epsilon} \geq 1-\epsilon$$

$$2^{24\epsilon} \geq \frac{1-\epsilon}{1+\epsilon}$$

$$2^{4\epsilon} \geq \sqrt{\frac{1-\epsilon}{1+\epsilon}}$$

$$2^{-4\epsilon} \leq \sqrt{\frac{1+\epsilon}{1-\epsilon}}$$

$$|A^y| \geq (1-\epsilon) 2^{4n} \cdot 2^{-4\epsilon} \quad (1-\epsilon) \leq (1+\epsilon) 2^{24\epsilon}$$

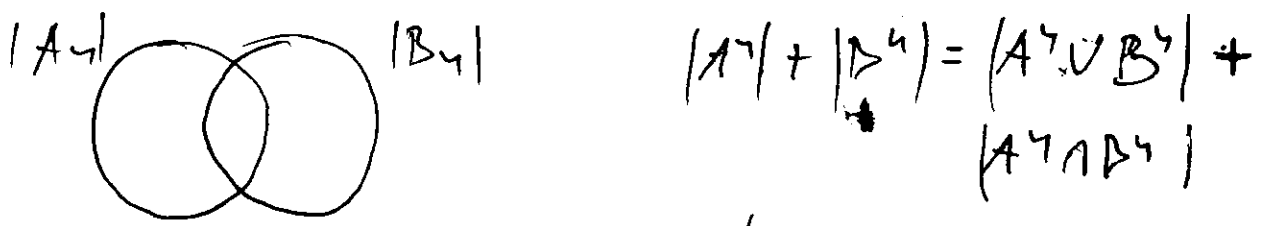
UNIVERSITY OF TORONTO SOLUTIONS:
 $\Pr\{x^y \in A^y \cap B^y\} = \Pr\{x^y \in A^y\} + \Pr\{x^y \in B^y\} - \Pr\{x^y \in A^y \cup B^y\}$
 $> (1-\epsilon) + (1-\epsilon_1) - \Pr\{x^y \in A^y \cup B^y\} \geq 1-\epsilon-\epsilon_1 \Rightarrow \Pr\{x^y \in A^y \cap B^y\} \geq 1-\epsilon-\epsilon_1$

→ STEP 1 OF THEOREM 3.1.2.2

(c) (ΣΥΜΠΛΗΡΩΣΤΙΚΟ ΟΥΚΑΝΟ ΣΤΟ ΒΥΡΟΝ ΚΕΙΜΕΝΟ ΚΑΤ' ΕΞΕΛΙΞΗ)

$|A^c \cap B^c| \leq 2^4(1+\epsilon)$ ΠΑΡΑ ΣΗΜΑΝΣΗ ΜΕΤΡΩΣΤΩΝ

$|A^c \cap B^c| = |A^c| + |B^c| - |A^c \cup B^c|$



~~$|A^c \cap B^c| = |A^c| + |B^c| - |A^c \cup B^c|$~~

$|A^c| + |B^c| \geq |A^c \cap B^c|$ $|A^c \cap B^c| \leq |A^c| + |B^c| \leq |A^c \cup B^c|$

$|A^c \cap B^c| = |A^c| + |B^c| - |A^c \cup B^c| = |B^c| \leq |A^c \cup B^c|$
 $\leq |A^c| + |A^c \cup B^c| - |A^c \cup B^c| = |A^c|$

$|A^c \cap B^c| \leq |A^c|$

ΠΟΥΛΕ ΚΑ ΛΟΜΑΡΑ
 ΚΑ ΕΤΤΟ ΤΑΥΤΑ !!!

ΣΕΓΑ ΚΑΤΙ ΜΕΤΡΩΝΟΤ ΠΟΥΛΕ. ΣΟΦΙΣΜΟ 3.1.2.3

$|A^c| \leq 2^4(1+\epsilon) \Rightarrow |A^c \cap B^c| \leq |A^c| \leq 2^4(1+\epsilon)$

$|A^c \cap B^c| \leq 2^4(1+\epsilon)$

(d) $|A^c \cap B^c| \geq \frac{1}{2} 2^4(1-\epsilon)$ $|A^c| \geq (1-\epsilon) 2^4(1-\epsilon)$

$|A^c \cap B^c| = |A^c| + |B^c| - |A^c \cup B^c| \geq \frac{|A^c \cup B^c| \leq 2|B^c|}{|A^c| \geq |B^c|}$
 $\geq |A^c| + |B^c| - 2|B^c| = |A^c| - |B^c|$

$|A^c| \geq (1-\epsilon) 2^4(1-\epsilon) \quad |Pr\{A^c\}| \geq (1-\epsilon)$

$|A^c| \geq$

(c) UNIVERSITY OF TORONTO SOLUTION

$$1 \geq \sum_{x^N \in A^N \cap B^N} \gamma(x^N) \geq \sum_{x^N \in A^N \cap B^N} 2^{-N(H+\epsilon)} = |A^N \cap B^N| \cdot 2^{-N(H+\epsilon)}$$

$$\Rightarrow |A^N \cap B^N| \leq 2^{N(H+\epsilon)} \quad \text{PROVED!!!}$$

(d) UNIVERSITY OF TORONTO SOLUTIONS

OB (c) ELEMENTARY BENA AND $|A^N \cap B^N| \leq 2^{N(H+\epsilon)}$

$|C^N| \leq 2^{N(H+\epsilon)} \Rightarrow C^N \in \text{TIPICO MODO}$

$\rightarrow P\{C^N\} \rightarrow 1 \Rightarrow P\{A^N \cap B^N\} \rightarrow 1$

$$\frac{1}{2} \leq P\{A^N \cap B^N\} = \sum_{x^N \in A^N \cap B^N} \gamma(x^N) \leq \sum_{x^N \in A^N \cap B^N} 2^{-N(H-\epsilon)} = |A^N \cap B^N| 2^{-N(H-\epsilon)}$$

$$|A^N \cap B^N| \geq \frac{1}{2} 2^{N(H-\epsilon)} \quad \text{PROVED!!!}$$

3.5 SETS DEFINED BY MODALITIES. Let X_1, X_2, \dots be an i.i.d. sequence of discrete random variables with entropy $H(X)$. Let

$$C_N(t) = \{x^N \in \mathcal{X}^N : \gamma(x^N) \geq 2^{-Nt}\}$$

Denote the subset of N -sequences with modalities $\geq 2^{-Nt}$

- (a) SHOW THAT $|C_N(t)| \leq 2^{Nt}$
- (b) FOR WHAT VALUES OF "t" DOES $P\{x^N \in C_N(t)\} \rightarrow 1$

(a) $2^{-N(H(X)+\epsilon)} \leq \gamma(x^N) \leq 2^{-N(H(X)-\epsilon)} \quad / \text{ld}$

$$1 = \sum_{x^N \in \mathcal{X}^N} \gamma(x^N) \geq \sum_{x^N \in C_N(t)} \gamma(x^N) \geq |C_N(t)| 2^{-Nt} \Rightarrow$$

$$|C_N(t)| \leq 2^{Nt}$$

(b) $t = H(X) + \epsilon \Rightarrow |C_N(t)| \leq 2^{N(H(X)+\epsilon)}$

$-N(H(X)+\epsilon) \leq \text{ld } \gamma(x^N) \leq -N(H(X)-\epsilon) \quad (-1)$

$N(H(X)+\epsilon) \geq -\text{ld } \gamma(x^N) \geq N[H(X)-\epsilon] \Rightarrow$

$$\ln(H(x) - \epsilon) \leq -\frac{1}{4} \ln(p(x)) \leq \ln(H(x) + \epsilon)$$

$$1 - \epsilon = \sum_{x \in \mathcal{X}^n} p(x^n) \geq \sum_{x \in \mathcal{X}^n} 2^{-nt} = |\mathcal{C}_n(t)| \cdot 2^{-nt}$$

$$2^{-nt} \leq \frac{1 - \epsilon}{|\mathcal{C}_n(t)|} \quad ; \quad -nt \leq \ln \frac{1 - \epsilon}{|\mathcal{C}_n(t)|} \quad ;$$

$$t \geq \frac{1}{n} \ln \frac{|\mathcal{C}_n(t)|}{1 - \epsilon}$$

• UNIVERSITY OF TORONTO SOLUTIONS

$$\frac{1}{n} \ln |\mathcal{B}_\delta^{(n)}| > H - \delta \quad (\mathcal{B}_\delta^{(n)})^{1/n} > 2^{(H - \delta)}$$

$$|\mathcal{B}_\delta^{(n)}| > 2^{n(H - \delta)} \Rightarrow |\mathcal{C}_n(t)| > 2^{n(H - \epsilon)}$$

- FROM PART (a) $|\mathcal{C}_n(t)| \leq 2^{nt} \Rightarrow$

$$t > H - \epsilon$$

IS NECESSARY CONDITION SUCH $\Pr\{x^n \in \mathcal{C}_n(t)\} \rightarrow 1$

$\{t \geq H + \epsilon$ IS SUFFICIENT CONDITION ($\epsilon > 0$)
 SUCH $A_\epsilon^{(n)} \subseteq \mathcal{C}_n(t)$ WHILE $\Pr\{x^n \in \mathcal{C}_n(t)\} \rightarrow 1$

□ WHETHER $t = H$ IS SUFFICIENT CONDITION (IT DEPENDS TO THE DISTRIBUTION).

□ E.G. $\Pr\{x=0\} = \Pr\{x=1\} = \frac{1}{2} \quad t = H = 1$

$$H = \frac{1}{2} \ln 2 + \frac{1}{2} \ln 2 = 1$$

ALL SEQUENCES HAVE PROBABILITY 2^{-n}

x_1	00
x_2	01
x_3	10
x_4	11

$$P(x_i) = \frac{1}{2^n} = \frac{1}{4} \quad i = 1, 2, 3, 4$$

MINIMUM ENTROPY !!!

□ AND $t = H$

$$\Pr\{x=0\} = \gamma \quad \Pr\{x=1\} = 1 - \gamma \quad 0 < \gamma < 1$$

IS NOT SUFFICIENT !!!

- NA 14/162 ZA
 $x_1 = 00$
 $x_2 = 01$
 $x_3 = 10$
 $x_4 = 11$

$n = 2$
 $p(t_1) = p^2 \cdot p(t_2) = p(1-p)$
 $p(t_3) = p(1-p) \quad p(t_4) = (1-p)^2$

EG. $p = 0.6 \quad q = 0.4 \quad p(t_1) = (0.6)^2 = 0.36$
 $p(t_2) = 0.6 \cdot 0.4 = 0.24 \quad p(t_3) = 0.24 \quad p(t_4) = 0.16$

$H = -0.6 \log 0.6 - 0.4 \log 0.4 = 0.971$
 $2^{-nH} = 2^{-2H} = 2^{-1.942} = \underline{\underline{0.26}}$

$p(t_1) > 0.26$
 $p(t_2 \dots t_4) < 0.26$

- ZNAZI NEODHOVEN SLOV PR $\{x_n \in C_n(t)\} \rightarrow 1 \in$
 $t > t - \epsilon$

36 AEP-like limit. Let x_1, x_2, \dots be i.i.d
 DRAWN ACCORDING TO PROBABILITY MASS FUNCTION
 $p(x)$. Find: $\lim_{n \rightarrow \infty} [p(x_1, x_2, \dots, x_n)]^{1/n}$

$2^{-n(H+\epsilon)} \leq p(x_1, x_2, \dots, x_n) \leq 2^{-n(H-\epsilon)}$

$-\frac{1}{n} \log p(x_1, x_2, \dots, x_n) \rightarrow H(x)$
 $\frac{1}{n} \log p(x_1, x_2, \dots, x_n) \rightarrow -H(x)$

$2^{-n(H+\epsilon)} \leq p(x_1, x_2, \dots, x_n) \leq 2^{-n(H-\epsilon)}$

$\lim_{n \rightarrow \infty} p(x_1, x_2, \dots, x_n)^{1/n} = 2^{-H(x)}$

• THMITI College DIXIN SECTION:

$\lim_{n \rightarrow \infty} p(x_1, x_2, \dots, x_n) = \lim_{n \rightarrow \infty} 2^{\frac{1}{n} \log p(x_1, x_2, \dots, x_n)} =$

$= \lim_{n \rightarrow \infty} 2^{\frac{1}{n} \sum_{i=1}^n \log p(x_i)} =$ **STRONG LAW OF LARGE NUMBERS**
 $\lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \log p(x_i) = E[\log p(X)]$

$= \lim_{n \rightarrow \infty} 2^{E[\log p(X)]} = 2^{-H(x)}$ **STRONG LAW: ALMOST SURELY**
 $\frac{1}{n} \sum_{i=1}^n \log p(x_i) \xrightarrow{A.S.} E[\log p(X)]$

3.7 ACF AND SOURCE CODING: A DISCRETE MEMORYLESS

SOURCE EMITS A SEQUENCE OF STATISTICALLY INDEPENDENT BINARY DIGITS WITH PROBABILITIES

$p(1) = 0.005$ AND $p(0) = 0.995$. THE DIGITS ARE TAKEN 100 AT TIME AND BINARY CODEWORD IS PROVIDED FOR EVERY SEQUENCE OF 100 DIGITS CONTAINING THREE OR FEWER 1'S.

(a) ASSUMING THAT ALL CODEWORDS ARE THE SAME LENGTH, FIND THE MINIMUM LENGTH REQUIRED TO PROVIDE CODEWORDS FOR ALL SEQUENCES WITH THREE OR FEWER 1'S.

(b) CALCULATE THE PROBABILITY OF OBSERVING A SOURCE SEQUENCE FOR WHICH NO CODEWORD HAS BEEN ASSIGNED.

(c) USE CRESSHEV'S INEQUALITY TO BOUND THE PROBABILITY OF OBSERVING A SOURCE SEQUENCE FOR WHICH NO CODEWORD HAS BEEN ASSIGNED. COMPARE THE BOUND WITH ACTUAL PROBABILITY COMPUTED IN PART (b).

2	SLUČAJ VO	1000	SEKVENCI I	} MAKE EXHAUSTIVE MULTIPLE MULTICOMMUNICATIONS
9	SLUČAJ VO	5000	SEKVENCI II	
18	SLUČAJ VO	10000	SEKVENCI II	
85	SLUČAJ VO	50000	SEKVENCI I	
- VEROVATNOST DA SE GENERALNA SEKVENCA SO > 3 EDICI E:				

$$\frac{5000}{4991} \approx 1$$

$$p = 1 - p(s_3 + s_2 + s_1 + s_0)$$

$s_i \rightarrow$ SEQUENCE WITH i ONES.

$$p(s_0 + s_1 + s_2 + s_3) = \binom{100}{3} p^3 \cdot q^{97} + \binom{100}{2} p^2 \cdot q^{98} + \binom{100}{1} p \cdot q^{99} + \binom{100}{0} p^0 \cdot q^{100}$$

- 000
- 001
- 010
- 011
- 100
- 101
- 110
- 111

$$\binom{3}{2} = \frac{6}{2 \cdot 1} = 3$$

$$\binom{3}{1} = \frac{6}{2! \cdot 1!} = 3$$

$p(0) = 0.995$

$p(1)$

$$H(p) = p \log \frac{1}{p} + q \log \frac{1}{q} = 0.0454$$

$$H(X) = 100 \cdot 0.0454 = 4.545 \text{ BITS}$$

VEROVATNOSTA DA SE GENERALNA SEKVENCA SO PONER OD 3 EDICI E ISTA SO VEROVATNOSTA DA SE GENERALNA SEKVENCA ZA KOZANE E GENERALAN KOPEN ZDOR

$$p = 1 - \binom{100}{3} p^3 \cdot q^{97} - \binom{100}{2} p^2 \cdot q^{98} - \binom{100}{1} p \cdot q^{99} - \binom{100}{0} p^0 \cdot q^{100} = 0.0017$$

(a) MINIMALNA POLŽITA NA KOPITE ŽAOKOVI €

$$u \cdot H(\gamma) = 100 \cdot \left(\gamma \log \gamma + (1-\gamma) \log \frac{1}{(1-\gamma)} \right) \quad \underline{\underline{\gamma = 0.995}}$$

$$u \cdot H(\gamma) = 100 \cdot 0.0454 = 4.5415 \approx 5 \text{ BITS}$$

(c) $\mu = \gamma \cdot 0 + 1 \cdot 1 = 0.005$

$$\sigma^2 = \sum_{x=0}^1 (x-\mu)^2 \gamma(x) = (1-0.005)^2 \cdot 0.005 + (0-0.005)^2 \cdot 0.995$$

$$\sigma^2 = (0.995)^2 \cdot 0.005 + (0.005)^2 \cdot 0.995 = 0.005 = (0.995 + 0.005) \cdot 0.005 = 0.005 \cdot 0.995$$

$$P\{|X-\mu| > \epsilon\} \leq \frac{\sigma^2}{\epsilon^2} \quad \forall \epsilon > 0 \quad \text{CS DRUŽI ŽAOKOVI.}$$

$$P\{|X-\mu| > \epsilon\} \leq \frac{0.005}{\epsilon^2}$$

$$P\{|X-0.005| > \epsilon\} \leq \frac{0.005}{\epsilon^2}$$

OVA JE EKVIVALENTNO NA $P\{X=1\} = 0.005$

NAJGOLEMA VREDNOST ZA " ϵ " € $\epsilon = 0.995$

$$P\{|X-0.005| > \epsilon\} \leq \frac{0.005}{0.995^2} = 0.0051$$

ZNAJI NIJE BOUND € : 0.005.

SE SVI OVA KOPITA = 1

□ SOLUTION 1st EDITION

(a) $\binom{100}{0} + \binom{100}{1} + \binom{100}{2} + \binom{100}{3} = 166751$

$\log(166751) = 17.24 \approx 18$ MINIMUM LEGIT BITI POZITIVNO LEGIT

$u \cdot H(\gamma) = 4.5 \approx 5$ BITI MINIMUM LEGIT (18 IS QUITE LARGER THAN $u \cdot H(\gamma)$)

(b) $P = 1 - \sum_{i=0}^3 \binom{100}{i} \gamma^{100-i} (1-\gamma)^i = 0.00167$

(c) S_n - RANDOM VARIABLE THAT IS SUM OF " n " i.i.d. CHERNIKOV INEQUALITY IS: $P\{|S_n - \gamma n| > \epsilon\} \leq \frac{4 \cdot \sigma^2}{\epsilon^2}$

$$n = 100$$

$$\mu = 1 \cdot 0.005 + 0 \cdot 0.995 = 0.005$$

$$\sigma^2 = |pp.53| = 0.005 \cdot 0.995$$

$S_{100} \geq 4$ IF AND ONLY IF: $|S_{100} - 100 \cdot 0.005| \geq 3.5$

$|S_{100} - 0.5| \geq 3.5 \Rightarrow \boxed{\epsilon = 3.5}$ \rightarrow NAPOZORNOST V KODU JE VEĆE NEGO 4 BITA KODIRANJE 5-BITNA SEKVENCA. STANJE JE 90 NADIMANJE OVA ϵ

$P\{|S_n - n \cdot \mu| \geq \epsilon\} \leq \frac{n \sigma^2}{\epsilon^2}$ OVA DEFIKICIJA JE 90 VGRADUVA VEROVATNOŠĆA DA NEMA ERORIME ZA SEKVENCU ZA KOJE JE ϵ DEFINIRAN KODEN EROR.

SO TO DA OVA JE NIJE VEROVATNOŠĆA ZA SEKVENCU DATA IMA 4 ICI KOJE JE DIO.

$$P\{\text{NEMA KODEN EROR}\} \leq \frac{100 \cdot (0.005 \cdot 0.995)}{3.5^2} = \frac{0.5 \cdot 0.995}{3.5^2} = 0.04061$$

$0.04061 \gg 0.0167$ } LOŠ BOUND DAVA CHEDIŠTEVATA NEKONTRIVOSTI.

3.8 PRODUCTS

LET

$$X = \begin{cases} 1, & \text{WITH PROBABILITY } 1/2 \\ 2, & \text{WITH PROBABILITY } 1/4 \\ 3, & \text{WITH PROBABILITY } 1/4 \end{cases}$$

MMV IMA JOŠ UŠAVA ŠKOLA ČIJA VO MAKE. MultibuytimeCapl. WIKI

LET X_1, X_2, \dots, X_n BE DRAWN I.I.D ACCORDING TO THIS DISTRIBUTION. FIND THE LIMITING BEHAVIOR OF THE PRODUCT

$(X_1 \cdot X_2 \cdot \dots \cdot X_n)^{1/n}$ BY LAW OF LARGE NUMB.

$$\begin{aligned} (X_1 \cdot X_2 \cdot \dots \cdot X_n)^{1/n} &= 2^{\frac{1}{n} \log(X_1 \cdot X_2 \cdot \dots \cdot X_n)} = 2^{\frac{1}{n} \sum_{i=1}^n \log(X_i)} \\ \text{a.s. } \rightarrow 2 & E[\log(X)] = 2^{\frac{1}{4} (1 + \log 3)} = 1.56508 \end{aligned}$$

$$\begin{aligned} H(X) &= \frac{1}{2} \log 2 + \frac{1}{4} \log 4 + \frac{1}{4} \log 4 = \frac{1}{2} + \frac{1}{2} + \frac{1}{2} = \frac{3}{2} \\ E[\log(X)] &= \log(1) \cdot \frac{1}{2} + \log 2 \cdot \frac{1}{4} + \log 3 \cdot \frac{1}{4} = \frac{1}{4} (1 + \log 3) = \frac{1}{4} \log 6 \end{aligned}$$

3.9 AEP. Let X_1, X_2, \dots be INDEPENDENT, IDENTICALLY DISTRIBUTED RANDOM VARIABLES TAKEN ACCORDING TO THE PROBABILITY MASS FUNCTION $q(x)$, $x \in \{1, 2, \dots, m\}$. THUS, $q(x_1, x_2, \dots, x_n) = \prod_{i=1}^n q(x_i)$. WE KNOW THAT $-\frac{1}{n} \log q(x_1, x_2, \dots, x_n) \rightarrow H(X)$ IN PROBABILITY.

LET $z(x_1, x_2, \dots, x_n) = \prod_{i=1}^n z(x_i)$, WHERE z^* IS ANOTHER PROBABILITY MASS FUNCTION ON $\{1, 2, \dots, m\}$

- (a) EVALUATE $\lim_{n \rightarrow \infty} \left(-\frac{1}{n} \log z(x_1, x_2, \dots, x_n) \right)$ WHILE x_1, x_2, \dots, x_n ARE i.i.d. $\sim q(x)$
- (b) NOW EVALUATE THE LIMIT OF THE LOG-LIKELIHOOD RATIO $\frac{1}{n} \log \frac{z(x_1, x_2, \dots, x_n)}{q(x_1, x_2, \dots, x_n)}$ WHEN x_1, x_2, \dots ARE i.i.d. $\sim q(x)$. THUS, THE OPDS FAVORING z^* ARE EXponentially SMALL WHEN z^* IS TRUE.

$$\begin{aligned} \text{(a)} \quad -\frac{1}{n} \log z(x_1, x_2, \dots, x_n) &= -\frac{1}{n} \sum_{i=1}^n \log z(x_i) \rightarrow -E[\log z] \\ &= -\sum_{i=1}^n q(x_i) \log z(x_i) = \underline{H_z(X)} \\ -\frac{1}{n} \log \frac{z(x_1, x_2, \dots, x_n)}{q(x_1, x_2, \dots, x_n)} &= -\frac{1}{n} \left(\sum_{i=1}^n \log z(x_i) - \sum_{i=1}^n \log q(x_i) + \sum_{i=1}^n \log \frac{q(x_i)}{z(x_i)} \right) \\ &= -\frac{1}{n} \sum_{i=1}^n \log \frac{z(x_i)}{q(x_i)} - \frac{1}{n} \sum_{i=1}^n \log q(x_i) \rightarrow -E \left[\log \frac{z(x_i)}{q(x_i)} \right] - E[\log q(x_i)] \\ &= E \left[\log \frac{q(x_i)}{z(x_i)} \right] + H(X) = \underline{D(Y||Z) + H(X)} \end{aligned}$$

$$\begin{aligned} \text{(b)} \quad \frac{1}{n} \log \frac{z(x_1, x_2, \dots, x_n)}{q(x_1, x_2, \dots, x_n)} &\rightarrow -E \left[\log \frac{z(x_1, \dots, x_n)}{q(x_1, \dots, x_n)} \right] = -D(Y||Z) \\ \text{LTO GO ESISTE VO EMISSION A SO MALA VARIANZA} \\ \lim_{n \rightarrow \infty} \frac{1}{n} \log \frac{z(x_1, x_2, \dots, x_n)}{q(x_1, x_2, \dots, x_n)} &= -\lim_{n \rightarrow \infty} \sum_{i=1}^n \log \frac{z(x_i)}{q(x_i)} \text{ w.p. 1} = -\sum q(x) \log \frac{z(x)}{q(x)} \\ &= +\sum q(x) \log \frac{q(x)}{z(x)} = \underline{D(Y||Z) + H(Y)} \end{aligned}$$

(6) $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{z(x_1, x_2, \dots, x_n)}{f(x_1, x_2, \dots, x_n)} = \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \ln \frac{z(x_i)}{f(x_i)}$ w.r.t. = $E \left[\ln \frac{z(x_i)}{f(x_i)} \right]$

090700865

= -D(y||g)

3.10 RANDOM BOX SIZE AN n-DIMENSIONAL RECTANGULAR BOX WITH SIDES x_1, x_2, \dots, x_n IS TO BE CONSTRUCTED. THE VOLUME IS $V_n = \prod x_i$. THE EDGE LENGTH l OF A n-CUBE WITH THE SAME VOLUME AS A RANDOM BOX IS $l = V_n^{1/n}$. LET x_1, x_2, \dots, x_n BE UNIFORM RANDOM VARIABLES OVER UNIT INTERVAL $[0, 1]$. FIND $\lim_{n \rightarrow \infty} V_n^{1/n}$ AND COMPARE TO $(E[V_n])^{1/n}$.

CLEARLY THE EXPECTED EDGE LENGTH DOES NOT CAPTURE THE IDEA OF THE VOLUME OF THE BOX. THE GEOMETRIC MEAN, RATHER THAN THE ARITHMETIC MEAN, CHARACTERIZES THE BEHAVIOR OF THE PRODUCTS.

$$H(x) = - \int f(x) \ln f(x) dx = \int_0^1 f(x) \ln \frac{1}{f(x)} dx = 0$$

$$f(x) = \left[\frac{1}{2}, \frac{1}{2} \right]$$

$$H(x) = \frac{1}{2} \ln 2 + \frac{1}{2} \ln 2 = 1$$

$$f(x) = \left[\frac{1}{4}, \frac{1}{4}, \frac{1}{4}, \frac{1}{4} \right]$$

$$H(x) = \left(\frac{1}{4} \cdot \ln 4 \right) 4 = \left(\frac{1}{2} \right) \cdot 4 = 2$$

(*)
NB!
GLEBO!

$$f(x) = \left[\frac{1}{8}, \frac{1}{8}, \dots, \frac{1}{8} \right]$$

$$H(x) = \left(\frac{1}{8} \cdot 3 \right) 8 = \frac{24}{8} = 3$$

$$\int \frac{dx}{x} = \ln x \quad (\ln x)' = \frac{1}{x}$$

$$\left(\ln \left(\frac{1}{x} \right) \right)' = \left(\frac{1}{x} \right)' = x \cdot x^{-1-1} \cdot (-1) = -x \cdot \frac{1}{x^2} = -\frac{1}{x}$$

• x_i ZAMISLI SI KAKO DA SE PDF ZODI TO SE SO VNEKOST VO PANG $[0, 1]$ (NE MOGA DA ODISI SO) OVAH HISTORICIZ

① $\lim_{n \rightarrow \infty} (x_1 \cdot x_2 \cdot \dots \cdot x_n)^{1/n} = \lim_{n \rightarrow \infty} \frac{1}{n} \ln(x_1 \cdot x_2 \cdot \dots \cdot x_n) = \lim_{n \rightarrow \infty} \left[\frac{1}{n} \sum_{i=1}^n \ln x_i \right]$

A.S. $\rightarrow e^{E[\ln x_i]} = e^{-1} = \frac{1}{e}$

PP. 63

e^x E KONTINUIRNA FUNKCIJA SA ZADOM LIM ODI VO CASO-MENTOT. (CONTINUITY THEOR.)

$$E[V_n]^{1/n} = E[x_1 \cdot x_2 \dots x_n]^{1/n}$$

- 1ST REASON MAKE VO EDITION 1 SOLUTIONS
 SAMO (*) MI E, GREDO T.E. NEKA TOREKA NA PROSJETKA NA HA,

$$E[\ln(x)] = \int_0^1 \ln(x) dx = \int_0^1 \frac{1}{y} dy \quad \left| \begin{array}{l} u = \ln x \\ du = \frac{1}{x} dx \\ v = \int 1 \cdot dx = x \end{array} \right. = x \cdot \ln x - \int x \cdot \frac{dx}{x} = x \ln x - x$$

$$E[\ln(x)] = \int_0^1 (x \ln x - x) dx = -1$$

TEOREMA ZA
 SLEDI VREDNOSTI
 IZ $\int f(x) \cdot \ln(x) dx$
 OD VEZOVAN

$$\lim_{n \rightarrow \infty} V_n^{1/n} = e^{\lim_{n \rightarrow \infty} \frac{1}{n} \ln V_n} = \frac{1}{e} < \frac{1}{2}$$

ALT. ④: $\lim_{n \rightarrow \infty} \ln(V_n)^{1/n} = \lim_{n \rightarrow \infty} \frac{1}{n} \ln(x_1 \cdot x_2 \dots x_n) = \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \ln(x_i)$

A.S. $\xrightarrow{n \rightarrow \infty} E[\ln(x)] = -1$

OD OVA SLEDI DEKA EFektivATA POLZANA NA
 NA OVA TEGO E: e^{-1}

$$E[x_1 \cdot x_2 \dots x_n] = \prod_{i=1}^n E[x_i] = \left(\frac{1}{2}\right)^n$$

$$\int \int \dots \int_{x_1, x_2, \dots, x_n} x_1 \cdot x_2 \dots x_n \cdot f(x_1, \dots, x_n) dx_1 dx_2 \dots dx_n = f(x_1) \cdot f(x_2) \dots f(x_n)$$

$$E[x_i] = \frac{1}{2} = \frac{0+1}{2} \quad \left. \vphantom{E[x_i]} \right\} \text{ARITHMETIC MEAN}$$

$$\textcircled{4} \Rightarrow (x_1 \cdot x_2 \dots x_n)^{1/n} = 1/e < \left(\frac{1}{2}\right)$$

DEFINITION FOR GEOMETRIC MEAN

$$\left(\frac{1}{e}\right)^n < \left(\frac{1}{2}\right)^n \Rightarrow \text{SPEDEN VOLUMEN NA VARNOST GEOMETRISKO VEZO}$$

- SAVA NA VAZE OBUA EFektivATA (SLOVATA)
 POLZITA PODIENA KAKO GEOMETRISKA SREDNA
 S POKAZA OD ONTA. DVA SVA KAKO ARITHMETICNA
 SREDNA.

3.11

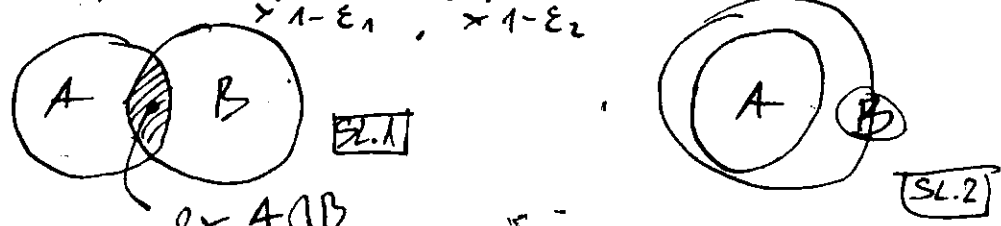
PROOF OF THEOREM 3.3.1.

075426898

THIS PROBLEM SHOWS THAT THE SIZE OF THE SMALLEST "INDISTINGUISHABLE" SET IS ABOUT 2^{4n} . LET $X_1, X_2, \dots, X_n \in$ i.i.d $\sim p(x)$. LET $B_\delta^{(n)} \subset X^n$ SUCH THAT $Pr(B_\delta^{(n)}) \geq 1 - \delta$. FIT $\epsilon < \frac{1}{2}$.

(a) GIVEN ANY TWO SETS A, B SUCH AS $Pr(A_\epsilon^{(n)}) > 1 - \epsilon_1$ AND $Pr(B) > 1 - \epsilon_2$, SHOW THAT $Pr(A \cap B) > 1 - \epsilon_1 - \epsilon_2$. HENCE $Pr(A_\epsilon^{(n)} \cap B_\delta^{(n)}) \geq 1 - \epsilon - \delta$

$$Pr(A+B) = \underbrace{Pr(A)}_{\geq 1-\epsilon_1} + \underbrace{Pr(B)}_{\geq 1-\epsilon_2} - Pr(A \cap B) > 1 - \epsilon_1 + 1 - \epsilon_2 - Pr(A \cap B)$$



$$Pr(A \cap B) = Pr(A) + Pr(B) - Pr(A+B) \geq 1 - \epsilon_1 + 1 - \epsilon_2 - Pr(A+B)$$

$$\geq 1 - \epsilon_1 + 1 - \epsilon_2 - \epsilon \geq 1 - \epsilon_1 - \epsilon_2 \quad \text{PROVED!!!}$$

(b) JUSTIFY THE STEPS IN THE CHAIN OF INEQUALITIES:

$$1 - \epsilon - \delta \leq Pr(A_\epsilon^{(n)} \cap B_\delta^{(n)}) = \sum_{A_\epsilon^{(n)} \cap B_\delta^{(n)}} p(x^n) \leq \sum_{A_\epsilon^{(n)} \cap B_\delta^{(n)}} e^{-\gamma(n-\epsilon)}$$

$$2^{-\gamma(n+\epsilon)} \leq p(x^n) \leq 2^{-\gamma(n-\epsilon)}$$

ASYMPTOTIC EQUIVALENCE THEOREM

$$= |A_\epsilon^{(n)} \cap B_\delta^{(n)}| e^{-\gamma(n-\epsilon)} \leq |A_\epsilon^{(n)}| e^{-\gamma(n-\epsilon)} \leq |B_\delta^{(n)}| e^{-\gamma(n-\epsilon)}$$

↑ PRESENT + POKAZ VEDANOV OD SAMOHO MOZESNO [SL.2]

$$|B_\delta^{(n)}| \geq (1 - \epsilon_1 - \epsilon_2) \cdot e^{\gamma(n-\epsilon)}$$

(c) COMPLETE THE PROOF OF THE THEOREM

$$(1 - \epsilon - \delta) \leq |B_\delta^{(n)}| e^{-\gamma(n-\epsilon)} / |A_\epsilon^{(n)}|$$

$$\ln(1 - \epsilon - \delta) \leq \ln |B_\delta^{(n)}| - \gamma(n-\epsilon)$$

$$\begin{aligned} \ln |B_{\delta}^{(n)}| &\geq \ln(1-\epsilon-\delta) + 4(H-\epsilon) \geq \ln(1-\frac{1}{2}-\delta) + 4(H-\epsilon) \\ &= 4H - 4\epsilon + \ln(\frac{1}{2}-\delta) \geq 4H - \frac{4}{2} + \ln(\frac{1}{2}-\delta) \\ \frac{1}{4} \ln |B_{\delta}^{(n)}| &\geq H + \frac{1}{4} \underbrace{\ln(\frac{1}{2}-\delta)}_{\leq 0} - \frac{1}{2} = H - \delta' \end{aligned}$$

$$\boxed{\delta' = \frac{1}{2} + \frac{1}{4} |\ln(\frac{1}{2}-\delta)|} \quad \delta' > 0$$

• EDITION 1 SOLUTION (WITH COMPLEMENTS)
 - AC KOMPONENT OF A (CLENOKI OD \mathcal{X}^n KOI
 NB PHANATAT NA A.

$$\begin{aligned} P(A^c \cup B^c) &= P(A^c) + P(B^c) - P(A^c \cap B^c) \\ &\leq P(A^c) + P(B^c) \end{aligned}$$

$$P(A) \geq 1 - \epsilon_1 \Rightarrow \begin{matrix} P(A^c) = 1 - P(A) \leq 1 - 1 + \epsilon_1 \\ \boxed{P(A^c) \leq \epsilon_1} \end{matrix} \Rightarrow$$

$$P(B) \geq 1 - \epsilon_2 \Rightarrow \begin{matrix} P(B^c) = 1 - P(B) \leq 1 - 1 + \epsilon_2 = \epsilon_2 \\ \boxed{P(B^c) \leq \epsilon_2} \end{matrix}$$

$$P(A \cap B) = 1 - P(A^c \cup B^c) \geq 1 - P(A^c) - P(B^c) \geq \underline{\underline{1 - \epsilon_1 - \epsilon_2}}$$

$$X_n \in \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$$

$$A \in \{1, 2, 3, 4, 5, 6\}$$

$$A^c \in \{7, 8, 9, 10\}$$

$$B \in \{1, 2, 3, 4, 5, 6, 7\}$$

$$B^c \in \{8, 9, 10\}$$

$$\boxed{A^c \cup B^c = A^c + B^c - A^c \cap B^c} = \{7, 8, 9, 10\}$$

$$A \cap B = \{1, 2, 3, 4, 5, 6\}$$

$$\rightarrow \boxed{P(A \cap B) = 1 - P(A^c \cup B^c)}$$

EMILISUM
 DOUBANO !!!

$$\boxed{N.4. \textcircled{*}} \quad |A_{\epsilon}^{(n)} \cap B_{\delta}^{(n)}| e^{-4(H-\epsilon)} \leq |B_{\delta}^{(n)}| e^{-4(H-\epsilon)}$$

SLEDI OD FAKTOR STO :

$$\boxed{A_{\epsilon}^{(n)} \cap B_{\delta}^{(n)} \subseteq B_{\delta}^{(n)}}$$

3.12

MONOTONIC CONVERGENCE OF EMPIRICAL DISTRIBUTION. LET \hat{p}_n DENOTE EMPIRICAL PROBABILITY MASS FUNCTION CORRESPONDING TO x_1, x_2, \dots, x_n

i.i.d $\sim \gamma(x), x \in \mathcal{X}$, SPECIFICALLY,

$$\hat{p}_n(x) = \frac{1}{n} \sum_{i=1}^n I(x_i = x)$$

IS THE PROPORTION OF TIMES $x_i = x$ IN THE FIRST "n" SAMPLES, WHERE I IS INDICATOR FUNCTION.

(a) SHOW FOR \mathcal{X} FINITE THAT $E[D(\hat{p}_n || \gamma)] \leq E[D(\hat{p}_1 || \gamma)]$

EMPIRICAL SE DOCTRINE TO SAMPLE: $H_n \sim H_{n-1} + \frac{1}{n} \log \frac{1}{n}$

THUS EXPECTED RELATIVE ENTROPY "DISTANCE" FROM THE EMPIRICAL DISTRIBUTION TO THE TRUE DISTRIBUTION DECREASES WITH SAMPLE SIZE.

(Hint: WRITE $\hat{p}_n = \frac{1}{2} \hat{p}_n + \frac{1}{2} \hat{p}_n$ AND USE THE CONVEXITY OF D.)

(b) SHOW FOR AN ARBITRARY DISCRETE \mathcal{X} THAT

$$E[D(\hat{p}_n || \gamma)] \leq E[D(\hat{p}_{n-1} || \gamma)]$$

(Hint: WRITE \hat{p}_n AS THE AVERAGE OF "n" EMPIRICAL MASS FUNCTIONS WITH EACH OF THE "n" SAMPLES DELETED IN TURN.)

$$E[D(\hat{p}_{2n} || \gamma)] \rightarrow \frac{1}{n} D(\hat{p}_{2n} || \gamma) \quad n \rightarrow \text{BROT NA HOLENTA}$$

$$D(\hat{p}_{2n} || \gamma) = \sum_{i=1}^n p_n \left(\log \frac{p_{2n}(i)}{p(i)} \right)$$

$$\hat{p}_{2n} = [p_n, p_n]$$

• ZA KONVEXNI. FURTHER USE JENSEN'S INEQUALITY.

$$E[D(\hat{p}_{2n} || \gamma)] \geq D(E[\hat{p}_{2n} || \gamma])$$

$$D(\hat{p} || \gamma) = \hat{p}_1 \cdot \log \frac{\hat{p}_1}{\gamma_1} + \hat{p}_2 \cdot \log \frac{\hat{p}_2}{\gamma_2} = E_{p_n} \left[\log \frac{\hat{p}_n}{\gamma} \right]$$

$$D(\hat{p}_4 || \gamma) = \frac{1}{4} \sum_{i=1}^4 I(x_i=1) \cdot \text{Ld} \left(\frac{\frac{1}{4} \sum_{i=1}^4 I(x_i=1)}{\gamma_1} \right) +$$

$$+ \frac{1}{4} \sum_{i=1}^4 I(x_i=2) \text{Ld} \left(\frac{\frac{1}{4} \sum_{i=1}^4 I(x_i=2)}{\gamma_2} \right)$$

$$D(\hat{\gamma}_{24} || q) = \frac{1}{24} \sum_{i=1}^{24} I(x_i=1) \cdot \text{Ld} \left(\frac{\frac{1}{24} \sum_{i=1}^{24} I(x_i=1)}{\gamma_1} \right) +$$

$$+ \frac{1}{24} \sum_{i=1}^{24} I(x_i=2) \text{Ld} \left(\frac{\frac{1}{24} \sum_{i=1}^{24} I(x_i=2)}{\gamma_2} \right) =$$

$$= \frac{1}{2} \left[\frac{1}{4} \sum_{i=1}^4 I(x_i=1) \text{Ld} \frac{\frac{1}{24} \sum_{i=1}^{24} I(x_i=1)}{\gamma_1} + \right.$$

$$\left. + \frac{1}{4} \sum_{i=1}^4 I(x_i=2) \text{Ld} \frac{\frac{1}{24} \sum_{i=1}^{24} I(x_i=2)}{\gamma_2} \right] +$$

$$+ \frac{1}{24} \sum_{i=11+1}^{24} I(x_i=1) \text{Ld} \frac{\frac{1}{24} \sum_{i=1}^{24} I(x_i=1)}{\gamma_1} +$$

$$+ \frac{1}{24} \sum_{i=17+1}^{24} I(x_i=2) \text{Ld} \frac{\frac{1}{24} \sum_{i=1}^{24} I(x_i=2)}{\gamma_2}$$

$$\hat{p}_{24} = \frac{1}{24} \sum_{i=1}^{24} I(x_i=x) = \frac{1}{2} \frac{1}{4} \sum_{i=1}^4 I(x_i=x) + \frac{1}{24} \sum_{i=11+1}^{24} I(x_i=x)$$

$$\hat{\gamma}_{24} = \frac{1}{2} \gamma_4 + \frac{1}{2} \gamma_4' \quad \hat{\gamma}_{24} \rightarrow \in [\gamma_4] \gamma_4 \quad \textcircled{*}$$

$$D(p_{24} || \gamma) = \underline{p_{24}(1)} \cdot \text{Ld} \frac{p_{24}(1)}{\gamma(1)} + p_{24}(0) \cdot \text{Ld} \frac{p_{24}(0)}{\gamma(0)}$$

$$= \left(\frac{1}{2} \gamma_4(1) + \frac{1}{2} \gamma_4'(1) \right) \text{Ld} \left[\frac{\frac{1}{2} \gamma_4(1) + \frac{1}{2} \gamma_4'(1)}{\gamma(1)} \right] + p_{24}(0) \text{Ld} \frac{p_{24}(0)}{\gamma(0)}$$

$$\begin{aligned}
D(\gamma_{2n} || \gamma) &= \frac{1}{2} (\gamma_1(n) + \gamma_1'(n)) \ln \frac{\gamma_1(n) + \gamma_1'(n)}{2\gamma(n)} + \frac{1}{2} (\gamma_1(0) + \gamma_1'(0)) \ln \frac{\gamma_1(0) + \gamma_1'(0)}{2\gamma(0)} \\
&= \frac{1}{2} (\gamma_{1n} + \gamma_{1n}') \ln \frac{\gamma_{1n} + \gamma_{1n}'}{2\gamma_1} + \frac{1}{2} (1 - \gamma_{1n} + 1 - \gamma_{1n}') \cdot \ln \left[\frac{2 - \gamma_{1n} - \gamma_{1n}'}{2(1 - \gamma_1)} \right] \\
&= \frac{1}{2} (\gamma_1 + \gamma_1') \ln \frac{\gamma_1 + \gamma_1'}{2\gamma} + \frac{1}{2} (2 - \gamma_1 - \gamma_1') \ln \frac{2 - \gamma_1 - \gamma_1'}{2(1 - \gamma)} = \\
&= \frac{1}{2} (\gamma_1 + \gamma_1') \ln \frac{\gamma_1 + \gamma_1'}{2\gamma} - \frac{1}{2} (\gamma_1 + \gamma_1') \ln \frac{2 - \gamma_1 - \gamma_1'}{2(1 - \gamma)} + \ln \frac{2 - \gamma_1 - \gamma_1'}{2(1 - \gamma)} \\
&= \frac{1}{2} (\gamma_1 + \gamma_1') \ln \frac{(\gamma_1 + \gamma_1') 2(1 - \gamma)}{2\gamma(2 - \gamma_1 - \gamma_1')} + \ln \frac{2 - \gamma_1 - \gamma_1'}{2(1 - \gamma)}
\end{aligned}$$

$$D(\gamma_1 || \gamma) = \gamma_1(n) \ln \frac{\gamma_1(n)}{\gamma} + \gamma_1(0) \ln \frac{\gamma_1(0)}{1 - \gamma}$$

$$D(\gamma_1 || \gamma) = \gamma_1 \ln \frac{\gamma_1}{\gamma} + (1 - \gamma_1) \ln \frac{1 - \gamma_1}{1 - \gamma}$$

for $n \rightarrow \infty$ $\boxed{\gamma_1 = \gamma_1'}$

$$D(\gamma_{2n} || \gamma) = \frac{1}{2} (\gamma_1 + \gamma_1') \ln \frac{(\gamma_1 + \gamma_1')(1 - \gamma)}{\gamma(2 - \gamma_1 - \gamma_1')} + \ln \frac{2 - \gamma_1 - \gamma_1'}{2(1 - \gamma)}$$

IF: $\gamma_1 = \gamma_1'$ $D(\gamma_{2n} || \gamma) = \gamma_1 \ln \frac{2\gamma_1(1 - \gamma)}{2\gamma(1 - \gamma_1)} + \ln \frac{1 - \gamma_1}{1 - \gamma}$

$$D(\gamma_{2n} || \gamma) = \gamma_1 \ln \frac{\gamma_1}{\gamma} + (1 - \gamma_1) \ln \frac{1 - \gamma_1}{1 - \gamma} = D(\gamma_1 || \gamma)$$

$$D(\hat{P}_{2n} || \gamma) = \frac{1}{2n} [\hat{P}_{1n}(n) + \hat{P}_{1n}'(n) + \dots + \hat{P}_{1n}(n)] \ln \frac{1}{2n} [\hat{P}_{1n}(n) + \dots + \hat{P}_{1n}(n)]$$

$$+ \frac{1}{2n} [\hat{P}_{1n}(0) + \hat{P}_{1n}'(0) + \dots + \hat{P}_{1n}(0)] \ln \frac{1}{2n} [\hat{P}_{1n}(0) + \dots + \hat{P}_{1n}(0)]$$

$$\rightarrow \in [P_{1n}(n)] \ln \frac{\in [P_{1n}(n)]}{\gamma(n)} + \dots + \in [P_{1n}(0)] \ln \frac{\in [P_{1n}(0)]}{\gamma(0)}$$

(6) ARBITRARY DISCRETE "X" $x \in \{x_1, x_2, \dots, x_n\}$
 $E[D(\hat{p}_n \| q)] \leq E[D(p_{n-1} \| q)]$

$$D(\hat{p}_n \| q) = \hat{p}_n(x_1) \log \frac{p_n(x_1)}{q(x_1)} + \hat{p}_n(x_2) \log \frac{p_n(x_2)}{q(x_2)} + \dots$$

$$+ \hat{p}_n(x_n) \log \frac{p_n(x_n)}{q(x_n)} = E_{p_n} \left[\log \left(\frac{\hat{p}_n(x)}{q(x)} \right) \right]$$

- RETURN TO P.67 *

$$D(\hat{p}_{2n} \| q) = \hat{p}_{2n}(1) \log \frac{p_{2n}(1)}{q(1)} + \hat{p}_{2n}(0) \log \frac{p_{2n}(0)}{q(0)}$$

$$\hat{p}_{2n}(1) = \frac{1}{2} p_n(1) + \frac{1}{2} p_n^{-}(1) \xrightarrow{\text{A.S.}} E[p_n(1)] = q(1)$$

$$p_n(1) = \frac{1}{n} \sum_{i=1}^n I(X_i=1) \left. \begin{array}{l} \text{KAKO KOSTE "n" TAKA} \\ \text{E VOJIMO DA KOSTE 1} \\ \text{p_n(1) POSTO MOGA} \\ \text{POVEĆE EDICIJA.} \end{array} \right\}$$

SOGLAZNO ZANODOT ZA GOLEMI BROJEVI PROBE
 NEKOE GOLEMO "n" BROJOT NA EDICIJA
 PORAZEN SO DOZETATA NA PRIMEROKOT ĆE
 BIDE EDIKVA NA VEROVATNOSTA NA POTAVNA
 NGE NA EDICIJA.

$$D(\hat{p}_n \| q) \xrightarrow[n \uparrow]{} 0 \Rightarrow \underline{D(\hat{p}_{2n} \| q) < D(\hat{p}_n \| q)}$$

OLASIA

$$D(\lambda p_1 + (1-\lambda)p_2 \| \lambda q_1 + (1-\lambda)q_2) \leq \lambda D(p_1 \| q_1) + (1-\lambda)D(p_2 \| q_2)$$

CONVEXITY OF RELATIVE ENTROPY

solu-2-229 SOLUTION

$$p_n^-(x) = \frac{1}{n} \sum_{i=1}^n I(X_{n+i}=x) \text{ FOR } n, x \in X$$

- N.67 * $\Rightarrow \hat{p}_{2n} = \left(\hat{p}_n + \hat{p}_n^- \right) \cdot \frac{1}{2}$

BT CONVEXITY OF RELATIVE ENTROPY

$$D(\hat{p}_{2n} \| q) = D\left(\frac{1}{2} \hat{p}_n + \frac{1}{2} \hat{p}_n^- \| \frac{1}{2} q + \frac{1}{2} q\right) \leq \frac{1}{2} D(\hat{p}_n \| q) + \frac{1}{2} D(\hat{p}_n^- \| q) \in [\dots]$$

$$E[D(\hat{y}_n || \gamma)] = \frac{1}{2} [E[D(\hat{y}_n^{(1)} || \gamma)]] + \frac{1}{2} E[D(\hat{y}_n^{(2)} || \gamma)] = E[D(\hat{y}_n^{(1)} || \gamma)]$$

$$E[D(\hat{y}_n || \gamma)] = E[D(\hat{y}_n^{(1)} || \gamma)]$$

$\hat{y}_n^{(1)}, \hat{y}_n^{(2)}, \hat{y}_n^{(i)}$ ARE RANDOM VARIABLES TAKING VALUES IN THE SPACE OF PROBABILISTIC DISTRIBUTIONS OF \mathcal{X}

(6) FOR $1 \leq i \leq n$, LET $\gamma_n^{(i)}$ DENOTE THE EMPIRICAL DISTRIBUTION OF $(n-1)$ RANDOM VARIABLES X_j $1 \leq j \leq n$ $j \neq i$, I.E.

$$\hat{\gamma}_n^{(i)}(x) = \frac{1}{n-1} \sum_{j=1, j \neq i}^n 1(X_j = x) \text{ FOR } x \in \mathcal{X}$$

$$\hat{\gamma}_n = \frac{1}{n} \sum_{i=1}^n \hat{\gamma}_n^{(i)}$$

BY USING CONVEXITY OF RELATIVE ENTROPY FUNCTION

$$D(\hat{\gamma}_n || \gamma) \leq \frac{1}{n} \sum_{i=1}^n D(\hat{\gamma}_n^{(i)} || \gamma)$$

क्या लोकार्थ
कालो बा लोस
सोत 24²

$$E[D(\hat{\gamma}_n || \gamma)] \leq \frac{1}{n} \sum_{i=1}^n E[D(\hat{\gamma}_n^{(i)} || \gamma)] = E[D(\hat{\gamma}_n^{(1)} || \gamma)]$$

$$E[D(\hat{\gamma}_n^{(i)} || \gamma)] = E[D(\hat{\gamma}_n^{(1)} || \gamma)] \text{ FOR } \forall i$$

RECALL: LOG-SUM INEQUALITY

$$\sum_{i=1}^n a_i \log \frac{a_i}{b_i} \geq \left(\sum_{i=1}^n a_i \right) \log \frac{\sum_{i=1}^n a_i}{\sum_{i=1}^n b_i}$$

THEOREM 2.7.2 (CONVEXITY OF RELATIVE ENTROPY)

$$D(\lambda \gamma_1 + (1-\lambda) \gamma_2 || \lambda \gamma_1' + (1-\lambda) \gamma_2') \leq \lambda D(\gamma_1 || \gamma_1') + (1-\lambda) D(\gamma_2 || \gamma_2')$$

$$(\lambda \gamma_1 + (1-\lambda) \gamma_2) \log \frac{\lambda \gamma_1 + (1-\lambda) \gamma_2}{\lambda \gamma_1' + (1-\lambda) \gamma_2'} \leq \lambda \gamma_1 \log \frac{\gamma_1}{\gamma_1'} + (1-\lambda) \gamma_2 \log \frac{\gamma_2}{\gamma_2'}$$

$(\gamma_1 || \gamma_1')$ & $(\gamma_2 || \gamma_2')$ ARE TWO PAIRS OF PROBABILISTIC MASS FUNCTIONS.

3.13 CALCULATION OF TYPICAL SET: TO CLARIFY THE

NOTION OF TYPICAL SET $A_\epsilon^{(n)}$ AND THE SMALLEST SET OF HIGH PROBABILITY 2^H WE WILL CALCULATE THE SET FOR A SIMPLE EXAMPLE. CONSIDER A SEQUENCE OF I.I.D BINARY RANDOM VARIABLES, X_1, X_2, \dots, X_n WHERE THE PROBABILITY THAT $X_i = 1$ IS 0.6 ($p=0.6, q=0.4$)

- (a) CALCULATE $H(X)$
- (b) WITH $n = 25$ AND $\epsilon = 0.1$, WHICH SEQUENCES FALL IN THE TYPICAL SET $A_\epsilon^{(n)}$? WHAT IS THE PROBABILITY OF THE TYPICAL SET? HOW MANY ELEMENTS ARE THERE IN THE TYPICAL SET? (THIS INVOLVES COMPUTATION OF THE TABLE OF PROBABILITIES FOR SEQUENCES WITH k 1'S, $0 \leq k \leq 25$, AND FINDING THOSE SEQUENCES THAT ARE IN THE TYPICAL SET.)
- (c) HOW MANY ELEMENTS ARE THERE IN THE SMALLEST SET THAT HAS PROBABILITY 0.9 ?
- (d) HOW MANY ELEMENTS ARE THERE IN THE INTERSECTION OF THE SETS IN (b) AND (c)? WHAT IS THE PROBABILITY OF THIS INTERSECTION.

$$2^{-n(H+\epsilon)} \leq P(X^n) \leq 2^{-n(H-\epsilon)}$$

$$Pr\{A_\epsilon^{(n)}\} \geq 1 - \epsilon \quad |A_\epsilon^{(n)}| \leq 2^{n(H+\epsilon)} \quad |A_\epsilon^{(n)}| \geq (1-\epsilon)2^{n(H-\epsilon)}$$

$$(H-\epsilon) \leq -\frac{1}{n} \log P(X^n) \leq (H+\epsilon) \quad \text{ALL SEQUENCES FOR WHICH THIS IS TRUE, } \epsilon \text{ TYPICAL SET}$$

(a) $H(X) = p \log \frac{1}{p} + q \log \frac{1}{q} = 0.6 \log \frac{1}{0.6} + 0.4 \log \frac{1}{0.4} = \underline{0.970951}$

(b) $n = 25 \quad 1 \leq k \leq 25$
 TOTAL NUMBER OF SEQUENCES WITH GIVEN "k" IS
 $\binom{25}{k}$ e.g. $\binom{25}{1} = \frac{25!}{24!} = 25$ THAT ARE 25 SEQUENCES WITH "1" ONE

PROBABILITY OF THOSE SEQUENCES IS
 $\binom{25}{k} \cdot p^k \cdot q^{25-k}$ e.g. $25 \cdot (0.6)^1 \cdot (0.4)^{24} = 4.2212 \cdot 10^{-9}$

$$-\frac{1}{n} \log ((0.6)^1 \cdot (0.4)^{24}) = \underline{1.29853}$$

• $A_\epsilon^{(n)}$ VO OVA MROZESTVO MISLAM DENA IADANAT SITE SEKVENCII DO $25 \cdot 0.6 = 15$ EDINICI NO

$$0.970951 - 0.1 \leq -\frac{1}{25} \log ((0.6)^{15} \cdot (0.4)^{10}) = 0.970951 \leq 0.970951 + 0.1$$

$$H - \epsilon \leq -\frac{1}{4} \ln \gamma(x^*) \leq H + \epsilon$$

$$-\epsilon \leq -\frac{1}{4} \ln \gamma(x^*) - H \leq \epsilon$$

$$\left| -\frac{1}{4} \ln \gamma(x^*) - H \right| \leq \epsilon$$

SKAŽTE NA FOKA NA
USCROT ZA TIKRO
MROSTVO

$$\Pr\{A_\epsilon^{(n)}\} \geq 1 - \epsilon \geq 0.9$$

OVA SO ISPOLUVIATA
SITE SEQUENCI SO:
11, 12, ..., 19 EKOICI

PROBABILITY OF TYPICAL SET IS:

$$P_A = \sum_{i=11}^{19} \binom{25}{i} p^i 2^{25-i} = 0.936246$$

ENAGGLAVATA HOVE-
KVALI! ISTA ZA
PLAVI I VO UNIVE-
RSITE OF TORONTO.

$$P_A = 0.936246 \geq 0.9 = 1 - \epsilon$$

THE NUMBER OF ELEMENTS IN TYPICAL SET IS:

$$N_A = \sum_{i=11}^{19} \binom{25}{i} = 26,366,510 \text{ ELEMENTS}$$

$$|A_\epsilon^{(n)}| \leq 2^{4(H+\epsilon)} = 2^{25(1.070951)} = 114,727,700$$

$$26,366,510 \leq 114,727,700$$

$$8 + 7 + 6 + 4 + 2 = 15 + 10 + 5 = 30$$

MINIMAL SET [12, 19] $P_B = 0.83284 \approx 0.1$

$$N_B = \sum_{i=12}^{19} \binom{25}{i} = 21,909,110$$

$$P_B = 0.9$$

SUM ZEMAL
DA BIDE " = " +
TRIDA

NUMBER OF ELEMENTS IN THE INTERSECTION

$$|A_\epsilon^{(n)} \cap B_\epsilon^{(n)}| = |B_\epsilon^{(n)}| = 21,909,110 \text{ ELEMENTS}$$

$$\frac{1}{4} \ln |B_\epsilon^{(n)}| > H - \delta' \quad (H = 0.970951)$$

$$\frac{1}{25} \ln(21,909,110) = 0.97540$$

$$\delta' = \frac{1}{2} + \frac{1}{4} \left| \ln\left(\frac{1}{2} - \delta\right) \right| = \frac{1}{2} + \frac{1}{25} \left| \ln(0.4) \right| = 0.53665$$

$$H - \delta' = 0.970951 - 0.53665 = 0.43430$$

$$25+24+23+\dots+2 = \frac{4 \cdot (4+1)}{2} - 1$$

PASPALEK Pi

$$S = 1 + 2 + 3 + \dots + n$$

$$S = n + (n-1) + (n-2) + \dots + 1$$

$$2S = n+1 + n+1 + n+1 + \dots + (n+1)$$

$$2S = (n+1) \cdot n$$

$$S = \frac{(n+1) \cdot n}{2}$$

• VREZ OROV NA SIMULACIJE VO MANJE OD 10 DENA MINIMIZEN SET E SO SOČINUVANAT SEKVENCITE ISTO KAKO KAKO:

12, 13, 14, ..., 20 EDICI

SUMATA NA VEROVATNOSTITE NA OVE SEKVENCII E ≥ 0.9 , VODI TO E MNOŽESTVO SO DVAZ NAOT NA ELEMENTI ISTO SO IŠTOLUVA OVOJ USLOV.

$$\sum_{\lambda=12}^{20} \binom{25}{\lambda} p^{\lambda} q^{25-\lambda} = 0.9127 >$$

$$\sum_{\lambda=12}^{20} \binom{25}{\lambda} = 21,962,240$$

$$|A_{\lambda}^{(n)} \cap B_{\lambda}^{(n)}| = \sum_{\lambda=12}^{19} \binom{25}{\lambda} = 21,909,110$$

NAVISTIA PRES EKOT E OGRANEN

- V_0 (#) SUCUČNO SUM ZA DVA IŠTATA VREDOST

A NE SUM DVAZ MNOŽESTVO DA KVA NA MINIMIZEN PRO NA ELEMENTI.

WRIGHT COLLEGE SOLUTION

(c) $N_0 p^{12} q^{13} \leq 0.9 - \sum_{j=13}^{25} \binom{25}{j} p^j q^{25-j}$ } ZEMAJ OVOJ SEKVENCI SO 12 EDICI ZA DA VUKATA VEROVATNOST BIDE 0.9

$$N \leq \frac{0.9 - \sum_{j=13}^{25} \binom{25}{j} p^j q^{25-j}}{p^{12} q^{13}}$$

$$N = 3.68067 \cdot 10^6$$

- ZNAČI VUKAJOT DVAZ NA ELEMENTI VO MINIMACIJO MNOŽESTVO E:

$$P = N \cdot p^{12} q^{13} + \sum_{j=13}^{25} \binom{25}{j} p^j q^{25-j}$$

$$|B_{\delta}^{(n)}| = N + \sum_{j=13}^{25} \binom{25}{j} = 2,045,710$$

$$|B_8^{(n)}| = |A_E^{(n)}| = 2^{n \cdot H} = 2^{25 \cdot 0.970951} = \underline{\underline{2.0283 \cdot 10^7}}$$

- And se slovedl mltava presvedca (BARANA NA MALE) I PREKETAJA OD WRIGHT COLLEGE, SE GLEDA DEKA (I) VS (II) N.72):

$$(I) = 21,90911 \cdot 10^6$$

$$(II) = 20,4579 \cdot 10^6$$

$$2^{n \cdot H} = 20,283 \cdot 10^6$$

$$2^{n(H+E)} = 114,7377 \cdot 10^6$$

$$n \cdot H = 25 \cdot 0,97 =$$

DEFINITIVNO (II) E POKRO VEŠČIE KOE ŽAS NE ŠO VODIV VO MALE ŽOŠTO FAZGLEDAVAN OZI EKUI NA SEKVENCI, A NE DEC OD SEKVENCI. SE PABOT ZA OPTIMIZACIEN PROBLEM.

$$(b) |A_E^{(n)} \cap B_8^{(n)}| = |A_{0..n}^{(n)} \cap B_{0..n}^{(n)}| =$$

$$A_E^{(n)} \quad k \in \{11, 12, \dots, 19\}$$

$$B_8^{(n)} \quad k \in \{\text{part of } 12, 13, \dots, 25\}$$

$$|A_{0..n}^{(n)} \cap B_{0..n}^{(n)}| \Rightarrow k \in \{\text{part of } 12, 13, \dots, 19\}$$

$$|A_{0..n}^{(n)} \cap B_{0..n}^{(n)}| = N + \sum_{i=12}^{19} \binom{25}{i} = 3,68067 \cdot 10^6 + 16,702210 = 2.03895 \cdot 10^7$$

$$P_{AND} = N \cdot q^{12} q^{13} + \sum_{i=12}^{19} \binom{25}{i} q^i q^{25-i} = \underline{\underline{0.87064}}$$

• SUDVITAVA ŽOŠTO $B_{0..n}^{(n)}$: $k \in \{\text{part of } 12, 13, \dots, 25\}$

E ŽAŠOA ŠFO $q > 0$ FA ŽAŠOA ŽATODROVK VO OD $k=25$ FA ŽUMIJA WAPOKU. TAKA IODLEGU (ZA IONAL ISAP NA ŽEROV) SE ŠTĚA VO VIVIA VEROTAČNOST 0.9. (MMV)

CHAPTER 4 ENJOY THE IDEAS OF ERGODIC PROCESSES
 AEP ESTABLISH THAT $n \cdot H(X)$ BITS SUFFICE ON AVERAGE TO DESCRIBE " n INDEPENDENT AND IDENTICALLY DISTRIBUTED RANDOM VARIABLES.

- If X_1, X_2, \dots, X_n are DEPENDENT we will show that the entropy $H(X_1, X_2, \dots, X_n)$ grows (asymptotically) linearly with n at a rate $H(X)$ which we will call the ENTROPY RATE of the process.

4.1 MARKOV CHAINS

A STOCHASTIC PROCESS $\{X_i\}$ is an INDEXED SEQUENCE OF RANDOM VARIABLES. In general there can be arbitrary dependence among the random variables. Process is characterized by the joint probability mass functions:

$$P\{(X_1, X_2, \dots, X_n) = (x_1, x_2, \dots, x_n)\} = p(x_1, x_2, \dots, x_n),$$

$$(x_1, x_2, \dots, x_n) \in \mathcal{X}^n \text{ for } n = 1, 2, \dots$$

DEFINITION A stochastic process is said to be STATIONARY if the joint distribution of any subset of the sequence of random variables is invariant with respect to shifts in TIME INDEX, that is:

$$P\{X_1 = x_1, X_2 = x_2, \dots, X_n = x_n\} = P\{X_{1+k} = x_1, X_{2+k} = x_2, \dots, X_{n+k} = x_n\}$$

for every n and every shift k , and for all $x_1, x_2, \dots, x_n \in \mathcal{X}$

DEFINITION A discrete stochastic process X_1, X_2, \dots is said to be a MARKOV CHAIN or MARKOV PROCESS if for $n = 1, 2, \dots$

$$P\{X_{n+1} = x_{n+1} | X_n = x_n, X_{n-1} = x_{n-1}, \dots, X_1 = x_1\} = P\{X_{n+1} = x_{n+1} | X_n = x_n\}$$

for all $x_1, x_2, \dots, x_n, x_{n+1} \in \mathcal{X}$

$$p(x_1, x_2, \dots, x_n) = p(x_1) p(x_2 | x_1) \cdot p(x_3 | x_2) \cdot \dots \cdot p(x_n | x_{n-1})$$

$$p(x_1, x_2, \dots, x_n) = p(x_1) \cdot p(x_2, x_3, \dots, x_n | x_1) = p(x_1) \cdot p(x_2 | x_1) \cdot p(x_3, \dots, x_n | x_1, x_2)$$

$$= p(x_1) \cdot p(x_2 | x_1) \cdot p(x_3 | x_1, x_2) \cdot p(x_4, \dots, x_n | x_1, x_2, x_3) = p(x_1) \cdot p(x_2 | x_1)$$

$$\cdot p(x_3 | x_2) \cdot p(x_4 | x_1, x_2, x_3) \cdot p(x_5, \dots, x_n | x_1, x_2, x_3, x_4) \dots$$

FOR TIME INVARIANT MARKOV CHAIN:

$P_r(x_{t+1}=6 | x_t=9) = P_r\{x_t=6 | x_1=9\}$ for all $9, 6 \in X$

• If $\{x_i\}$ is a Markov Chain, x_t is called the state at time t . A time invariant Markov chain is characterized by its initial state and probability transition matrix

$P = [P_{ij}]$, $i, j \in \{1, 2, \dots, m\}$ $P_{ij} = P_r\{x_{t+1}=j | x_t=i\}$

If it is possible to go with positive probability from any state of Markov chain to any other state in the finite number of steps, the Markov chain is said to be **irreducible**.

• If the probability mass function of random variable at time t is $f(x_t)$ the probability mass function at time $(t+1)$ is:

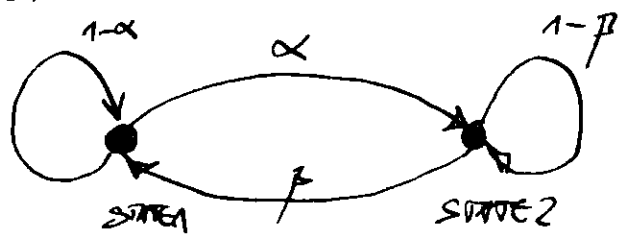
$f(x_{t+1}) = \sum_{x_t} f(x_t) \cdot P_{x_t, x_{t+1}}$

A distribution of states such that distribution at time $t+1$ is the same as the distribution at t is called **stationary distribution**.

EXAMPLE 4.1.1

CONSIDER A TWO STATE MARKOV CHAIN WITH PROBABILITY TRANSITION MATRIX

✓ $P = \begin{bmatrix} 1-\alpha & \alpha \\ \beta & 1-\beta \end{bmatrix}$



μ -VECTOR REPRESENTING STATIONARY PROBABILITIES OF STATES 1 & 2, RESPECTIVELY.

$\mu \cdot P = \mu$

FIND STATIONARY PROBABILITIES BY SOLVING OF THIS EQUATION.

$[\mu_1 \mu_2] \begin{bmatrix} 1-\alpha & \alpha \\ \beta & 1-\beta \end{bmatrix} = [\mu_1(1-\alpha) + \mu_2 \beta, \alpha \mu_1 + (1-\beta)\mu_2]$
 $= [\mu_1, \mu_2]$

$\mu_1(1-\alpha) + \mu_2 \beta = \mu_1$

$-\alpha \mu_1 + \mu_2 \beta = 0$

$\alpha \mu_1 = \mu_2 \beta$

$2\mu_1 + (1-\beta)\mu_2 = \mu_2$

$2\mu_1 - \beta \mu_2 = 0$

$2\mu_1 = \beta \mu_2$

$\mu_1 + \mu_2 = 1$

$\mu_1 = 1 - \mu_2$

$2(1-\mu_2) = \beta \mu_2$

$$\alpha - \alpha p_2 = p p_2 \quad (\alpha + p) p_2 = \alpha$$

$$p_2 = \frac{\alpha}{\alpha + p}$$

$$p_1 = \frac{p}{\alpha} \cdot p_2 = \frac{p}{\alpha + p}$$

- IF PROCESS IS STATIONARY THE ENTROPY OF THE STATE X_n AT TIME n IS:

$$H(X_n) = H\left(\frac{p}{\alpha + p}, \frac{\alpha}{\alpha + p}\right) = \frac{p}{\alpha + p} \log \frac{\alpha + p}{p} + \frac{\alpha}{\alpha + p} \log \frac{\alpha + p}{\alpha}$$

$$H(Y|X) = \sum_{x \in X} p(x) H(Y|X=x)$$

$$H(Y|X) = \sum_{x \in X} p(x, Y) \log \frac{1}{p(Y|x)}$$

4.2 ENTROPY RATE

IF WE HAVE A SEQUENCE OF n RANDOM VARIABLES, A NATURAL QUESTION TO ASK IS: HOW DOES THE ENTROPY OF THE SEQUENCE GROW WITH n ? WE DEFINE THE ENTROPY RATE AS THIS RATE OF GROWTH AS FOLLOWS.

DEFINITION: THE ENTROPY OF STATIONARY PROCESS $\{X_n\}$ IS DEFINED BY:

$$H(X) = \lim_{n \rightarrow \infty} \frac{1}{n} H(X_1, X_2, \dots, X_n)$$

WHEN THE LIMIT EXISTS.

EXAMPLES:

1.) TELETYPEWRITER: M - EQUALLY LIKELY OUTPUT LETTERS
 M^n - SEQUENCES OF LENGTH n , ALL EQUALLY LIKELY.

$$H(X_1, X_2, \dots, X_n) = \sum_{i=1}^n \frac{1}{M^n} \cdot \log M^n = n \cdot \frac{1}{M^n} \cdot \log M^n$$

$$H(X_1, X_2, \dots, X_n) = \log M^n \quad \text{THE ENTROPY RATE IS:}$$

$$H(X) = \lim_{n \rightarrow \infty} \frac{1}{n} H(X_1, X_2, \dots, X_n) = \frac{1}{n} \log M^n = \log M$$

2.) X_1, X_2, \dots, X_n ARE I.I.D RANDOM VARIABLES. THEN

$$H(X) = \lim_{n \rightarrow \infty} \frac{H(X_1, X_2, \dots, X_n)}{n} = \lim_{n \rightarrow \infty} \frac{n H(X_1)}{n} = H(X_1)$$

3. SEQUENCE OF INDEPENDENT BUT NOT IDENTICALLY DISTRIBUTED RANDOM VARIABLES

$$H(X_1, X_2, \dots, X_n) = \sum_{i=1}^n H(X_i)$$

E.G. SEQUENCE OF PISOX ACTION SUCH AS:

$$\lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n H(x_i) \text{ DOES NOT EXIST}$$

$p_i = P(x_i=1)$ IS CONSTANT BUT FUNCTION OF $i =$

$$p_i = \begin{cases} 0.9 & \text{IF } 2k \leq i \leq 2k+1 \\ 0 & \text{IF } 2k+1 \leq i \leq 2k+2 \end{cases}$$

FOR $k=0,1,2,\dots$

$H(x_i)$ WILL OSCILLATE BETWEEN 0 & 1 , AND WILL NOT HAVE THE LIMIT.

QUANTITY OF ENTROPY RATE: (Alternative Definition)

$$H'(X) = \lim_{n \rightarrow \infty} H(x_n | x_{n-1}, x_{n-2}, \dots, x_1)$$

$H(X) \sim$ PER SYMBOL ENTROPY OF n RANDOM VARIABLES.

$H'(X) \sim$ CONDITIONAL ENTROPY OF LAST RANDOM VARIABLE GIVEN THE PAST.

Theorem 4.2.1 FOR STATIONARY STOCHASTIC PROCESS THE LIMITS $H(X)$ AND $H'(X)$ EXIST AND ARE EQUAL:

$$H(X) = H'(X)$$

$$H(x_1, x_2, \dots, x_n) = H(x_1) + H(x_2 | x_1) + H(x_3 | x_2, x_1) + \dots$$

$$H(x_1, x_2, \dots, x_n) = \sum_{i=1}^n H(x_i | x_{i-1}, \dots, x_1)$$

$$\lim_{n \rightarrow \infty} \frac{1}{n} H(x_1, x_2, \dots, x_n) = \sum_{i=1}^n \lim_{n \rightarrow \infty} \frac{1}{n} H(x_i | x_{i-1}, \dots, x_1)$$

IF EXISTS

THIS LIMIT EXISTS ALSO

Theorem 4.2.2 FOR STATIONARY STOCHASTIC PROCESS

$H(x_n | x_{n-1}, \dots, x_1)$ IS NONINCREASING IN n AND HAS LIMIT $H'(X)$. CONDITIONING REDUCES ENTROPY!!!

PROOF: $H(x_{n+1} | x_1, x_2, \dots, x_n) \leq H(x_{n+1} | x_1, \dots, x_n) =$
 $= H(x_n | x_{n-1}, \dots, x_1)$

DUE TO DEFINITION OF STATIONARY STOCHASTIC PROCESS.

Theorem 4.2.3 (Cesaro mean) IF $a_n \rightarrow a$ AND $b_n = \frac{1}{n} \sum_{i=1}^n a_i$ THEN $b_n \rightarrow a$

PROOF: Let $\epsilon > 0$. Since $a_n \rightarrow a$ there exist NUMBER such that $|a_n - a| < \epsilon$ for all $n \geq N(\epsilon)$

Hence,

$$|b_n - a| = \left| \frac{1}{n} \sum_{i=1}^n (a_i - a) \right| \leq \frac{1}{n} \sum_{i=1}^n |a_i - a|$$

$$\leq \frac{1}{n} \sum_{i=1}^{N(\epsilon)} |a_i - a| + \frac{n - N(\epsilon)}{n} \cdot \epsilon \leq \frac{1}{n} \sum_{i=1}^{N(\epsilon)} |a_i - a| + \epsilon$$

First term goes to 0 as $n \rightarrow \infty$, we can make $|b_n - a| \leq 2\epsilon$ by taking n large enough. Hence, $b_n \rightarrow a$ as $n \rightarrow \infty$.

PROOF OF THEOREM 4.2.1 BY CHAIN RULE

$$H(x_1, x_2, \dots, x_n) = \sum_{i=1}^n H(x_i | x_{i-1}, \dots, x_1)$$

$$\frac{H(x_1, x_2, \dots, x_n)}{n} = \frac{1}{n} \sum_{i=1}^n H(x_i | x_{i-1}, \dots, x_1)$$

ENTROPY RATE IS AVERAGE OF CONDITIONAL ENTROPIES. (TIME)

(MHV)

$$\frac{1}{n} \sum_{i=1}^n H(x_i | x_{i-1}, \dots, x_1) = H(X_n | X_{n-1}, \dots, X_1)$$

$\rightarrow H(X) = \lim_{n \rightarrow \infty} \frac{1}{n} H(x_1, x_2, \dots, x_n) \stackrel{\text{T4.2.3}}{=} \lim_{n \rightarrow \infty} H(x_n | x_{n-1}, \dots, x_1) = H'(X)$

(PROOF IN 168)

ACT FOR STATIONARY ERGODIC PROCESS

FOR ANY STATIONARY ERGODIC PROCESS:

$$-\frac{1}{n} \log \gamma(x_1, x_2, \dots, x_n) \rightarrow H(X)$$

MARKOV CHAINS

FOR STATIONARY MARKOV CHAIN, THE ENTROPY RATE IS GIVEN BY:

$$H(X) = H'(X) = \lim_{n \rightarrow \infty} H(x_n | x_{n-1}, \dots, x_1) = \lim_{n \rightarrow \infty} H(x_n | x_{n-1}) = \lim_{n \rightarrow \infty} H(x_2 | x_1)$$

CONDITIONAL ENTROPY IS CALCULATED USING THE GIVEN STATIONARY DISTRIBUTION. STATIONARY DISTRIBUTION μ IS SOLUTION OF THE EQUATIONS

$$\mu_j = \sum_i \mu_i p_{ij} \text{ for all } j$$

$$[M_1 \ M_2] = [M_1 \ M_2] \begin{bmatrix} 1-\alpha & \beta \\ \alpha & 1-\beta \end{bmatrix} = [M_i] \cdot [P_{ij}]$$

$$M_i = \sum_j M_j P_{ij}$$

(MOŽA INTERPRETACIJA SUMA NA VELOJATNOŠTICE-TRANSICIJI KOJI OD BILIO KOJA SOSTOJA "1" VOĐAT VO SOSTOJA "j")

$$M_1 = M_1(1-\alpha) + M_2 \cdot \beta = \sum_{i=1}^2 P_{i1} \cdot M_i$$

THEOREM 4.2.4

Let $\{x_n\}$ be a STATIONARY MARKOV CHAIN WITH STATIONARY DISTRIBUTION AND TRANSITION MATRIX P .

Then the ENTROPY RATE is:

$$H(x) = - \sum_{i,j} p_{ij} \log p_{ij}$$

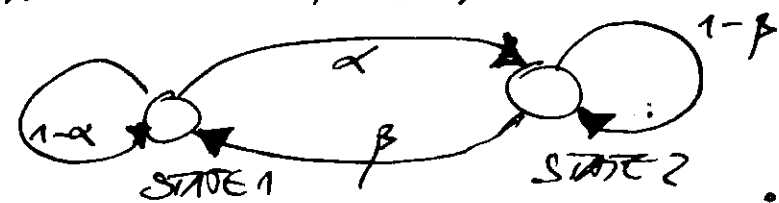
OVA E KLASIFA: PROCESOR MOŽE BITI U VO SOSTOJA "i" i e {1,2,...,n} I PEOGŽA VO SOSTOJA "j" SO VEROVATNOST P_{ij} I KO VOŠE E VO SOSTOJA "j"

PROOF: $H(x) = H(x_2|x_1) = \sum_i p_{i1} \sum_j P_{ij} \log \frac{1}{P_{ij}}$

$$H(Y|X=x) = - \sum_{Y \in Y} P(Y|X=x) \log P(Y|X=x)$$

$$H(Y|X) = \overline{H(Y|X=x)} = - \sum_{x \in X} P(x) \sum_{Y \in Y} P(Y|x) \log P(Y|x) = \sum_{x \in X} P(x) H(Y|X=x)$$

EXAMPLE 4.2.1 (TWO STATE MARKOV CHAIN) THE ENTROPY RATE OF TWO STATE MARKOV CHAIN IS:



$$H(x) = - \sum_{i,j} M_i P_{ij} \log P_{ij}$$

$$P_{ij} = \begin{bmatrix} 1-\alpha & \alpha \\ \beta & 1-\beta \end{bmatrix}$$

$$= \sum_i p_{i1} \sum_j P_{ij} \log \frac{1}{P_{ij}} = M_1 \sum_j P_{1j} \log \frac{1}{P_{1j}} + M_2 \sum_j P_{2j} \log \frac{1}{P_{2j}} = M_1 [\alpha \log \frac{1}{\alpha} + (1-\alpha) \log \frac{1}{1-\alpha}] + M_2 [\beta \log \frac{1}{\beta} + (1-\beta) \log \frac{1}{1-\beta}]$$

$$H(x) = \frac{\beta}{\alpha + \beta} H(\alpha) + \frac{\alpha}{\beta + \alpha} H(\beta)$$

Mod 1 izvisken pakek za markov proces so
 dve sostojbi $\{x_1, x_2\}$ $x_1 \in (0, 1)$ $P(x_1) = \begin{bmatrix} \frac{1}{2} & \frac{1}{2} \\ 0 & 0 \end{bmatrix}$
 $x_2 \in (0, 1)$ $P(x_2) = \begin{bmatrix} 0 & 0 \\ \frac{1}{3} & \frac{2}{3} \end{bmatrix}$

Procesor $\xi : \xi = x_1 \oplus x_2 \pmod 2$

$\xi \in [0, 1]$

x_1	x_2	ξ
0	0	0
0	1	1
1	0	1
1	1	0



x_1	x_2	ξ	μ
0	0	0	1

- ce ostane vo sostoji $\xi = 0$ ali:
- Novoto izveiceno $x_1 = 0$
 Novoto izveiceno $x_2 = 0$
 - Novoto izveiceno $x_1 = 1$
 Novoto $x_2 = 1$

Verjetnost da procesor ξ ne za spremeni
 sostojba vo "0" $\epsilon :$
 $P_0 = \frac{1}{2} \cdot \frac{1}{3} + \frac{1}{2} \cdot \frac{2}{3} = \frac{1}{6} + \frac{2}{6} = \frac{3}{6} = \frac{1}{2} = \mu_1$

ξ ce bide vo sostoji "1" ali:

- $x_1 = 0, x_2 = 1$ ali $x_1 = 1, x_2 = 0$

 $P_1 = \frac{1}{2} \cdot \frac{2}{3} + \frac{1}{2} \cdot \frac{1}{3} = \frac{2}{6} + \frac{1}{6} = \frac{3}{6} = \frac{1}{2} = \mu_2$

ξ ce spremeni vo sostoja 1 ali sostoja

	x_1	x_2
$x_1=0$	0	1
$x_1=1$	0	1

$$P(1|0) = P(00) [P(x_1=0) + P(x_2=1)] + P(00) [P(x_1=1) + P(x_2=0)] + P(11) [P(x_1=0) + P(x_2=1)] + P(11) [P(x_1=1) + P(x_2=0)]$$

$$= (P(11) + P(00)) [P(x_1=0) + P(x_2=1)] + (P(11) + P(00)) [P(x_1=1) + P(x_2=0)]$$

$$= \left(\frac{2}{6} + \frac{1}{6}\right) \left[\frac{1}{2} + \frac{2}{3}\right] + \left(\frac{2}{6} + \frac{1}{6}\right) \left[\frac{1}{2} + \frac{1}{3}\right] = \frac{1}{2} \cdot \frac{7}{6} + \frac{1}{2} \cdot \frac{5}{6} = \frac{12}{12}$$

	u	$u=1$
$x_1=0$	1	0
$x_2=1$	1	0
$x_1=1$	1	0
$x_2=0$	1	0

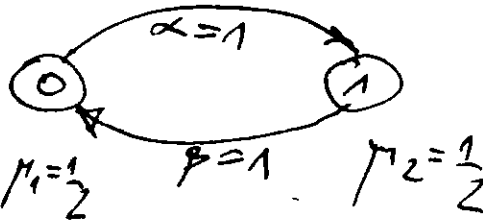
$$P(1|0) = P(0) [P(x_1=1) + P(x_2=1)] + P(0) [P(x_1=0) + P(x_2=0)] + P(1) [P(x_1=0) + P(x_2=0)] + P(1) [P(x_1=1) + P(x_2=1)]$$

$$P(1|0) = [P(0) + P(1)] \cdot [P(x_1=1) + P(x_2=1) + P(x_1=0) + P(x_2=0)]$$

$$= \left(\frac{2}{6} + \frac{1}{6}\right) \left(\frac{1}{2} + \frac{2}{3}\right) + \left(\frac{2}{6} + \frac{1}{6}\right) \left(\frac{1}{2} + \frac{1}{3}\right)$$

$$= \frac{1}{2} \frac{3+4}{6} + \frac{1}{2} \frac{3+2}{6} = \frac{12}{12}$$

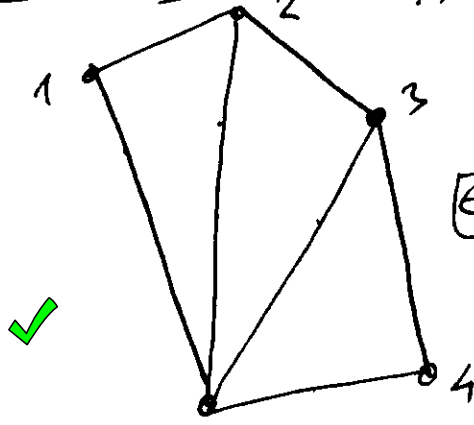
$$\alpha=1 \quad \beta=1$$



ANO GO
MORZIAM
VO MATLE
BI TRAZO
DA PODJAM
NAJEMICA
SERVENCA

4.3

EXAMPLE: ENTROPY RATE OF RANDOM WALK ON THE WEIGHTED GRAPH.



$$E=7$$

$$w = \frac{1}{7}$$

CONSIDER GRAPH WITH n NODES LABELED $\{1, 2, \dots, n\}$, WITH

WEIGHT $W_{ij} \geq 0$ ON THE EDGE JOINING i TO NODE j . $W_{ij} = W_{ji}$ } UNDIRECTED GRAPH

$W_{ij} = 0$ IF THERE IS NO EDGE CONNECTING i & j

A PARTICLE WALKS RANDOMLY FROM NODE TO NODE. THE RANDOM WALK $\{X_n\}$, $X_n \in \{1, 2, \dots, n\}$

GIVEN $X_n = i$, THE NEXT VERTEX j IS CHOSEN FROM THE NODES CONNECTED TO NODE i WITH PROBABILITY PROPORTIONAL TO THE WEIGHT OF THE EDGE CONNECTING i TO j

$$P_{ij} = \frac{W_{ij}}{\sum_k W_{ik}}$$

THE STATIONARY DISTRIBUTION FOR THIS MARKOV CHAIN ASSIGNS PROBABILITY TO NODE i . LET

$W_i = \sum_j W_{ij}$ BE TOTAL WEIGHT OF EDGES ORIGINATING FROM NODE i , AND LET $W = \sum_{i,j} W_{ij}$ BE THE SUM OF WEIGHTS OF ALL EDGES. THEN: $\sum W_i = 2W$

- NOW WE GUESS THAT STATIONARY DISTRIBUTION IS:

$$M_i = \frac{W_i}{2W}$$

WE CHECK THIS IS STATIONARY DISTRIBUTION BY CHECKING

$$\mu P = \mu$$

$$\sum_i \mu_i P_{ij} = \sum_i \frac{w_i}{2W} \frac{w_{ij}}{w_i} = \frac{w_j}{2W} = \mu_j$$

THUS, THE STATIONARY PROBABILITY OF STATE i IS PROPORTIONAL TO THE WEIGHT OF THE EDGES EMANATING FROM NODE i .

ENTROPY RATE IS:

$$H(x) = H(x_2 | x_1) = - \sum_i \mu_i \sum_j P_{ij} \log P_{ij} =$$

$$= - \sum_i \frac{w_i}{2W} \sum_j \frac{w_{ij}}{w_i} \log \frac{w_{ij}}{w_i} = - \sum_{i,j} \frac{w_{ij}}{2W} \log \frac{w_{ij}}{w_i}$$

$$= - \sum_{i,j} \frac{w_{ij}}{2W} \log \frac{w_{ij}}{2W} + \sum_{i,j} \frac{w_{ij}}{2W} \log \frac{w_i}{2W} =$$

$$= H(\dots, \frac{w_{ij}}{2W}, \dots) - H(\dots, \frac{w_i}{2W}, \dots)$$

IF ALL EDGES HAVE EQUAL WEIGHT, STATIONARY DISTRIBUTION PUTS WEIGHT $\frac{1}{2E}$ ON NODE i , WHERE E IS NUMBER OF EDGES EMANATING FROM NODE i AND E IS TOTAL NUMBER OF EDGES IN THE GRAPH.

$$- \sum_i \sum_j \frac{w_{ij}}{2W} \log \frac{w_{ij}}{2W} = \sum_i \sum_j \frac{w_{ij}}{2E} \log \frac{w_{ij}}{2E} =$$

$$\left(\frac{w_{ij} = w = \text{const}}{2E} \right) = - \frac{w}{2E} \sum_i \sum_j \log \frac{w}{2E} =$$

$$= \left(\frac{w = \frac{1}{E}}{2E} \right) = - \frac{1}{2E} \cdot 2E \cdot \log \frac{1}{2E} = \frac{1}{E} \log 2E^2 = \frac{2}{E} \log(E) \quad \text{FOR } w = \frac{1}{E}$$

OR IS TO EXAMPLE 4.2: (2E w = 1)

$$\sum_{i=1}^7 \sum_{j=1}^7 \frac{w}{2E} \log \frac{w}{2E} = \frac{w_{12}}{2E} \log \frac{w_{12}}{2E} + \frac{w_{15}}{2E} \log \frac{w_{15}}{2E} +$$

$$+ \frac{w_{21}}{2E} \log \frac{w_{21}}{2E} + \frac{w_{25}}{2E} \log \frac{w_{25}}{2E} + \frac{w_{27}}{2E} \log \frac{w_{27}}{2E} + \frac{w_{32}}{2E} \log \frac{w_{32}}{2E} +$$

$$\frac{w_{35}}{2E} \log \frac{w_{35}}{2E} + \frac{w_{34}}{2E} \log \frac{w_{34}}{2E} + \frac{w_{42}}{2E} \log \frac{w_{42}}{2E} + \frac{w_{45}}{2E} \log \frac{w_{45}}{2E} +$$

$$+ \frac{w_{51}}{2E} \log \frac{w_{51}}{2E} + \frac{w_{52}}{2E} \log \frac{w_{52}}{2E} + \frac{w_{53}}{2E} \log \frac{w_{53}}{2E} + \frac{w_{54}}{2E} \log \frac{w_{54}}{2E} = 2E \cdot \frac{1}{2E} \log \frac{1}{2E} = \log 2E$$

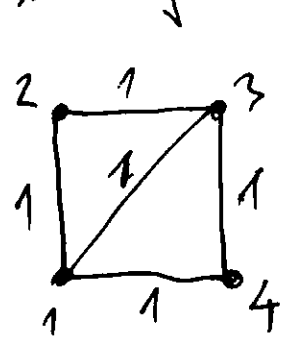
FOL $w=1 \Rightarrow - \sum_i \sum_j \frac{w_{ij}}{2W} \log \frac{w_{ij}}{2W} = \log 2E$

→ FOR EQUATE EDGES:

$$H(x) = \underbrace{\log(2E)}_{\text{PROOF PP.83}} - H\left(\frac{E_1}{2E}, \frac{E_2}{2E}, \dots, \frac{E_n}{2E}\right)$$

$$-H\left(\frac{E_1}{2E}, \frac{E_2}{2E}, \dots, \frac{E_n}{2E}\right) = \frac{E_1}{2E} \log \frac{2E}{E_1} + \frac{E_2}{2E} \log \frac{2E}{E_2} + \dots + \frac{E_n}{2E} \log \frac{2E}{E_n}$$

$$\sum_i \sum_j \frac{w_{ij}}{2W} \log \frac{w_{ij}}{2W} = \sum_{i=1}^4 \sum_{j=1}^4 \frac{1}{2E} \log \frac{w_{ij}}{2E} =$$



$$= \frac{w_{12}}{2E} \log \frac{w_{12}}{2E} + \frac{w_{14}}{2E} \log \frac{w_{14}}{2E} + \frac{w_{13}}{2E} \log \frac{w_{13}}{2E} + \frac{w_{21}}{2E} \log \frac{w_{21}}{2E} +$$

$$+ \frac{w_{23}}{2E} \log \frac{w_{23}}{2E} + \frac{w_{31}}{2E} \log \frac{w_{31}}{2E} + \frac{w_{32}}{2E} \log \frac{w_{32}}{2E} + \frac{w_{34}}{2E} \log \frac{w_{34}}{2E} +$$

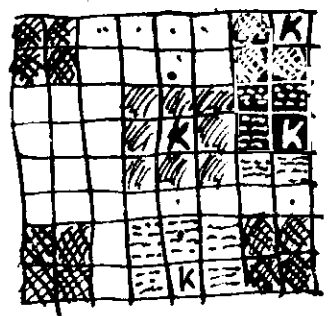
$$+ \frac{w_{41}}{2E} \log \frac{w_{41}}{2E} + \frac{w_{42}}{2E} \log \frac{w_{42}}{2E} = \textcircled{3} \frac{1}{2E} \log \frac{w_1}{2E} + \textcircled{2} \frac{1}{2E} \log \frac{w_2}{2E}$$

$$+ \textcircled{3} \frac{1}{2E} \log \frac{w_3}{2E} + \textcircled{2} \frac{1}{2E} \log \frac{w_4}{2E} = \frac{w_1}{2E} \log \frac{w_1}{2E} + \frac{w_2}{2E} \log \frac{w_2}{2E} +$$

$$+ \frac{w_3}{2E} \log \frac{w_3}{2E} + \frac{w_4}{2E} \log \frac{w_4}{2E} = \sum_{i=1}^4 \frac{w_i}{2E} \log \frac{w_i}{2E} =$$

$$= -H\left(\frac{w_1}{2E}, \frac{w_2}{2E}, \frac{w_3}{2E}, \frac{w_4}{2E}\right) \quad \text{DOCCENNO!!!}$$

EXAMPLE 4.3.1 (RANDOM WALK ON A CHESSBOARD) LET A KING MOVE AT RANDOM ON 8x8 CHESSBOARD.



$$H(x) = \log 2E - H\left(\frac{E_1}{2E}, \frac{E_2}{2E}, \dots, \frac{E_n}{2E}\right)$$

- i = — INTERIOR 4 КОМПАРАЦИИ ВНЕШЕГО
- e = — EDGES 2 КОМПАРАЦИИ ПО СРЕДНЕМУ
- c = — CENTER 4x4 КОМПАРАЦИИ = 4

РАЗЕТИ ≙ EDGE w_{ij}

ВНЕШНИХ РАБОТ НА EDGES (w_{ij}) : (i) $4 \times 8 = 32$ EDGES

(e) $8 \times 5 = 40$ EDGES (c) $4 \times 7 = 12$ EDGES

(84) $W = E = 32 + 40 + 12 = 72 + 12 = 84$ EDGES (8) ?? P.85.

$$H(x) = \log 168 + 4 \left(\frac{8}{168} \log \frac{8}{168} \right) + 8 \left(\frac{5}{168} \log \frac{5}{168} \right) + 4 \left(\frac{3}{168} \log \frac{3}{168} \right)$$

$$\mu_i = \frac{w_i}{2W} \quad \mu_1 = \frac{8}{168} \quad \mu_2 = \frac{5}{168} \quad \mu_3 = \frac{3}{168}$$

①	36 POZICII	X 8 LINKA	= 288 LINKA	} DVA E KAVIČO + NE \$
②	24 POZICII	X 5 LINKA	= 120 LINKA	
③	4 POZICII	X 3 LINKA	= 12 LINKA	
			<u>2E = TOTAL = 420 LINKA</u>	076105

ZNAČI:

$$\mu_1 = \frac{8}{420} \quad \mu_2 = \frac{5}{420} \quad \mu_3 = \frac{3}{420}$$

STACIONARITE VEŠTOBNOVI

$$H(x) = \log 420 + 36 \frac{8}{420} \log \frac{8}{420} + 24 \frac{5}{420} \log \frac{5}{420} + 4 \frac{3}{420} \log \frac{3}{420}$$

$$= \frac{1}{35} (72 \log 2 + 163 + 10 \log 5) = \frac{1}{35} \log(2^{72} \cdot 3 \cdot 5^{10}) =$$

$$= \frac{\log 420}{35} + \frac{288}{420} \log 8 - \frac{288}{420} \log 420 + \frac{120}{420} \log 5 - \frac{120}{420} \log 420 +$$

$$\frac{12}{420} \log 3 - \frac{12}{420} \log 3 = \frac{288}{420} \log 8 + \frac{120}{420} \log 5 + \frac{12}{420} \log 3$$

$$= \frac{288}{420} \log 8 + \frac{2}{7} \log 5 + \frac{2}{70} \log 3 = \frac{144}{210} \log 8 + \frac{2}{7} \log 5 + \frac{1}{35} \log 3$$

$$= \frac{72}{35} + \frac{2}{7} \log 5 + \frac{1}{35} \log 3 = \frac{72}{35} + \log(5^{\frac{2}{7}} \cdot 3^{\frac{1}{35}}) = \underline{2.76584}$$

$$= \frac{70}{35} + \frac{2}{35} \log 2 + \frac{40}{35} \log 5 + \frac{1}{35} \log 3 = 2 + \frac{1}{35} \log(4 \cdot 5^{40} \cdot 3) =$$

$$= 2 + \frac{1}{35} \log 12 + \frac{40}{35} \log 5 = 2 + \frac{1}{35} \log 12 + \frac{2}{7} \log 5 =$$

$$= 2 + \frac{1}{35} \log(12 \cdot 5^{40}) = \boxed{0.32 \cdot \log 8} \rightarrow \begin{matrix} \text{VO KMGATA} \\ \text{GO DOVEDUVAT} \\ \text{DO DVA FORMA.} \end{matrix}$$

ENTROPY RATE OF THE BOOKS

DVA E KA ŽABZI

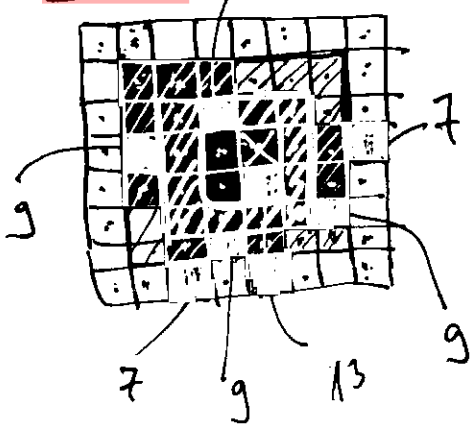
$$\mu = \frac{14}{2W}$$

$$2W = 14 \cdot 8^2 = 14 \cdot 64 = 896$$

$$\mu = \frac{14}{896}$$

$$H(x) = \log 896 + 64 \cdot \frac{14}{896} \cdot \log \frac{14}{896} = \log 896 \cdot \frac{14}{896} = \log 14 = \underline{3.8}$$

ДИСТОЯ



- I: 7 - MOVES (12V KONCENTRACIJA) 28 NODES
- II: 9 - MOVES 12+8=20 NODES
- III: 11 - MOVES 8+4=12 NODES
- IV: 13 - MOVES 4 NODES

$$\begin{array}{r} 11 \cdot 12 \\ \hline 22 \\ 11 \\ \hline 132 \end{array}$$

$$2E = 7 \cdot 28 + 9 \cdot 20 + 11 \cdot 12 + 13 \cdot 4 = 196 + 180 + 132 + 52 = 376 + 184 = 560 \text{ LINKS}$$

$$P_I = \frac{7}{560} \quad P_{II} = \frac{9}{560} \quad P_{III} = \frac{11}{560} \quad P_{IV} = \frac{13}{560}$$

$$H(x) = \log_2 560 + 28 \cdot \frac{7}{560} \log_2 \frac{7}{560} + 20 \cdot \frac{9}{560} \log_2 \frac{9}{560} + 12 \cdot \frac{11}{560} \log_2 \frac{11}{560} + 4 \cdot \frac{13}{560} \log_2 \frac{13}{560}$$

$$= \frac{196}{560} \log_2 7 + \frac{180}{560} \log_2 9 + \frac{132}{560} \log_2 11 + \frac{52}{560} \log_2 13$$

$$= \frac{1}{560} (196 \log_2 7 + 180 \log_2 9 + 132 \log_2 11 + 52 \log_2 13)$$

= 3.16053

СУБЖЕН

- I: 7+14=21 MOVES 28 NODES

$$\begin{array}{r} 21 \cdot 28 \\ \hline 168 \\ 42 \\ \hline 588 \end{array}$$

- II: 9+14=23 MOVES ; 20 NODES

- III: 11+14=25 MOVES ; 12 NODES

- IV: 13+14=27 MOVES ; 4 NODES

$$\begin{array}{r} 25 \cdot 12 \\ \hline 50 \\ 25 \\ \hline 300 \end{array}$$

$$2E = 21 \cdot 28 + 23 \cdot 20 + 25 \cdot 12 + 27 \cdot 4 = 588 + 460 + 300 + 108 = 1048 + 408 = 1456$$

$$P_I = \frac{21}{1456} \quad P_{II} = \frac{23}{1456} \quad P_{III} = \frac{25}{1456} \quad P_{IV} = \frac{27}{1456}$$

$$H(x) = \log_2 1456 + 28 \cdot \frac{21}{1456} \log_2 \frac{21}{1456} + 20 \cdot \frac{23}{1456} \log_2 \frac{23}{1456} + 12 \cdot \frac{25}{1456} \log_2 \frac{25}{1456} + 4 \cdot \frac{27}{1456} \log_2 \frac{27}{1456}$$

$$= \frac{1}{1456} (588 \log_2 21 + 460 \log_2 23 + 300 \log_2 25 + 108 \log_2 27) = 4.5125$$

86 КЕЗИЧАТА ИМА НАЈБОЉА ЕНТРОПИЈА. МОТ ТЕКУ-
ВАЈЕ. Е ПЕКА НАЈМЕНЕШНО Е КАДЕ ИЕ СЕ НАПРЕА С
ТОА ЕМА НАЈБОЉА СЛОБОДА НА ДУЗЕНЕ.

4.4. SECOND LAW OF THERMODYNAMICS

ENTROPY OF AN ISOLATED SYSTEM IS NON DECREASING.
 - IN STATISTICAL DYNAMICS ENTROPY IS \log^2 OF NUMBER OF MICROSTATES IN THE SYSTEM

① RELATIVE ENTROPY $D(\mu_n || \mu'_n)$ DECREASES WITH n .
 LET μ_n & μ'_n BE TWO PROBABILITY DISTRIBUTIONS ON THE STATE SPACE OF MARKOV CHAIN AT TIME n , AND μ_{n+1} AND μ'_{n+1} BE THE CORRESPONDING DISTRIBUTIONS AT TIME $n+1$. LET CORRESPONDING MASS FUNCTIONS BE DENOTED AS p & q . THUS:

$$p(x_n, x_{n+1}) = p(x_n) r(x_{n+1} | x_n) \quad q(x_n, x_{n+1}) = q(x_n) r(x_{n+1} | x_n)$$

WHERE $r(o|o)$ IS PROBABILITY TRANSITION FUNCTION FOR MARKOV CHAIN.

$$D(p(x_n, x_{n+1}) || q(x_n, x_{n+1})) = D(p(x_n) || q(x_n)) + D(p(x_{n+1} | x_n) || q(x_{n+1} | x_n)) + D(p(x_n | x_{n+1}) || q(x_n | x_{n+1}))$$

$$p(x_{n+1} | x_n) = q(x_{n+1} | x_n) = r(x_{n+1} | x_n) \Rightarrow$$

$$D(p(x_{n+1} | x_n) || q(x_{n+1} | x_n)) = 0$$

$$D(p(x_n) || q(x_n)) \geq D(p(x_{n+1}) || q(x_{n+1})) \quad \text{i.e.}$$

$$D(\mu_n || \mu'_n) \geq D(\mu_{n+1} || \mu'_{n+1})$$

CONSEQUENTLY THE DISTANCE BETWEEN THE PROBABILITY MASS FUNCTIONS IS DECREASING WITH TIME n FOR ANY MARKOV CHAIN. PRIMER: AND SE PREDSTAVI POU

PROBLEM OF SYSTEM ZA POKRETNOST NA SOPTVENOTT
 E IST VO UK, CANADA, SO TEK NA VREMJE
 DISTRIBUCIJA NA POZASTIVOD VO TE DVE ZEMJI
 JE STANE SLICNA.

② RELATIVE ENTROPY $D(\mu_n || \mu)$ BETWEEN A DISTRIBUTION μ_n ON THE STATES AT TIME n AND A STATIONARY DISTRIBUTION μ DECREASES WITH n .
 e.g. $\mu'_n = \mu \quad D(\mu_n || \mu) \geq D(\mu_{n+1} || \mu)$

WHICH IMPLIES THAT ANY STATE DISTRIBUTION CONVERGES TOWARDS STATIONARY DISTRIBUTION AS TIME PASSES.

5. ENTROPY INCREASES IF STATIONARY DISTRIBUTION IS UNIFORM.

IN GENERAL, THE FACT THAT RELATIVE ENTROPY DECREASES DOESN'T IMPLY THAT ENTROPY INCREASES.

IF WE START THE MARKOV CHAIN FROM UNIFORM DISTRIBUTION, THE DISTRIBUTION WILL TEND TO THE STATIONARY DISTRIBUTION, HENCE ENTROPY DECREASES WITH TIME.

IF STATIONARY DISTRIBUTION IS UNIFORM

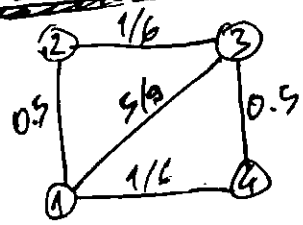
$$D(\gamma||\xi) = \sum_{i=1}^n \gamma_i \ln \frac{\gamma_i}{\xi_i} = \sum_{i=1}^n \gamma_i \ln \frac{\gamma_i}{\frac{1}{|Z|}} = |Z| + \sum_{i=1}^n \gamma_i \ln \gamma_i$$

$$D(\gamma||\xi) = |Z| - H(\gamma)$$

$$D(\mu||\pi) = |X| - H(\mu) = (|X| - H(\mu))$$

SINCE FOR MARKOV CHAIN WE HAVE MONOTONIC DECREASE OF RELATIVE ENTROPY IT IMPLIES MONOTONIC INCREASE OF ENTROPY. THIS EXPLANATION TIES MOST CLOSEST TO STATISTICAL THERMODYNAMIC WHERE ALL MICROSTATES ARE EQUALLY LIKELY.

PROBLEMA DA DISTRIBUIÇÃO MARKOV (EX. 2)



$E=5$

$2E=10$

$\mu_i = ? \quad \mu_1 = \frac{3}{10} \quad \mu_2 = \frac{2}{10} \quad \mu_3 = \frac{3}{10} \quad \mu_4 = \frac{2}{10}$

$P_{ij} = 0.2$

$\mu_j = \sum_i \mu_i \cdot P_{ij}$

$j=2 \quad \mu_2 = \mu_1 \cdot P_{12} + \mu_3 \cdot P_{32} = \frac{3}{10} \cdot 0.2 + \frac{3}{10} \cdot 0.2 = 0.2 \cdot 2 \cdot \frac{3}{10} = 0.4 \cdot \frac{3}{10} = \frac{1.2}{10}$

$\mu_2 = \mu_1 \cdot P + \mu_3 \cdot P = \frac{3}{10} \cdot P \quad P = \frac{10}{3} \cdot \mu_2 = \frac{10}{3} \cdot \frac{2}{10} = \frac{2}{3}$

$\mu_3 = \mu_1 \cdot P + \mu_2 \cdot P + \mu_4 \cdot P = \left(\frac{3}{10} + \frac{4}{10} \right) \cdot P = \frac{7}{10} \cdot \frac{2}{3} = \frac{14}{30}$

$[\mu_1 \mu_2 \mu_3 \mu_4] = [\mu_1 \mu_2 \mu_3 \mu_4] \begin{bmatrix} 0 & 0.2 & 0.2 & 0.2 \\ 0.2 & 0 & 0.2 & 0 \\ 0.2 & 0.2 & 0 & 0.2 \\ 0.2 & 0 & 0.2 & 0 \end{bmatrix}$

$$\mu_1 = 0.2(\mu_2 + \mu_3 + \mu_4)$$

$$\mu_2 = 0.2(\mu_1 + \mu_3)$$

$$\mu_3 = 0.2(\mu_1 + \mu_2 + \mu_4)$$

$$\mu_4 = 0.2(\mu_1 + \mu_2)$$

$$5 \cdot \mu_1 = 0.2(\mu_1 + \mu_2) + \mu_3 + 0.2(\mu_1 + 0.2(\mu_1 + \mu_2)) =$$

$$= \frac{0.2\mu_1}{\cancel{\quad}} + \frac{0.2\mu_2}{\cancel{\quad}} + \mu_3 + \frac{0.2\mu_1}{\cancel{\quad}} + \frac{0.04\mu_1}{\cancel{\quad}} + \frac{0.04\mu_2}{\cancel{\quad}} =$$

$$= (2 \cdot 0.2 + 0.04)\mu_1 + (1.2\mu_2 + 0.04\mu_2) = 0.44\mu_1 + 1.24\mu_2$$

$$\boxed{4.56\mu_1 = 1.24\mu_2}$$

$$[0.3, 0.2, 0.3, 0.2] = [0.3, 0.2, 0.3, 0.2] \begin{bmatrix} p_{11} & p_{12} & p_{13} & p_{14} \\ p_{21} & p_{22} & p_{23} & p_{24} \\ p_{31} & p_{32} & p_{33} & p_{34} \\ p_{41} & p_{42} & p_{43} & p_{44} \end{bmatrix}$$

$$= [0.3, 0.2, 0.3, 0.2] \begin{bmatrix} 0 & p_{12} & p_{13} & p_{14} \\ p_{12} & 0 & p_{23} & 0 \\ p_{13} & p_{23} & 0 & p_{34} \\ p_{14} & 0 & p_{34} & 0 \end{bmatrix} =$$

$$= [0.3, 0.2, 0.3, 0.2] \begin{bmatrix} 0 & p_1 & p_2 & p_3 \\ p_1 & 0 & p_3 & 0 \\ p_2 & p_3 & 0 & p_4 \\ p_3 & 0 & p_4 & 0 \end{bmatrix}$$

$$\frac{18}{10} = \frac{9}{5}$$

$$0.3 = 0.2 p_1 + 0.3 p_2 + 0.2 p_3$$

$$0.3 = 0.3 p_2 + 0.2 p_3 + 0.2 p_4$$

$$0.2 = 0.3 p_1 + 0.3 p_3$$

$$0.2 = 0.3 p_3 + 0.3 p_4$$

$$p_1 = p_4, \quad p_2 = 0.556, \quad p_3 = 0.67 - p_4, \quad \text{MARK}$$

$$\boxed{p_4 = 0.5 \quad p_1 = 0.5, \quad p_2 = 0.556 \quad p_3 = 0.17}$$

$$\frac{0.15}{0.051} = 0.02$$

$$0.3 \cdot p_1 + 0.3 p_2 = 0.5 \cdot 0.3 + 0.3 \cdot 0.17 = 0.15 + 0.051$$

$$\boxed{p_1 = 0.5; \quad p_2 = \frac{5}{9}; \quad p_3 = \frac{1}{6}; \quad p_4 = 0.5}$$

$$x^4 = \{x_1, x_2, x_3, x_4\}$$

$$P(x_2|x_1) = 0.5; \quad P(x_3|x_2) = 1/6; \quad P(x_3|x_1) = 5/9;$$

$$P(x_4|x_1) = 1/6; \quad P(x_3|x_4) = 0.5$$

$$P_{ij} = \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 0.5 & 1/6 & 1/6 \\ 0.5 & 0 & 1/6 & 0 \\ 5/6 & 1/6 & 0 & 0.5 \\ 1/6 & 0 & 0.5 & 0 \end{bmatrix}$$

$$\mu = [0.3, 0.2, 0.3, 0.2]$$

$$H(x) = -\sum_i \mu_i \sum_j P_{ij} \ln P_{ij}$$

$$H(x) = 1.2135 \quad \text{MARKA / MARKAS}$$

$$x_1 = \{1, 0, 0, 0\}$$

$$x_3 = \{0, 0, 1, 0\}$$

$$x_2 = \{0, 1, 0, 0\}$$

$$x_4 = \{0, 0, 0, 1\}$$

$$P(x_1) = \mu_1 (1-\mu_2)(1-\mu_3)(1-\mu_4) = 0.3 \cdot 0.8 \cdot 0.7 \cdot 0.8 = 0.1344$$

$$P(x_2) = \mu_2 (1-\mu_1)(1-\mu_3)(1-\mu_4) = 0.2 \cdot 0.7 \cdot 0.7 \cdot 0.8 = 0.0784$$

$$P(x_3) = \mu_3 (1-\mu_1)(1-\mu_2)(1-\mu_4) = 0.3 \cdot 0.7 \cdot 0.8 \cdot 0.8 = 0.1344$$

$$P(x_4) = \mu_4 (1-\mu_1)(1-\mu_2)(1-\mu_3) = 0.2 \cdot 0.7 \cdot 0.8 \cdot 0.7 = 0.0784$$

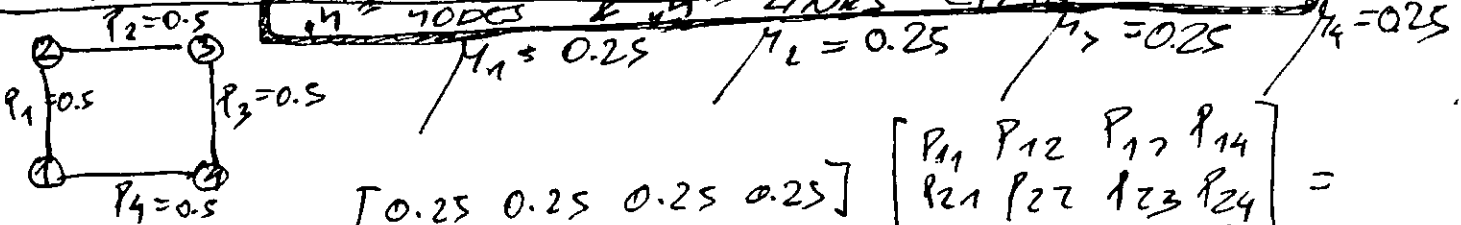
$$H(x_1, x_2, x_3, x_4) = -\sum_{i=1}^4 P(x_i) \ln P(x_i) = 0.93867$$

$$H(x) = \frac{1}{4} H(x_1, x_2, x_3, x_4) = \frac{1.35421}{4} = 0.23467$$

VALJA NE
DVA POKA
PROJKA:
VIDI PPT
KADJE OPAK
SO U=4
U=1

• AKO SE ŽELJE TOGA J $P(x_i) = \mu_i$

$$-\sum_{i=1}^4 \mu_i \ln \mu_i = 1.97095 \quad \frac{H(x_1, x_2, x_3, x_4)}{4} = \frac{1.97095}{4} = 0.5$$



HOZOV POKREK E
LEGLENA KEZ
KADJE PROJEZOT
E POKREK STOMATA!!

$$= [0.25 \ 0.25 \ 0.25 \ 0.25] \begin{bmatrix} 0 & P_1 & 0 & P_4 \\ P_1 & 0 & P_2 & 0 \\ 0 & P_2 & 0 & P_3 \\ P_4 & 0 & P_3 & 0 \end{bmatrix} = [0.25, 0.25, 0.25, 0.25]$$

$$\begin{aligned} 0.25 P_1 + 0.25 P_4 &= 0.25 & 0.25 P_2 + 0.25 P_3 &= 0.25 \\ 0.25 P_1 + 0.25 P_2 &= 0.25 & 0.25 P_3 + 0.25 P_4 &= 0.25 \\ P_1 + P_4 &= 1 & P_2 + P_3 &= 1 \\ P_1 + P_2 &= 1 & P_3 + P_4 &= 1 \end{aligned}$$

- КВАДРАТ: $P_1 = P_3$ $P_2 = 1 - P_3$, $P_4 = 1 - P_3$
- ПЕРИОДИЧ: $P_1 = 1 - P_3$; $P_2 = P_3$; $P_3 = 1 - P_2$; $P_4 = 1 - P_3$; $P_5 = P_3$

FOR $P_3 = 0.5$ $P_1 = 0.5$ $P_2 = 0.5$ $P_3 = 0.5$ $P_4 = 0.5$

$P_3 = 0.5 \Rightarrow P_1 = 0.5$ $P_2 = 0.5$ $P_3 = 0.5$ $P_4 = 0.5$ $P_5 = 0.5$

- ИСХОДНОСТЬ: $P_1 = P_3$; $P_2 = 1 - P_3$; $P_3 = P_3$; $P_4 = 1 - P_3$; $P_5 = P_3$; $P_6 = 1 - P_3$

- ПЕРИОДИЧ

$[P_{11} P_{21} P_{31} P_{41} P_{51}] \cdot \begin{bmatrix} 1 & 2 & 3 & 4 & 5 \\ 1 & 0 & 0.5 & 0 & 0.5 \\ 2 & 0.5 & 0 & 0.5 & 0 \\ 3 & 0 & 0.5 & 0 & 0.5 \\ 4 & 0 & 0 & 0.5 & 0 \\ 5 & 0.5 & 0 & 0 & 0 \end{bmatrix} = [0.2, 0.2, 0.2, 0.2, 0.2]$

- ТРИГОНОМЕТРИЧ

$[P_{11} P_{21} P_{31}] \cdot \begin{bmatrix} 1 & 2 & 3 \\ 1 & 0.5 & 0.5 \\ 2 & 0.5 & 0.5 \\ 3 & 0.5 & 0 \end{bmatrix}$

$H(X) = 1$

$[P_{11} P_{21} P_{31}] * (P_{ij} \cdot \log P_{ij}) = 1$ — СЕРИОЗНАЈ ЗА НЕКО КОЛКАВА МАТРИЦА

$-\sum_{i=1}^N \frac{1}{N} \cdot \log \frac{1}{N} = + \frac{N}{N} \cdot \log N = \log N = H(x_1, x_2, \dots, x_N)$
 $H(x_1, x_2, \dots, x_N) = \frac{\log 4}{4} \rightarrow 0$

$p(x_i) = p_1(1-p_2)(1-p_3) \dots (1-p_n) = p(1-p)^{n-1} = p(x_2) = \dots = p(x_n)$

$-\sum_{i=1}^N p(x_i) \log p(x_i) = -N \cdot p(x_i) \log p(x_i) = -N \cdot p(1-p)^{n-1} \log p(1-p)^{n-1}$
 $\lim_{N \rightarrow \infty} -N p(1-p)^{n-1} \log p(1-p)^{n-1} = -\lim_{N \rightarrow \infty} p(1-p) \log p(1-p)$

НЕ ТРЕБА ДА ПИШЕМО
 $N = 100$ ПИБЛЕШЕ
 4 ТРИГОМЕТРИЧНА МАТРИЦА

$-\lim_{p \rightarrow 0} p(1-p)^{n-1} \log p = -\lim_{n \rightarrow \infty} p(1-p)^{n-1} \log(1-p)^{n-1} = \left[\lim_{n \rightarrow \infty} (n-1)(1-p)^{n-1} \right] \cdot p \cdot \log(1-p)$

$$\lim_{n \rightarrow \infty} n(1-p)^{n-1} \log(n-p)^{n-1} = \lim_{n \rightarrow \infty} \frac{1}{n} \left(1 - \frac{1}{n}\right)^{n-1} \log\left(1 - \frac{1}{n}\right)^{n-1}$$

$$= \lim_{n \rightarrow \infty} \frac{1}{n} \left(\frac{n-1}{n}\right)^{n-1} \log\left(\frac{n-1}{n}\right)^{n-1} = \lim_{n \rightarrow \infty} \left(\frac{n-1}{n}\right)^n \cdot \log\left(\frac{n-1}{n}\right)$$

$$\lim_{n \rightarrow \infty} \frac{1}{n} \left(1 - \frac{1}{n}\right)^{n-1} \log \frac{1}{n} = \lim_{n \rightarrow \infty} \frac{1}{n} \left(\frac{n-1}{n}\right)^{n-1} \log \frac{1}{n}$$

• Ако се од формулата

$$H(X) = \log 2e - H\left(\frac{e_1}{2e}, \frac{e_2}{2e}, \dots, \frac{e_n}{2e}\right) \quad 4.667. \text{Sor. ORL}$$

$$= 0.5 \log \frac{2e}{0.5} = 0.5 \log 4e$$

- ТРИГОНИК \rightarrow бројито во модел сеува $N_{ij} = 0.5$
 $\rightarrow e_1 = 2 \cdot 0.5 = 1$

$$H(X) = 0.5 \log 4 \cdot 3 + 3 \cdot \left(\frac{1}{6} \log \frac{1}{6}\right) =$$

$$= 0.5 \log 12 + 0.5 \log \frac{1}{6} = 0.5 \log 2 + 0.5 \log 6 - 0.5 \log 6$$

$$= 1 + 0.5 \log 3$$

- СЕРПТИГОНИК

$$H(X) = 0.5 \log(4 \cdot 4) + 4 \cdot \left(\frac{1}{8} \log \frac{1}{8}\right) = 0.5 \log 2 + 0.5 \log 8 - 0.5 \log 8$$

$$= 0.5$$

▶ НЕМАТИ ЗЕМЕНО ВО ПРЕДИО ПЕНА $2e = 2 \cdot (N_e) \cdot 0.5 = N_e$

$$H(X) = -2 \cdot N_e \frac{0.5}{N_e} \log \frac{0.5}{N_e} - H\left(\frac{e_1}{N_e}, \frac{e_2}{N_e}, \dots, \frac{e_n}{N_e}\right) =$$

$$= \log 2(N_e) - H\left(\frac{e_1}{N_e}, \frac{e_2}{N_e}, \dots, \frac{e_n}{N_e}\right) \quad \text{MAV} \quad \text{OK!!!}$$

- ТРИГОНИК $H(X) = \log 2 \cdot 3 + 3 \cdot \frac{1}{3} \cdot \log \frac{1}{3} = \log 2 + \log 3 - \log 3 = 1$

- КВАДРАТ $H(X) = \log 2 \cdot 4 + 4 \cdot \frac{1}{4} \log \frac{1}{4} = \log 2 + \log 4 - \log 4 = 1$

- ПЕРПЕНДИКУЛАР $H(X) = \log 2 \cdot 5 + 5 \cdot \frac{1}{5} \log \frac{1}{5} = \log 2 + \log 5 - \log 5 = 1$

Пр. 90 ★ $X \in \{1, 2, 3, 4\}$ $m=4$ ДОЗНА НА КОДНАТА АЗБУКА.

- ИЗБОРОТ ГЕНЕРАТОР ПО ЕДЕН ЗНАК СО ДОЗНА $4=1$

$$P(X) = \{p_1, p_2, p_3, p_4\} = \{0.1344, 0.0784, 0.1744, 0.0784\}$$

ЕНТРОПИЈА НА СЕКВЕНЦИЈА Е:

$$H(x) = \frac{1}{4} H(x_1) = \frac{1}{1} H(x) = - \sum_{i=1}^4 p_i \log p_i = 1.35421$$

DP. 71 (★)

$$1.35421 \approx 1.2135$$

ZNAZI ZA PARAMETROT 5 LINKA 4 ZAZI DODIVAN TOVLA NA TEOLGYATA 4.24 I.E. REZULTATOT DOBIEN TAHA OGGOVARA NA REZULTATOT DOBIEN SOGLASNO DEFT-NICIZATA ZA ENTROPY RATE. SEPAK ZA PARAMETRE "LINKA & "NODES NE USIEVAM PA SO DODIVAN TOV.

DEFINITION A PROBABILITY TRANSITION MATRIX $[P_{ij}]$, IS CALLED DOUBLY STOCHASTIC

$$P_{ij} = P\{X_{t+1} = j | X_t = i\}$$

$$\sum_j P_{ij} = 1 \quad j = 1, 2, 3, \dots \text{ AND } \sum_i P_{ij} = 1 \quad i = 1, 2, 3, \dots$$

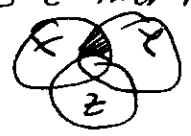
REMARK: THE UNIFORM DISTRIBUTION IS STATIONARY DISTRIBUTION OF P IF AND ONLY IF THE PROBABILITY TRANSITION MATRIX IS DOUBLY STOCHASTIC (THEOREM 4.1)

④ THE CONDITIONAL ENTROPY $H(X_t | X_1)$ INCREASES WITH t FOR A STATIONARY MARKOV PROCESS. IF THE MARKOV PROCESS IS STATIONARY, $H(X_1)$ IS CONSTANT, SO THE ENTROPY IS NONDECREASING, HOWEVER $H(X_t | X_1)$ IS INCREASING WITH t . THUS THE CONDITIONAL UNCERTAINTY OF THE FUTURE INCREASES.

$$H(X_2 | X_1) \geq H(X_2 | X_1, X_2) \text{ (CONDITIONING REDUCES ENTROPY)} =$$

$$= H(X_2 | X_2) = H(X_{t-1} | X_1)$$

BY MARKOVITY BY STATIONARITY



$$H(X_t | X_1) \geq H(X_{t-1} | X_1)$$

LOGICNO MI E!!
JO TEK NA VREM NE ZEMEZOSTA PASTE.

-INTELLIGENCE OF DATA-PROCESSING INEQUALITY

$$I(X; Y | Z) = H(X | Z) - H(X | Z, Y) = H(Y | Z) - H(Y | X, Z)$$

$$I(X; Y | Z) = -I(X; Z) + I(Z; Z | X)$$

$$I(X, Z; Y) = I(X; Y) + I(Z; Y | X) = I(Z; Y) + I(X; Y | Z)$$

$$I(X, Y; Z) = I(X; Z) + I(Y; Z | X) = I(Y; Z) + I(X; Z | Y)$$

$$I(Y; Z) \geq I(X; Z)$$

$$I(X; Y, Z) = I(X; Y) + I(X; Z | Y) = I(X; Z) + I(X; Y | Z)$$

$$I(X; Y) \geq I(X; Z)$$

DATA PROCESSING INEQUALITY 93

$$x_1 \rightarrow x_{n-1} \rightarrow x_n$$

$$I(x_1; x_{n-1}) \geq I(x_1; x_n)$$

$$\begin{aligned} H(x_1) - H(x_1|x_{n-1}) &\geq \\ H(x_1) - H(x_1|x_n) & \\ \hline H(x_1|x_{n-1}) &\leq H(x_1|x_n) \end{aligned}$$

- 1 ORRADO:

$$H(x_{n-1}) - H(x_{n-1}|x_n) \geq H(x_n) - H(x_n|x_{n-1})$$

$$\binom{3}{2} = \frac{6}{2} = 3$$

$$\begin{matrix} 011 \\ 101 \\ 110 \end{matrix}$$

STATIONARITY

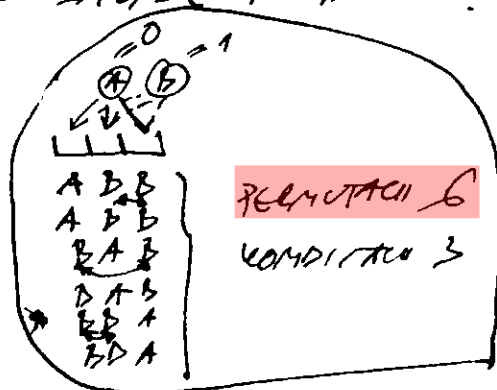
$$H(x_n) = H(x_{n-1}) \Rightarrow$$

$$H(x_{n-1}|x_n) \leq H(x_n|x_{n-1}) \quad \text{---} \quad H(x_n|x_n) \geq H(x_{n-1}|x_n)$$

⑤ SHUFFLES INCREASE ENTROPY. IF T IS A SHUFFLE (PERMUTATIONS) OF A DECK OF CARDS AND X IS THE INITIAL (RANDOM) POSITION OF THE CARDS IN THE DECK, AND IF THE CHOICE OF THE SHUFFLE T IS INDEPENDENT OF X , THEN

$$H(TX) \geq H(X)$$

WHERE TX IS PERMUTATION OF THE DECK INDUCED BY THE SHUFFLE T ON THE INITIAL PERMUTATION X . PROVEY 4.3 OUTLINES THE PROOF.



4.5 FUNCTIONS OF MARKOV CHAINS

LET x_1, x_2, \dots, x_n BE STATIONARY MARKOV CHAIN AND LET $y_i = \phi(x_i)$ BE A PROCESS EACH TERM OF WHICH IS A FUNCTION OF THE CORRESPONDING STATE IN THE MARKOV CHAIN. WHAT IS $H(Y)$? IT WOULD SIMPLIFY MATTERS IF y_1, y_2, \dots, y_n ALSO FORMED MARKOV CHAIN (MC), BUT IN MANY CASES IT IS NOT TRUE. SINCE MC IS STATIONARY SO IS x_1, x_2, \dots, x_n

$$H(x_n | x_{n-1}, x_{n-2}, \dots, x_1) \rightarrow H(X)$$

FOR STATIONARY PROCESSES $H(Y) = H(X)$ TH. 4.2.1

- FOR LOWER BOUND WE WILL USE

$$H(x_n | x_{n-1}, \dots, x_1) \text{ BASED ON THE}$$

IDEA THAT x_1 CONTAINS AS MUCH INFORMATION FOR x_n AS y_1, y_2, y_3, \dots

LEMMA 4.5.1

$$H(x_n | x_{n-1}, \dots, x_2, x_1) \leq H(Y)$$

PROOF: for $k=1, 2, \dots$ $H(Y_k | Y_{k-1}, \dots, Y_2, X_1) \stackrel{\text{SINCE } Y_1=f(X_1)}{=} H(Y_k | Y_{k-1}, \dots, Y_2, X_1)$

$H(Y_k | Y_{k-1}, \dots, Y_2, X_1) = \textcircled{*}$

$H(Y, Z | X) = H(Y | X) + H(Z | X, Y)$

$Y = f(X) \quad H(Y | X) = 0$

$H(Z | X, Y) = H(Y, Z | X) - H(Y | X)$

$H(Z, Y | X) = H(Z | X) + H(Y | X, Z) \stackrel{0}{=} H(Z | X) + H(Y | Z)$

$H(Z | X, Y) = H(Z | X) - H(Y | X, Z) = H(Z, Y | X)$

$H(Y_k | X_1, Y_1) = H(Y_k | X_1) + H(Y_1 | Y_k)$

$H(Y_k | X_1) = H(Y_k, Y_1 | X_1) - H(Y_1 | X_1, Y_k)$
IN DIRECTIO $= H(Y_k | X_1, Y_1)$ TWO GO ZNAES X1 GO ZNAES Y1

$H(Y_k, Y_1 | X_1) = H(Y_k | X_1) + H(Y_1 | X_1, Y_k) = H(Y_k | X_1)$
 $H(Y_k, Y_1 | X_1) = H(Y_k | X_1) + H(Y_1 | X_1, Y_k) = H(Y_k | X_1)$
 $H(Y_k | X_1, Y_1) = H(Y_k | X_1) + H(Y_1 | X_1, Y_k)$

USE EMAS: $H(Y_k | X_1, Y_1) = H(Y_k, Y_1 | X_1) - H(Y_1 | X_1)$ SO POKAZATE NA ZENOVITE:

$H(Y_k, Y_1 | X_1) = H(Y_k | X_1) + H(Y_1 | X_1, Y_k)$

$H(Y_k | X_1, Y_1) = H(Y_k | X_1)$ IF $Y_1 = f(X_1)$ MMV

$\textcircled{*} = H(Y_k | Y_{k-1}, \dots, Y_1, X_1, X_0, X_{-1}, \dots, X_{-k}) =$
DUE TO MARKOVITY OF Z
 $= H(Y_k | Y_{k-1}, \dots, Y_1, X_1, X_0, X_{-1}, \dots, X_{-k}, X_0, X_{-1}, \dots, X_{-k}) \leq$
 $\leq H(Y_k | Y_{k-1}, \dots, Y_1, X_0, X_{-1}, \dots, X_{-k}) = H(Y_{k+k+1} | Y_{k+k}, \dots, Y_1)$ SINCE $Y_i = f(X_i)$ STATIONARITY + k+1 CONDITIONING FOR ALL k
 $H(Y_k | Y_{k-1}, \dots, Y_1, X_1) \leq \lim_{k \rightarrow \infty} H(Y_{k+k+1} | Y_{k+k}, \dots, Y_1) = H(Y)$ 95

LEMMA 4.5.2

$$H(Y_n | Y_{n-1}, \dots, Y_1) - H(Y_n | Y_{n-1}, \dots, Y_1, X_1) \rightarrow 0$$

PROOF:

$$H(Y_n | Y_{n-1}, \dots, Y_1) - H(Y_n | Y_{n-1}, \dots, Y_1, X_1) = I(X_1; Y_n | Y_{n-1}, \dots, Y_1)$$

$$I(X_1; Y_1, Y_2, \dots, Y_n) = I(X_1; Y_1) + I(X_1; Y_2 | Y_1) + \dots + I(X_1; Y_n | Y_{n-1}, \dots, Y_1) = \sum_{i=1}^n I(X_1; Y_i | Y_{i-1}, \dots, Y_1)$$

$$I(X_1; Y) = H(X_1) - H(X_1 | Y) \quad I(X_1; Y) \leq H(X_1)$$

$$I(X_1; Y_1, Y_2, \dots, Y_n) \leq H(X_1)$$

$I(X_1; Y_1, Y_2, \dots, Y_n) \rightarrow$ INCREASES WITH n THUS $\lim_{n \rightarrow \infty} I(X_1; Y_1, \dots, Y_n)$ EXISTS AND

$$\lim_{n \rightarrow \infty} I(X_1; Y_1, Y_2, \dots, Y_n) \leq H(X_1)$$

$$H(X_1) \geq \lim_{n \rightarrow \infty} I(X_1; Y_1, Y_2, \dots, Y_n) = \lim_{n \rightarrow \infty} \sum_{i=1}^n I(X_1; Y_i | Y_{i-1}, \dots, Y_1)$$

$$H(X_1) \geq \sum_{i=1}^{\infty} I(X_1; Y_i | Y_{i-1}, \dots, Y_1)$$

SINCE THIS INFINITE SUM IS FINITE AND THE TERMS ARE NON-NEGATIVE, TERMS MUST TEND TO 0; I.E.:

$$\lim_{n \rightarrow \infty} I(X_1; Y_n | Y_{n-1}, \dots, Y_1) = 0 \quad \text{ZNAČI:}$$

$$H(Y_n | Y_{n-1}, \dots, Y_1) - H(Y_n | Y_{n-1}, \dots, Y_1, X_1) \rightarrow 0 \quad \text{DOKAZANO!!!}$$

COMBINING LEMMA 4.5.1 AND 4.5.2 WE HAVE:

THEOREM 4.5.1 If X_1, X_2, \dots, X_n FORM A STATIONARY MARKOV CHAIN, AND $Y_i = \phi(X_i)$, THEN:

$$H(Y_n | Y_{n-1}, \dots, Y_1, X_1) \leq H(Y) \leq H(Y_n | Y_{n-1}, \dots, Y_1)$$

AND: $\lim_{n \rightarrow \infty} H(Y_n | Y_{n-1}, \dots, Y_1, X_1) = H(Y) = \lim_{n \rightarrow \infty} H(Y_n | Y_{n-1}, \dots, Y_1)$

IN GENERAL WE COULD ALSO CONSIDER CASE WHERE Y_i IS STOCHASTIC FUNCTION (AS OPPOSED TO A DETERMINISTIC FUNCTION OF X_i).

CONSIDER A MARKOV PROCES x_1, x_2, \dots, x_n , AND DEFINE A NEW PROCESS y_1, y_2, \dots, y_n , WHERE EACH y_i IS DRAWN ACCORDING $p(y_i | x_i)$, CONDITIONALLY INDEPENDENT OF ALL THE OTHER $x_j, j \neq i$; THAT IS:

$$p(x^n, y^n) = p(x_1) \prod_{i=1}^{n-1} p(x_{i+1} | x_i) \prod_{i=1}^n p(y_i | x_i)$$

SUCH PROCESSES ARE CALLED GIBBS, MARKOV MIXTURE (HMM) AND IS USED IN SPEECH AND HANDWRITING RECOGNITION.

SUMMARY

ENTROPY RATE: TWO DEFINITIONS

$$H(X) = \lim_{n \rightarrow \infty} \frac{1}{n} H(x_1, x_2, \dots, x_n)$$

$$H'(X) = \lim_{n \rightarrow \infty} H(x_n | x_{n-1}, \dots, x_1)$$

- FOR STATIONARY STOCHASTIC PROCESSES

$$H(X) = H'(X)$$

ENTROPY RATE OF STATIONARY MARKOV CHAIN

$$H(X) = - \sum_{i,j} p_i p_{ij} \log p_{ij}$$

SECOND LAW OF THERMODYNAMICS. FOR MARKOV CHAIN:

- 1.) RELATIVE ENTROPY $D(p_n || p_1)$ DECREASES WITH TIME
- 2.) RELATIVE ENTROPY BETWEEN DISTRIBUTION AND THE STATIONARY DISTRIBUTION DECREASES WITH TIME.
- 3.) ENTROPY INCREASES IF $D(p_n || p)$ IS UNIFORM
- 4.) $H(x_n | x_1)$ INCREASES WITH TIME.
- 5.) $H(x_0 | x_n)$ INCREASES FOR ANY MARKOV CHAIN.

FUNCTIONS OF MARKOV CHAIN. IF x_1, x_2, \dots, x_n FORMS A STATIONARY MARKOV CHAIN AND $y_i = \phi(x_i)$

$$H(y_n | y_{n-1}, \dots, y_1, x_1) \leq H(Y) \leq H(x_n | x_{n-1}, \dots, x_1)$$

$$\lim_{n \rightarrow \infty} H(x_n | x_{n-1}, \dots, x_1, x_1) = H(Y) = \lim_{n \rightarrow \infty} H(y_n | y_{n-1}, \dots, y_1)$$

PROBLEMS

4.1 DOUBLY STOCHASTIC MATRICES. An $n \times n$ matrix

$P = [p_{ij}]$ is said to be doubly stochastic if $p_{ij} \geq 0$ and $\sum_j p_{ij} = 1$ for all i and $\sum_i p_{ij} = 1$ for all j .

An $n \times n$ matrix P is said to be PERMUTATION MATRIX if it is doubly stochastic and there is precisely one $p_{ij} = 1$ in each row and each column. It can be shown that every doubly stochastic matrix can be written as the convex combination of permutation matrices.

(a) Let $a^t = (a_1, a_2, \dots, a_n)$, $a_i \geq 0$, $\sum a_i = 1$ be a probability vector. Let $b = a \cdot P$ where P is doubly stochastic. Show that b is probability vector and that

$$H(b_1, b_2, \dots, b_n) \geq H(a_1, a_2, \dots, a_n)$$

Thus stochastic mixing increases the entropy.

(b) Show that stationary distribution μ^* for doubly stochastic matrix P is the uniform distribution.

(c) Conversely, prove that if the uniform is stationary distribution for Markov transition matrix P , then P is doubly stochastic.

(a)

$$[a_1, a_2, \dots, a_n] \cdot \begin{bmatrix} p_{11} & p_{12} & \dots & p_{1n} \\ p_{21} & p_{22} & \dots & p_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ p_{n1} & p_{n2} & \dots & p_{nn} \end{bmatrix} = \left[\underbrace{\sum_{i=1}^n a_i p_{i1}}_{b_1}, \underbrace{\sum_{i=1}^n a_i p_{i2}}_{b_2}, \dots, \underbrace{\sum_{i=1}^n a_i p_{in}}_{b_n} \right]$$

$$b_i = \sum_{i=1}^n a_i p_{i1} = a_1 \cdot p_{11} + a_2 p_{21} + \dots + a_n p_{n1}$$

$$p_{11} = 1 - \sum_{i=2}^n p_{i1} \quad p_{21} = 1 - \sum_{i=1, i \neq 2}^n p_{i1} \quad \dots \quad p_{n1} = 1 - \sum_{i=1}^{n-1} p_{i1}$$

b is probability vector hence $b_1 + b_2 + \dots + b_n = 1$

$$\sum_{i=1}^n a_i p_{i1} + \sum_{i=1}^n a_i p_{i2} + \dots + \sum_{i=1}^n a_i p_{in} = \sum_{j=1}^n \sum_{i=1}^n a_i p_{ij} = \sum_{i=1}^n a_i \underbrace{\sum_{j=1}^n p_{ij}}_1 = \sum_{i=1}^n a_i = 1 \quad \text{PROVED!}$$

1 (DUE TO DOUBLY STOCHASTIC) 1 (DUE TO PROBABILITY VECTOR)

(b)

$$\sum_{i=1}^n \mu_i P_{ij} = \mu_j$$

$$\mu_1 = \sum_{i=1}^n \mu_i P_{i1}$$

$$[\mu_1 \ \mu_2] \begin{bmatrix} P_{11} & P_{12} \\ P_{21} & P_{22} \end{bmatrix} = [\mu_1 \ \mu_2]$$

$$\begin{aligned} \mu_1 &= \mu_1 P_{11} + \mu_2 P_{21} \\ \mu_2 &= \mu_1 P_{12} + \mu_2 P_{22} \end{aligned}$$

$$\begin{aligned} \mu_1 (1 - P_{11}) &= \mu_2 P_{21} \\ \mu_2 (1 - P_{22}) &= \mu_1 P_{12} \end{aligned}$$

$$\mu_1 = \frac{\mu_2 P_{21}}{1 - P_{11}}$$

$$\mu_2 (1 - P_{22}) = \frac{\mu_2 P_{21} P_{12}}{(1 - P_{11})}$$

~~$P_{21} + P_{22} = 1 \Rightarrow P_{21} = 1 - P_{22}$
 $P_{11} + P_{12} = 1 \Rightarrow P_{12} = 1 - P_{11}$
 $\mu_1 P_{12} = \mu_2 P_{21}$
 $\mu_2 P_{21} = \mu_1 P_{12}$
 $\mu_1 (1 - P_{21}) = \mu_2 P_{21}$~~

\Rightarrow $P_{11} + P_{21} = 1$ $P_{21} = 1 - P_{11}$ $\mu_1 (1 - P_{11}) = \mu_2 (1 - P_{11})$

OVA + DOKAZ ZA 2x2 MATEMATIKA
 - VO GENERALIZACIJI

$$\mu_1 = \sum_{i=1}^n \mu_i P_{i1} = \mu_1 P_{11} + \mu_2 P_{21} + \mu_3 P_{31} + \dots + \mu_n P_{n1}$$

$$\mu_1 (1 - P_{11}) = \mu_2 P_{21} + \mu_3 P_{31} + \dots + \mu_n P_{n1}$$

$$1 - P_{11} = P_{21} + P_{31} + \dots + P_{n1}$$

$$\mu_1 P_{21} + \mu_1 P_{31} + \mu_1 P_{n1} = \mu_2 P_{21} + \mu_3 P_{31} + \dots + \mu_n P_{n1}$$

$$\Rightarrow \mu_1 = \mu_2 \quad \mu_1 = \mu_3 \quad \mu_1 = \mu_n$$

CONTINUE ON TP.100

$$[\mu_1 \ \mu_2 \ \dots \ \mu_n] \begin{bmatrix} P_{11} & P_{12} & \dots & P_{1n} \\ P_{21} & P_{22} & \dots & P_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ P_{n1} & P_{n2} & \dots & P_{nn} \end{bmatrix} = [\mu_1 \ \mu_2 \ \dots \ \mu_n]$$

$$\mu \cdot \sum_{i=1}^4 P_{i1} = \mu$$

$$\mu \cdot \sum_{i=1}^4 P_{i2} = \mu$$

$$\mu \cdot \sum_{i=1}^4 P_{i4} = \mu$$

$$\sum_{i=1}^4 P_{i1} = 1$$

$$\sum_{i=1}^4 P_{i2} = 1$$

$$\sum_{i=1}^4 P_{i4} = 1$$

$$\sum_{i=1}^4 P_{ij} = 1 \quad \text{for } j=1, \dots, 4$$

② EVEN GREATER GENERALIZATION:

1199 \$

$$\mu_1 P_{21} + \mu_1 P_{31} + \dots + \mu_1 P_{41} = \mu_2 P_{21} + \mu_3 P_{31} + \dots + \mu_4 P_{41}$$

$$\mu_1 \sum_{i=2}^4 P_{i1} = \sum_{i=2}^4 \mu_i P_{i1} \quad \sum_{i=2}^4 \mu_i P_{i1} - \sum_{i=2}^4 \mu_1 P_{i1} = 0$$

$$\sum_{i=2}^4 (\mu_i - \mu_1) P_{i1} = 0 \Rightarrow \underline{\mu_i = \mu_1 \quad i=2 \dots 4}$$

$$\textcircled{a} H(b_1, b_2, \dots, b_n) = - \left(\sum_{i=1}^n a_i P_{i1} \log \sum_{i=1}^n a_i P_{i1} + \sum_{i=1}^n a_i P_{i2} \log \sum_{i=1}^n a_i P_{i2} \right.$$

$$\left. + \dots + \sum_{i=1}^n a_i P_{in} \log \sum_{i=1}^n a_i P_{in} \right) \geq \text{CONCAVE FUNCTION}$$

$$\sum_{i=1}^n a_i P_{i1} = E[a_i] \quad \log E[a_i] \geq E[\log a_i]$$

$$\sum_{i=1}^n a_i P_{i1} \log \sum_{i=1}^n a_i P_{i1} \geq \sum_{i=1}^n a_i P_{i1} \left(\sum_{i=1}^n \log(a_i) P_{i1} \right)$$

$$\log a_1 P_{11} + \log a_2 P_{21} + \dots + \log a_n P_{n1}$$

$$H(a_1, a_2, \dots, a_n) = - \sum_{i=1}^n a_i \log a_i$$

LOG SUM INEQUALITY -
 $\sum_{i=1}^n a_i \log a_i \geq (\sum_{i=1}^n a_i) \log \frac{\sum_{i=1}^n a_i}{n}$

$$\begin{aligned} &\geq - \sum_{i=1}^n (P_{i1} a_i) \log (P_{i1} a_i) - \sum_{i=1}^n (P_{i2} a_i) \log (P_{i2} a_i) - \dots - \sum_{i=1}^n (P_{in} a_i) \log (P_{in} a_i) \\ &= - \sum_{j=1}^n \sum_{i=1}^n P_{ij} a_i \log (P_{ij} a_i) = - \sum_{j=1}^n \sum_{i=1}^n P_{ij} a_i \log P_{ij} - \sum_{j=1}^n \sum_{i=1}^n P_{ij} a_i \log a_i \end{aligned}$$

$$= - \sum_{i=1}^n a_i \sum_{j=1}^n p_{ij} \ln p_{ij} - \sum_{i=1}^n a_i \ln a_i \sum_{j=1}^n p_{ij} =$$

$$= \underbrace{\sum_{i=1}^n p_{ij} \ln \frac{1}{p_{ij}}}_{\geq 0} + \sum_{i=1}^n a_i \ln \frac{1}{a_i} \geq \sum_{i=1}^n a_i \ln \frac{1}{a_i} = H(a_1, a_2, \dots, a_n) !!!$$

PROVED !!!

EXERCISE 1 SOLUTIONS

SKORO IDENTIČNIM ST PONOVITE

(a) $H(b) - H(a) = - \sum_{i=1}^n (b_i) \ln b_i + \sum_{i=1}^n a_i \ln a_i =$

$$= - \sum_{j=1}^n \sum_{i=1}^n a_i p_{ij} \ln \sum_{k=1}^n a_k p_{kj} + \sum_{i=1}^n a_i \ln a_i =$$

POVORUVAJI "K"
ZA DA NE JEMEMO
SO "1" (100 N 100)

$$\sum_{i=1}^n a_i \ln a_i = \sum_{i=1}^n a_i \ln a_i \sum_{j=1}^n p_{ij} = \sum_{i=1}^n \sum_{j=1}^n a_i p_{ij} \ln a_i$$

$$= \sum_{i=1}^n \sum_{j=1}^n a_i p_{ij} \ln \frac{a_i}{\sum_{k=1}^n a_k p_{kj}} \geq \left(\sum_{i=1}^n a_i p_{ij} \right) \ln \frac{\sum_{i=1}^n a_i}{\sum_{i=1}^n b_i} =$$

LOG-SUM INEQUALITY:

$$\sum_{i=1}^n a_i \ln \frac{a_i}{b_i} \geq \left(\sum_{i=1}^n a_i \right) \ln \frac{\sum_{i=1}^n a_i}{\sum_{i=1}^n b_i}$$

$$= \sum_{i=1}^n a_i \sum_{j=1}^n p_{ij} \cdot \ln \frac{a_i}{\sum_{k=1}^n a_k p_{kj}} = 1 \cdot \ln 1 = 0$$

4.2 Times Arrow. Let $\{x_i\}_{i=-\infty}^{\infty}$ be a stationary stochastic process with

$$H(x_0 | x_{-1}, x_{-2}, \dots, x_{-n}) = H(x_0 | x_1, x_2, \dots, x_n)$$

$$H(x_{0+n} | x_1, x_2, \dots, x_n) = H(x_0 | x_1, x_2, \dots, x_n)$$

$$H(x_{4+n} | x_1, x_2, \dots, x_n) = H(x_0 | x_1, x_2, \dots, x_n)$$

$$H(x_{n+1} | x_1, \dots, x_n) = H(x_0 | x_1, \dots, x_n)$$

$$\boxed{H(x_1 | t_0) = H(x_0 | x_{-1})}$$

$$\begin{aligned} H(x_1, x_2, \dots, x_{n+1}) &= \sum_{i=1}^{n+1} H(x_i | x_{i-1}, \dots, t_n) = \\ &= H(x_1) + H(x_2 | x_1) + H(x_3 | x_1, x_2) + \dots + H(x_{n+1} | t_1, t_2, \dots, t_n) \end{aligned}$$

$$\begin{aligned} H(x_0, x_1, \dots, x_n) &= \sum_{i=0}^n H(x_i | x_{i-1}, \dots, x_0) = \\ &= H(x_0) + H(x_1 | t_0) + \dots + H(x_n | x_{n-1}, \dots, x_0) = \\ &= H(x_1, \dots, x_n, x_0) = H(x_1) + H(x_2 | x_1) + H(x_3 | x_1, x_2) + \dots \\ &+ \dots + H(x_0 | x_1, x_2, \dots, x_n) \end{aligned}$$

$$\begin{aligned} H(x_0, x_1, \dots, x_n) &= H(x_0+t, x_0+t, \dots, x_n+t) = \left|^{t=1} \right| = \\ &= H(x_1, x_2, \dots, x_{n+1}) \end{aligned}$$

Znači: $H(x_1, x_2, \dots, x_{n+1}) = H(x_1, x_2, \dots, x_n, t_0)$

$$\begin{aligned} \Rightarrow H(x_1) + H(x_2 | x_1) + H(x_3 | x_1, x_2) + \dots + H(x_n | x_1, x_2, \dots, x_{n-1}) + \\ + H(x_{n+1} | x_1, x_2, \dots, x_n) = H(x_1) + H(x_2 | x_1) + H(x_3 | x_1, x_2) + \dots \\ + H(x_n | x_{n-1}, x_{n-2}, \dots, x_1) + H(t_0 | x_n, x_{n-1}, \dots, x_1) \end{aligned}$$

Znači: $\boxed{H(x_{n+1} | x_1, x_2, \dots, x_n) = H(x_0 | t_n, t_{n-1}, \dots, t_1)}$
PROVED !!! $= H(t_0 | t_1, t_2, \dots, t_n)$

T.e. videti $H(x_{n+1} | x_1, x_2, \dots, x_n) = H(x_0 | x_{-1}, x_{-2}, \dots, x_{-n})$

$$\Rightarrow \boxed{H(x_0 | x_{-1}, x_{-2}, \dots, x_{-n}) = H(x_0 | x_1, x_2, \dots, x_n)}$$

OPTIONAL SOLUTION: $H(x_0 | x_{-1}, x_{-2}, \dots, x_{-n}) = H(x_0, x_1, \dots, x_{-n})$

$$\begin{aligned} \Rightarrow H(x_{-1}, x_{-2}, \dots, x_{-n}) &= H(x_n, \dots, x_1, t_0) - H(x_1, x_2, \dots, x_n) \\ &= H(x_0 | x_1, x_2, \dots, x_n) \end{aligned}$$

4.3 Shuffles increase ENTROPY. ARGUE THAT FOR ANY DISTRIBUTION ON SHUFFLES π AND ANY DISTRIBUTION ON CARD POSITIONS X THAT:

$$H(\pi X) \geq H(\pi X | \pi) = H(\pi^{-1} \pi X | \pi) = H(X | \pi) = H(X)$$

IF X AND π ARE INDEPENDENT

$$X = \{x_1, x_2, \dots, x_i, \dots, x_{51}, x_{52}\} \quad x_i \in \{1, 2, \dots, 52\}$$

e.g. $X = \{3, 4, 1, 2, \dots, 52, 51\}$

INITIAL RANDOM POSITION OF THE CARDS IN THE DECK

$$x_1 \rightarrow x_2 \rightarrow x_3 \rightarrow \dots \rightarrow x_n$$

$$I(x_1; x_2) \geq I(x_1; x_n)$$

$$x_2 = \pi x_1$$

$$x_1 \rightarrow \pi x_1 \rightarrow \pi^2 x_1 \rightarrow \dots \rightarrow \pi^{n-1} x_1$$

$$I(x_1; \pi x_1) \geq I(x_1; \pi^{n-1} x_1)$$

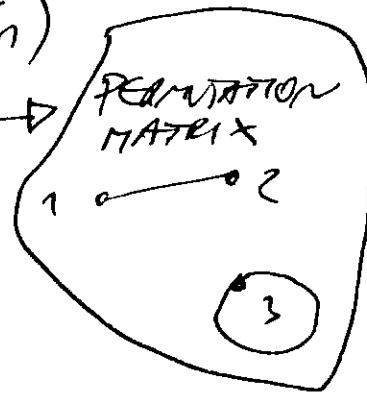
$$H(x_1) - H(x_1 | \pi x_1) \geq H(x_1) - H(x_1 | \pi^{n-1} x_1) \Rightarrow$$

$$H(x_1 | \pi x_1) \leq H(x_1 | \pi^{n-1} x_1)$$

THE CONDITIONAL ENTROPY DECREASES.

$$H(\pi x_1) - H(\pi x_1 | x_1) \geq H(\pi^{n-1} x_1) - H(\pi^{n-1} x_1 | x_1)$$

$$H(\pi x_1 | x_1) \leq H(\pi^{n-1} x_1 | x_1)$$



$$[1, 2, 3] \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} = [2, 1, 3]$$

$$[x_1, x_2, \dots, x_{52}] \begin{bmatrix} p_{11} & p_{12} & \dots & p_{1,52} \\ p_{21} & p_{22} & \dots & p_{2,52} \\ \dots & \dots & \dots & \dots \\ p_{52,1} & p_{52,2} & \dots & p_{52,52} \end{bmatrix} = [x_{\pi(1)}, x_{\pi(2)}, \dots, x_{\pi(52)}]$$

$\rightarrow f(x) = ?$

- AND 1 HAS 3 KATTA

$x_1 = 1, 2, 3; \quad x_2 = 2, 1, 3; \quad x_3 = 2, 3, 1; \quad x_4 = 3, 2, 1;$

$x_5 = 3, 1, 2; \quad x_6 = 1, 3, 2$

$P_k^y = \frac{y!}{(y-k)!} = \frac{3!}{0!} = 6$

$P_4^4 = \frac{4!}{0!} = 4!$

$H(X) = \sum_{i=1}^n \frac{1}{y!} \log(y!) = \log(y!)$

$H(X) = 225.581$ (with $y=52$)

$$T_1 = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$T_2 = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$$

$$T_3 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$T_4 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$$

$$T_5 = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$T_6 = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$$

$$T_7 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$T_8 = \begin{bmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{bmatrix}$$

$$T_9 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$$

$$T_{10} = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$$

$$T_{11} = \begin{bmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{bmatrix}$$

$$T_{12} = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$$

$$T_{13} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$$

$$T_{14} = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$$

$$T_{15} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$T_{16} = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

• bez potvorkovance

$$t_1 = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} \quad t_2 = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} \quad t_3 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \quad t_4 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$$

$$t_5 = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix} \quad t_6 = \begin{bmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{bmatrix}$$

$$P_2^4 = P_3^3 = \frac{3!}{0!} = 6$$

$$H(x) = \lg 4! = \lg 6 = 2.585.$$

• USTE POED OSOVEN PRIMER (n=2 miren RCO NA VESTI VOJITRO)

$$[x_1 \ x_2] \begin{bmatrix} T_{11} & T_{12} \\ T_{21} & T_{22} \end{bmatrix}$$

$$T_1 = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

$$T_2 = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$$

$$x_1^4 = [1 \ 2] \quad x_2^4 = [2 \ 1]$$

$$H(x^4) = \sum_{i=1}^4 \frac{1}{4!} \lg 4! = \lg 4!$$

$$H(x^4) = \lg 2 = 1$$

$$H(TX|T) = \underbrace{P(T=t_1)}_{1/2} \cdot H(x_1|t_1) + \underbrace{P(T=t_2)}_{1/2} \cdot H(x_2|t_2)$$

PERMUTACIONE MATRICI SE GENERALIZUJU DO UNIFORMNA RASILEDENA.

• $H(t_1 X | t_1) = ?$

$$X_1^u = \{1, 2\}$$

$$x_2^u = [2, 1]$$

$$t_1 = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

$$t_1 \cdot x_1^u = [2 \ 1]$$

$$t_1 \cdot x_2^u = [1 \ 2]$$

$$P(t_1 x_1^u) = \frac{1}{2!}$$

$$P(t_1 x_2^u) = \frac{1}{2!}$$

$$H(t_1 X | t_1) = \frac{1}{2!} \ln 2! + \frac{1}{2!} \ln 2! = \frac{1}{2} + \frac{1}{2} = 1$$

$$H(t_2 X | t_2) = \frac{1}{2} \ln 2 + \frac{1}{2} \ln 2 = 1$$

$$H(TX|T) = \frac{1}{2} + \frac{1}{2} = 1$$

• SE MOVAKAZUJE NA PRIMERU DO $4=3$ KARDI VO IPRUOT

$$H(TX|T) = \sum_{i=1}^6 P(T=t_i) H(t_i X | t_i)$$

IMO PERMUTACIONE MATRICI SE DO UNIFORMNA RASILE-
DENA:

$$H(TX|T) = \frac{1}{6} \sum_{i=1}^6 H(t_i X | t_i)$$

$$x^u = [x_1, x_2, x_3]$$

$$H(t_1 X | t_1) = ?$$

$$x_1^u = [1, 2, 3] \quad x_2^u = [1, 3, 2] \quad x_3^u = [2, 1, 3] \quad x_4^u = [2, 3, 1]$$

$$x_5^u = [3, 1, 2] \quad x_6^u = [3, 2, 1]$$

$$t_1 x_1^u = [2, 1, 3]; \quad t_1 x_2^u = [3, 1, 2]; \quad t_1 x_3^u = [1, 2, 3]; \quad t_1 x_4^u = [3, 2, 1]$$

$$t_1 x_5^u = [1, 3, 2]; \quad t_1 x_6^u = [2, 3, 1]$$

$$H(t_1 X | t_1) = \left(\frac{1}{6} \ln 6\right) \cdot 6 = \ln 6 = \ln(3!)$$

- DUK NE MOGA UNIFORMNO DA SE RASILEDENI
PERMUTACIONE MATRICI, KOJA

$$H(TX|T) = \sum_{i=1}^6 P(T=t_i) \cdot H(t_i X | t_i) = \ln 6 \cdot \sum_{i=1}^6 P(T=t_i) = \ln 6 \cdot 1 = \ln 6 = H(X)$$

OVAK DOKAZ MOZE DA SE MOZERI VO GENERALIZEN
SLUCAJ!!

$$H(TX|T) = \sum_{i=1}^{4!} P(T=t_i) H(TX|T=t_i) = \log(4!) \sum_{i=1}^{4!} P(T=t_i) = \log(4!) \cdot 1 = \log(4!) = H(X)$$

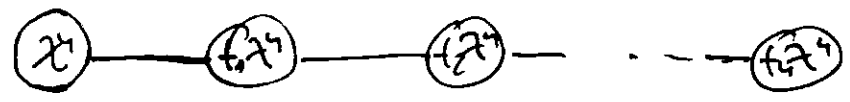
Ⓐ

$$H(TX|T) = \log(4!) = H(X)$$

$$H(X) \geq H(TX|T)$$

JA POVAZEVAN EP-AKUVOSTA !!!

- SO OGLASNA NA DOA STO SE PAMOBT ZA
 MARKOV KAZEL (MOZE I RANDOM WALK APRA-
 KLIZA PA SE NAKLAVI UADE 'T' C
 VEKINOST TAJ ZI CIGORATA MATRICA)



$$H(t_n x^2 | x^1) \leq H(t_n x^2 | x^2)$$



$$I(x; y) \geq I(x; z)$$

$$H(x) - H(x|y) \geq H(x) - H(x|z)$$

$$H(x|y) \leq H(x|z)$$

$$H(y) - H(y|x) \geq H(z) - H(z|x)$$

$H(y) = H(z) \rightarrow$ STATIONARNA MARKOV CIGORNA

$$H(y|x) \leq H(z|x)$$

$$H(TX) \geq H(TX|T) = H(T \cdot TX|T) = H(X|T) = H(X)$$

Ⓐ \Rightarrow CONDITIONING REDUCES ENTROPY

$$H(X, T) = H(X) + H(T|X)$$

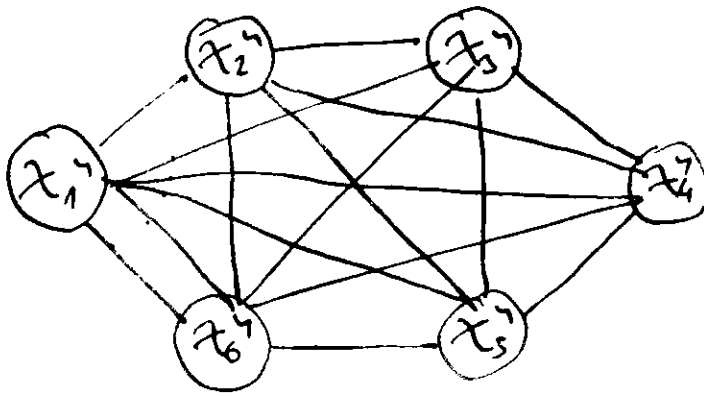
Ⓐ \Rightarrow X & T ARE STATISTICALLY INDEPENDENT

$$H(TX) = \sum_{t,x} p(t,x) \log p(t,x)$$

$$p(t,x) = p(t) \cdot p(x) = \frac{1}{4!} \cdot \frac{1}{4!}$$

$$H(TX) = \sum_{t,x} \frac{1}{(4!)^2} \log (4!)^2 = (4!)^2 \cdot \frac{1}{(4!)^2} \log (4!)^2$$

$$H(TX) = 2 \log(4!) = 2H(X)$$



$$x_1^1 = \{1, 2, 3\}$$

$$x_2^1 = \{1, 3, 2\}$$

$$x_3^1 = \{2, 1, 3\}$$

$$x_4^1 = \{2, 3, 1\}$$

$$x_5^1 = \{3, 1, 2\}$$

$$x_6^1 = \{3, 2, 1\}$$

$$P_{ij} = \sum_{i=1}^6 \frac{P_{ij}}{P_{ij}} \quad j = 1, 2, \dots, 4$$

$$P_{ij} = P(t_i)$$

$$H(X) = \sum_{i=1}^6 P_i \sum_{j=1}^6 P_{ij} \log P_{ij} = \sum_{i=1}^6 \sum_{j=1}^6 P_i P_{ij} \log P_{ij}$$

$$= \sum_{i=1}^6 \sum_{j=1}^6 \frac{1}{6!} \frac{1}{6!} \log 6! = \frac{6!}{6!^2} \log 6! = \log(6!)$$

• Ako odic to formulacija za HANJAN WAKK:

$$H(X) = \log 26 - H\left(\frac{E_1}{26}, \frac{E_2}{26}, \dots, \frac{E_7}{26}\right)$$

$$26 = 6 \cdot 6 = 36$$

$$\frac{E_1}{26} = \frac{E_2}{26} = \dots = \frac{E_6}{26} = \frac{6}{36} = \frac{1}{6}$$

$$H(X) = \log 36 - 6 \cdot \frac{1}{6} \cdot \log 6 = \log 36 - \log 6 = \log 6 + \log 6 - \log 6 = \log 6$$

ŠTITINSKATA KONSTANTATA KOTA SLEPI OD MA-
LIZADA NA PP. 105/106 E DEUA:

$$H(TX) = H(X|T)$$

ZAOA ŠTO SAMO SE POKREDENI SOSTOJATE, A
NE SE PROMENI VEROVATNOSTITE NA SOSTOJATE.

$$H(TX) \approx H(TX|T) = \sum_t \gamma(t) H(TX|T=t) =$$

$$= \sum_t \gamma(t) H(TX) = H(X) \sum_t \gamma(t) = H(X)$$

ISTA VREDNOST ZA
BICO KOT "t" = i.e. = H(X)

→ DUE TO CONDITIONING REDUCES ENTROPY.

- OVA E VŠUJNOST ROBINJAN DOKAZ NA ONA ŠTO SUVA 90
POKAZAL NA PP. 105-106

4.4 SECOND LAW OF THERMODYNAMICS. LET x_1, x_2, x_3

... BE A STATIONARY FIRST-ORDER MARKOV CHAIN. IN SECTION 4.4 IT WAS SHOWN THAT:

$$H(x_n | x_1) \geq H(x_{n-1} | x_1) \text{ for } n = 2, 3, \dots$$

THUS CONDITIONAL UNCERTAINTY ABOUT THE FUTURE GROWS WITH TIME. THIS IS TRUE ALTHOUGH THE UNCONDITIONAL UNCERTAINTY $H(x_n)$ REMAINS CONSTANT. HOWEVER, SHOW BY EXAMPLE THAT $H(x_n | x_1 = x_1)$ DOES NOT NECESSARILY GROW WITH "n" FOR EVERY x_1 .

$$H(x_n | x_1) = p(x_1 = x_1) H(x_n | x_1 = x_1) + p(x_1 = x_2) H(x_n | x_1 = x_2) + \dots + p(x_1 = x_n) H(x_n | x_1 = x_n)$$

EXAMPLE IF $n=3$

$$x_1 \rightarrow x_2 \rightarrow x_3 \quad I(x_1; x_2) \geq I(x_1; x_3)$$

$$H(x_1) - H(x_2 | x_1) \geq H(x_1) - H(x_3 | x_1)$$

STATIONARY

$$H(x_3 | x_1) \geq H(x_2 | x_1)$$

if $p(x_n) = \frac{1}{|x_n|} \Rightarrow$ UNIFORM DISTRO

$$H(x_2 | x_1) = \frac{1}{|x_1|} H(x_2 | x_1 = \frac{1}{|x_1|}) + \frac{1}{|x_1|} \cdot H(x_2 | x_1 = \frac{1}{|x_1|}) + \dots$$

$$H(x_2 | x_1) = \frac{1}{|x_1|} \cdot \frac{1}{|x_1|} H(x_2 | x_1 = \frac{1}{|x_1|}) = H(x_2 | x_1 = \frac{1}{|x_1|})$$

$$H(x_3 | x_1) = H(x_3 | x_1 = \frac{1}{|x_1|}) \Rightarrow H(x_3 | \frac{1}{|x_1|}) = H(x_2 | \frac{1}{|x_1|})$$

$$D(p \parallel q) = \sum_x p \log \frac{p}{q} = - \sum_x p \log \frac{1}{p} + \sum_x p \log |x|$$

(STATIONARY MARKOV CHAIN)

$$= \log |x| - H(x) = \log |x| - H(x_2) \quad \text{②}$$

POTENTIALS:

$$p_{n-1}, p_n; p_{n+1}, p_{n+2}$$

TRANSITION PROBABILITY

$$p(x_n, x_{n+1}) = p(x_n) \cdot p(x_{n+1} | x_n)$$

$$q(x_n, x_{n+1}) = q(x_n) \cdot r(x_{n+1} | x_n)$$

$$D(p(x_n, x_{n+1}) \parallel q(x_n, x_{n+1})) = D(p(x_n) \parallel q(x_n)) +$$

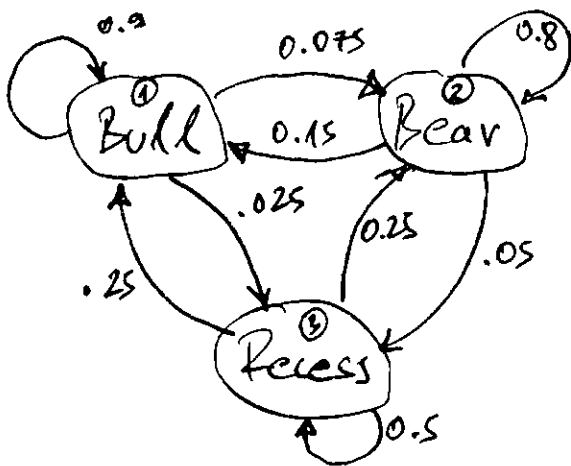
$$D(\underbrace{p(x_{t+1}|x_t)}_{\theta} \parallel \underbrace{q(x_{t+1}|x_t)}_{\theta}) = D(p(x_{t+1}) \parallel q(x_{t+1})) + D(p(x_t|x_{t+1}) \parallel q(x_t|x_{t+1}))$$

$$p(x_{t+1}|x_t) = q(x_{t+1}|x_t) = p(x_{t+1}|x_t) \Rightarrow \underline{D(p(x_{t+1}|x_t) \parallel q(x_{t+1}|x_t))} = \dots$$

$$\boxed{D(p(x_t) \parallel q(x_t)) \geq D(p(x_{t+1}) \parallel q(x_{t+1}))} \text{ i.e.}$$

$$D(\mu_t \parallel \mu_t) \geq D(\mu_{t+1} \parallel \mu_{t+1})$$

- Parameters za modelov citain od NIKIFENJA



• BULL WEEK IS FOLLOWED BY ANOTHER BULL WEEK 90% OF TIME

$$\mu_j = \sum_{i=1}^3 \mu_i P_{ij} \quad j=1$$

$$\mu_1 = \mu_1 \cdot P_{11} + \mu_2 \cdot P_{21} + \mu_3 \cdot P_{31}$$

$$\mu_1(1 - P_{11}) = \mu_2 P_{21} + \mu_3 P_{31}$$

$$\boxed{0.1 \mu_1 = 0.15 \mu_2 + 0.25 \mu_3}$$

$$\mu_2 = \sum_{i=1}^3 \mu_i P_{i2} = \mu_1 P_{12} + \mu_2 P_{22} + \mu_3 P_{32}$$

$$\mu_2(1 - P_{22}) = 0.075 \mu_1 + 0.25 \mu_3 \quad \boxed{0.2 \mu_2 = 0.075 \mu_1 + 0.25 \mu_3}$$

$$\mu_3 = \mu_1 \cdot P_{13} + \mu_2 P_{23} + \mu_3 P_{33}$$

$$\boxed{0.5 \mu_3 = 0.025 \mu_1 + 0.05 \mu_2}$$

$$\mu_1 = 0.625 \quad \mu_2 = 0.3125 \quad \mu_3 = 0.0625$$

$$D(\mu_t \parallel \mu_t) \geq D(\mu_{t+1} \parallel \mu_{t+1}) \quad \leftarrow \text{IF STATIONARY DISTRIBUTION IS UNIFORM}$$

$$\Rightarrow \log |x| - H(\mu_t) \geq \log |x| - H(\mu_{t+1})$$

IF STATIONARY DISTRIBUTION IS $p(x) = \frac{1}{2^n}$



$$D(\mu_t \parallel \mu) \geq D(\mu_{t+1} \parallel \mu)$$

$$D(M_{n+1}||\mu) = \sum_{x \in X} p(x) \log\left(\frac{p(x)}{2^n}\right) \quad I(x) = \frac{1}{2^n}$$

$$= -H(x^n) + \sum_{x \in X} p(x) \log 2^n = n - H(x^n)$$

CONST

$$D(M_{n+1}||\mu) = (n+1) - H(x^n)$$

$D(M_{n+1}||\mu) \geq D(M_n||\mu)$
NOT POSSIBLE!!!

$$H(x_n | x_1) \geq H(x_n | x_1, x_2) \stackrel{\uparrow \text{MARKOVITZ}}{=} H(x_n | x_2) \stackrel{\uparrow \text{MARKOVITZ}}{=} H(x_n | x_1)$$

$$H(x_n | x_{n-1}) = H(x_n | x_{n-1}, x_{n-2}) =$$

\downarrow DU TO KNOWING THE PRESENCE OF FUTURE DOESN'T DEPEND FROM PAST

$$= H(x_{n-1} | x_{n-2}, x_{n-3}) = H(x_{n-1} | x_{n-2})$$

STATIONARITY.

IF: $p(x_n) = p(x_{n-1}) \Rightarrow$

$$H(x_n | x_1) = H(x_{n-1} | x_{n-2})$$

NE E VOA
VOA !!!

IF: $x_{n-1} = f(x_n)$

$$H(x_{n-1} | x_n) = H(x_{n-1} | x_n, x_{n-1}) \stackrel{\uparrow \text{MARKOVITZ}}{=} H(x_{n-1} | x_n)$$

SE KADAM VO OVA !!!
POTANJE E OZI VO
GOD SEVAJ IMAJE
MILJEX CEMU

DOLAZIMO !!!

OPTION 1: $x_n \rightarrow x_{n-1} \rightarrow x_1$

$$I(x_n; x_{n-1}) \geq I(x_n; x_1) \quad H(x_{n-1}) - H(x_{n-1} | x_n) \geq$$

$$\geq H(x_n) - H(x_n | x_1) \quad H(x_n) = H(x_{n-1}) \Rightarrow$$

$$H(x_{n-1} | x_n) \leq H(x_n | x_1) \quad \text{i.e. } H(x_n | x_1) \geq H(x_{n-1} | x_n)$$

4.5 ENTROPY OF RANDOM TREE CONSIDER THE FOLLOWING METHOD OF GENERATING A RANDOM TREE WITH n NODES. FIRST EXPAND THE ROOT NODE:

THEN EXPAND ONE OF THE TERMINAL NODES AT RANDOM:



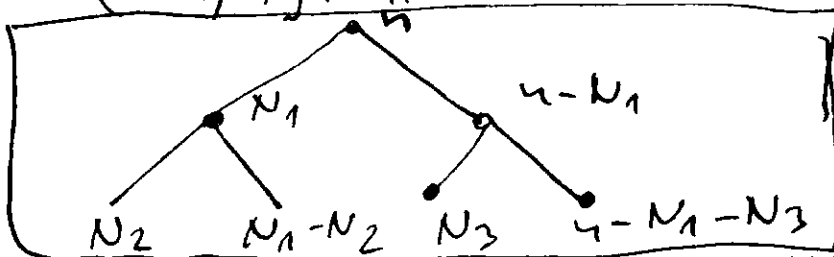
AT TIME k , CHOOSE ONE OF THE $k-1$ TERMINAL NODES ACCORDING TO THE UNIFORM DISTRIBUTION AND EXPAND IT. CONTINUE UNTIL u TERMINAL NODES HAVE BEEN GENERATED. THIS SEQUENCE LEADING TO A FIVE-NODE TREE MIGHT LOOK LIKE THIS:



SOMEHOW THE FOLLOWING METHOD OF GENERATING RANDOM TREES YIELDS THE SAME PROBABILITY DISTRIBUTION ON TREES WITH u TERMINAL NODES. FIRST CHOOSE AN INTEGER N_1 UNIFORMLY DISTRIBUTED ON $\{1, 2, \dots, u-1\}$. WE THEN HAVE STRUCTURE:



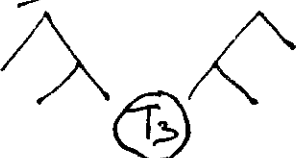
THEN CHOOSE AN INTEGER N_2 UNIFORMLY DISTRIBUTED OVER $\{1, 2, \dots, N_1-1\}$ AND INDEPENDENTLY CHOOSE ANOTHER INTEGER N_3 UNIFORMLY OVER $\{1, 2, \dots, (u-N_1)-1\}$. THE PICTURE IS NOW

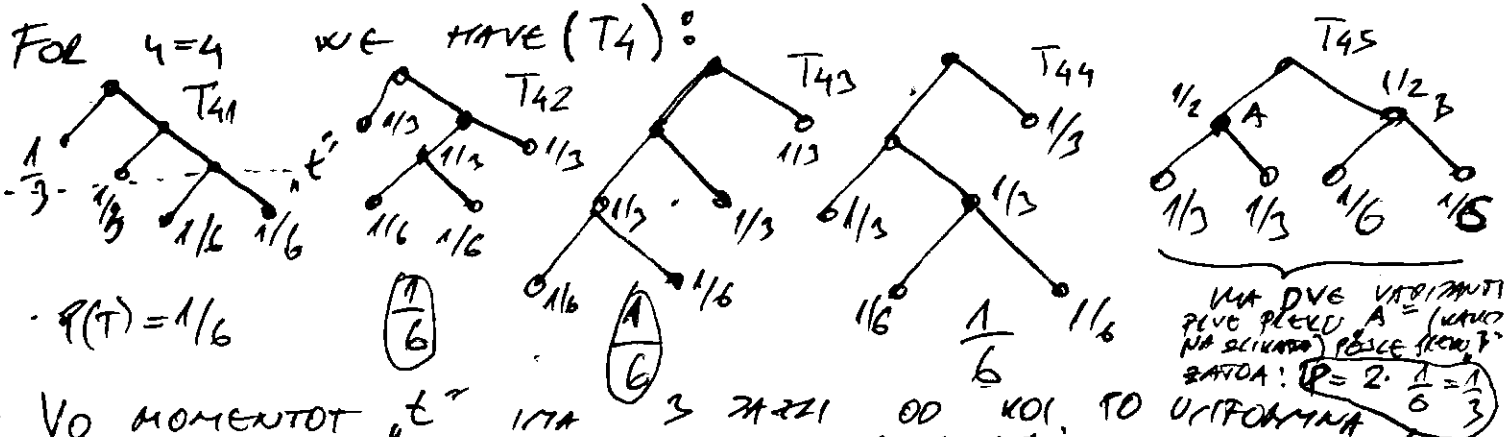


CONTINUE THE PROCESS UNTILL NO FURTHER SUBDIVISION CAN BE MADE. (THE EQUIVALENCE OF THESE TWO TREE GENERATION SCHEMES FOLLOWS, FOR EXAMPLE FROM POLYA'S URN MODEL.)

NOW LET T_u DENOTE A RANDOM u -NODE TREE GENERATED AS DESCRIBED. THE PROBABILITY OF SUCH TREES SEEMS DIFFICULT TO DESCRIBE, BUT WE CAN FIND THE ENTROPY OF THIS DISTRIBUTION IN RECURSIVE FORM. FIRST SOME EXAMPLES. FOR $u=2$ WE HAVE ONLY ONE TREE THUS $H(T_2) = 1 \cdot \log 1 = 0$. FOR $u=3$ WE HAVE TWO EQUALLY PROBABLE TREES:

THUS $H(T_3) = \left(\frac{1}{2} \log 2\right) \cdot 2 = \log 2$





$$8 = 4 \cdot \frac{1}{3} + \frac{1}{3} = \frac{2}{2} + \frac{1}{3} = 1$$

$P(T_4)$ NEMA VREMA SO GLAVNENESTO NA PIVOTO TUKO SO VIKURJOT PILEV NA DEVA I BRJOT NA PIVOTUVANEN NA 2

NOW FOR THE RECURRENCE RELATION. LET $N_1(T_n)$ DENOTE THE NUMBER OF TERMINAL NODES OF T_n IN THE RIGHT HALF OF THE TREE.

$$H(T_n) \stackrel{(a)}{=} H(N_1, T_n) \stackrel{(b)}{=} H(N_1) + H(T_n | N_1) \stackrel{(c)}{=} \log(n-1) + H(T_n | N_1) \stackrel{(d)}{=} \log(n-1) + \frac{1}{n-1} \sum_{k=1}^{n-1} (H(T_k) + H(T_{n-k})) \stackrel{(e)}{=} \log(n-1) + \frac{2}{n-1} \sum_{k=1}^{n-1} H(T_k)$$

(f) USE THIS TO SHOW THAT:

$$(n-1)H_n = 4H_{n-1} + (n-1)\log(n-1) - (n-2)\log(n-2) \text{ OR}$$

$$\frac{H_n}{n} = \frac{H_{n-1}}{n-1} + c_n$$

FOR APPROXIMATELY DEFINED c_n

SINCE $\sum c_n = C < \infty$, YOU HAVE PROVED THAT $\frac{1}{n} H(T_n)$ CONVERGES TO A CONSTANT. THUS, THE EXPECTED NUMBER OF BITS NECESSARY TO DESCRIBE RANDOM TREE T_n GROWS LINEARLY WITH n .

(e) BY USAGE OF CHAIN FORMULA

$$H(N_1, T_n) = H(N_1) + H(T_n | N_1)$$

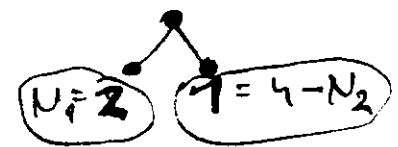
(c) N_1 IS UNIFORMLY DISTRIBUTED OVER $\{1, 2, \dots, n-1\}$

i.e. $P(N_1) = \frac{1}{n-1} \Rightarrow H(N_1) = \sum_{i=1}^{n-1} \frac{1}{n-1} \log(n-1)$

$N_1 \in \{1, 2, \dots, n-1\}$

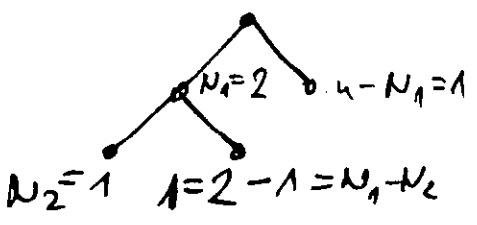
- PROVA NA RECURRENCE RELATIONOT ZA $n=3$

10) $N_1 \in \{1, 2\}$ $P(N_1) = \frac{1}{2}$ e.g. $N_1 = 2$



2° $N_2 \in \{1\}$

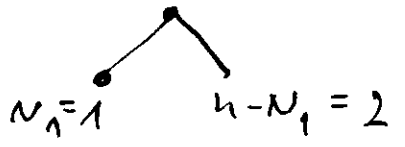
$N_3 \in \{1.. 3-2-1\} \in \{0\}$



• DLUQIOT PAT

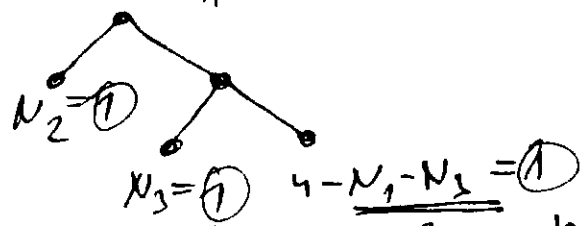
1°) $N_1 \in \{1, 2\}$

$P(N_1) = \frac{1}{2}$ $N_1 = 1$



2°) $N_2 \in \{1.. N_1-1\} \in \{\}$

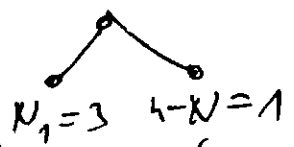
$N_3 = \{1.. \frac{3-1-1}{4-N_1-1}\} \in \{1\}$



• RECURRENCE

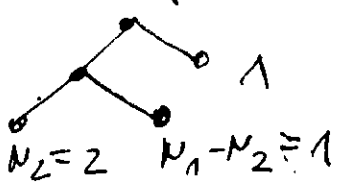
1°) $N_1 \in \{1, 2, 3\}$

$N_1 = 3$ $P(N_1) = \frac{1}{3}$



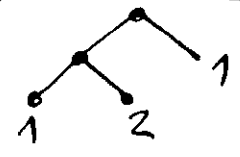
2a°) $N_2 \in \{1, 2\}$

$P(N_2) = \frac{1}{2}$; $N_2 \in \{1.. 4-3-1\} = \{\}$



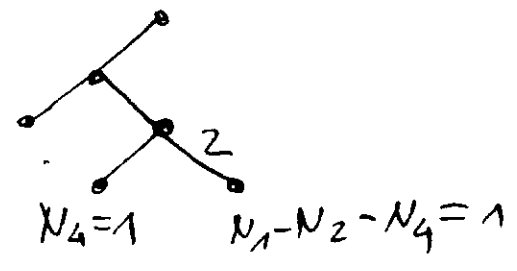
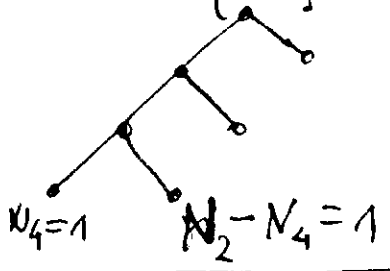
4.) $N_2 = 2$

6.) $N_2 = 1$

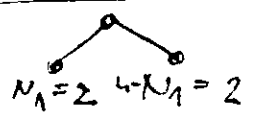


3a°) $N_4 \in \{1, 2, .. N_2-1\}$
 $N_4 \in \{1\}$

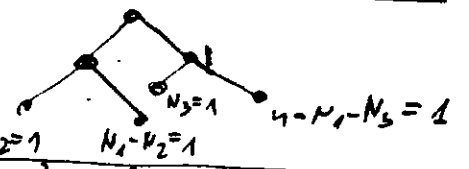
3c°) $N_4 \in \{1, 2, .. N_1-N_2-1\} \in \{1\}$



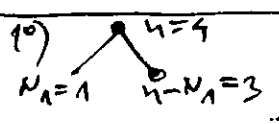
1°) $N_1 = 2$



2°) $N_2 \in \{1\}$
 $N_3 \in \{1\}$

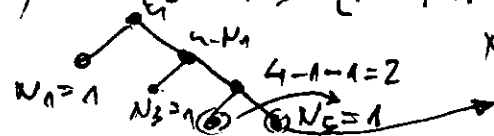


$N_1 = 1$

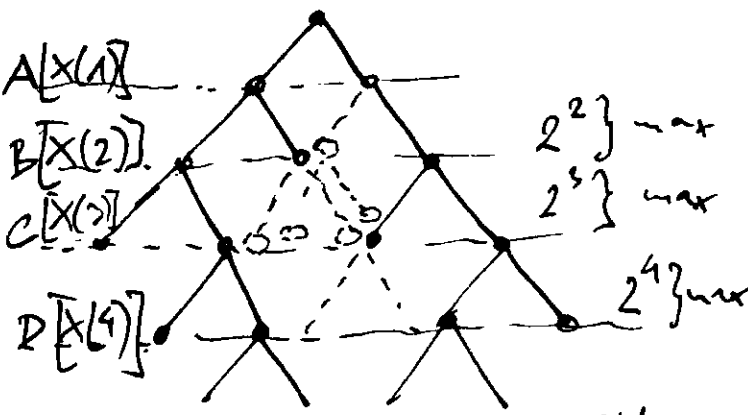


2°) $N_2 \in \{\}$

$N_3 \in \{1.. 4-1-1\} \in \{1, 2\}$ $N_3 = 1$



$N_4 \in \{\}$ $N_5 \in \{1.. 2-1\} = \{1\}$
 $4-N_1-N_3-N_5 = 4-2=1$



$$X' = \begin{bmatrix} 1 & 2 & 3 & 4 & 1 \\ 0 & 2 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 2 \end{bmatrix} \begin{matrix} \\ \\ \end{matrix} \left. \vphantom{\begin{matrix} \\ \\ \end{matrix}} \right\} \begin{matrix} \\ \\ \end{matrix} \begin{matrix} \\ \\ \end{matrix}$$

$$X'' = \begin{bmatrix} 0 & 1 & 2 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix} \begin{matrix} 1 \\ 2 \end{matrix}$$

$n=4$

$T_{41} =$

1	2	3	4	5	6	7	8
			1	3			
				1	2		
					1	1	

$T_{42} =$

1	2	3	4	5	6	7	8
			1	3			
				2	1		
			1	1			

$T_{45} =$

1	2	3	4	5	6	7	8
			2	2			
			1	1	1	1	

$$H(N_1) = P(N_1=1) \log \frac{1}{P(N_1=1)} + P(N_1=2) \log \frac{1}{P(N_1=2)} + P(N_1=3) \log \frac{1}{P(N_1=3)}$$

$$= \left(\frac{1}{3} \log 3\right) \cdot 3 = \log 3$$

PRIMARNA NA DUKOVATI

1 2 2 3 4 3 1 2 1 2 1 2 2 1 2 1 2 ; 3 1 2 ; 2 1 2

$T_{43} =$

1	2	3	4	5	6	7	8
				3	1		
				2			
				1	1		

$T_{44} =$

1	2	3	4	5	6	7	8
			3	1			
			1	2			
			1	1			

• G1 i G2 implementiraju T_3 i T_4 DA G1 GENERIŠE ŽRANJE.

(a) $H(N_1, T_4) = H(T_4) + H(N_1 | T_4) = \left| \begin{matrix} N_1 = f(T_4) \\ \Rightarrow H(N_1 | T_4) = 0 \end{matrix} \right| = H(T_4)$

$H(T_4) = H(N_1, T_4) = H(N_1) + H(T_4 | N_1)$

(d) $H(T_4 | N_1) = ?$ $H(T_4 | N_1) = \frac{1}{n-1} \sum_{k=1}^{n-1} (H(T_k) + H(T_{n-k}))$

e.g. $N_1=3$ $n=4$ (FIG ON PR. 112)

$$H(T_4 | N_1) = \frac{1}{3} \sum_{k=1}^3 H(T_k) + H(T_{4-k}) = \frac{1}{3} [H(T_1) + H(T_3) + H(T_2) + H(T_2) + H(T_1)] = \frac{2}{3} [H(T_1) + H(T_2) + H(T_3)]$$

• ДИРЕКТНА РЕКРЕТНА НА :

$$H(T_4|N_1=3) = H(T_{41}, T_{42}) = \frac{1}{2} \log_2 2 + \frac{1}{2} \log_2 2 = 1$$

$$T_{43} H(T_4|N_1=2) = 1 \log_2 1 = 0$$

$$H(T_4|N_1=1) = \frac{1}{2} \log_2 2 + \frac{1}{2} \log_2 2 = 1$$

T_{42}, T_{44}

$$H(T_4|N_1) = P(N_1=3) \cdot H(T_4|N_1=3) + P(N_1=2) \cdot H(T_4|N_1=2) + P(N_1=1) \cdot H(T_4|N_1=1)$$

$$= \frac{1}{3} \cdot 1 + \frac{1}{3} \cdot 1 + \frac{1}{3} \cdot 0 = \frac{2}{3}$$

$$H(T_4) = 4 \cdot \frac{1}{6} \log_2 6 + \frac{1}{3} \log_2 3 = \frac{2}{3} \log_2 2 + \frac{2}{3} \log_2 3 + \frac{1}{3} \log_2 3$$

$$= \frac{2}{3} + \log_2 3$$

$$H(T_4) = H(N_1) + H(T_4|N_1) = \left(\frac{1}{3} \log_2 3\right) \cdot 3 + \frac{2}{3} = \log_2 3 + \frac{2}{3}$$

ПРОБЛЕМА НА П. 117

- DREAM ABOUT CONDITIONAL PROBABILITY

MMV

$X \backslash Y$	\diamond	\heartsuit	\spadesuit	\clubsuit	$P(X)$
RED	$\frac{1}{2}$	$\frac{1}{2}$	0	0	$\frac{1}{2}$
BLACK	0	0	$\frac{1}{4}$	$\frac{1}{4}$	$\frac{1}{2}$
$P(Y)$	$\frac{1}{4}$	$\frac{1}{4}$	$\frac{1}{4}$	$\frac{1}{4}$	

X - COLOUR

Y - CARD FIGURE

e.g. $P(X, Y) = (RED, \diamond) =$

$$= P(RED) \cdot P(\diamond) = \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4}$$

$P(Y|X)$

$X \backslash Y$	\diamond	\heartsuit	\spadesuit	\clubsuit
RED	$\frac{1}{2}$	$\frac{1}{2}$	0	0
BLACK			$\frac{1}{2}$	$\frac{1}{2}$

• АКО ЗНАЕМ ДЕКА N_1 НЕ Е САРМО ЕА РЕЗУЛТАТ СВАКА - ОД ПРВОТО ЈЕ ПОЗИВА :

$$H(T_4|N_1=3) = H(T_{41}, T_{42}, T_{43}, T_{44}) = 4 \cdot \frac{1}{4} \log_2 4 = 2$$

• АКО ОРАМ СО ВЕРОЈАНОСТИ: $\frac{1}{5}, \frac{1}{5}, \frac{1}{5}, \frac{1}{5}, \frac{1}{5}$

$$H(T_4) = \left(\frac{1}{5} \log_2 5\right) 5 = \log_2 5$$

$$H(T_4|N_1=3) = \left(\frac{1}{2} \log_2 2\right) \cdot 2 = 1 \quad H(T_4|N_1=2) = 1 \log_2 1 = 0$$

$$H(T_4|N_1=1) = \left(\frac{1}{2} \log_2 2\right) 2 = 1$$

$$H(T_4, N_1) = P(N_1=3) \cdot H(T_4|N_1=3) + P(N_1=2) \cdot H(T_4|N_1=2) + P(N_1=1) \cdot H(T_4|N_1=1) = \frac{1}{3} + \frac{1}{3} = \frac{2}{3}$$

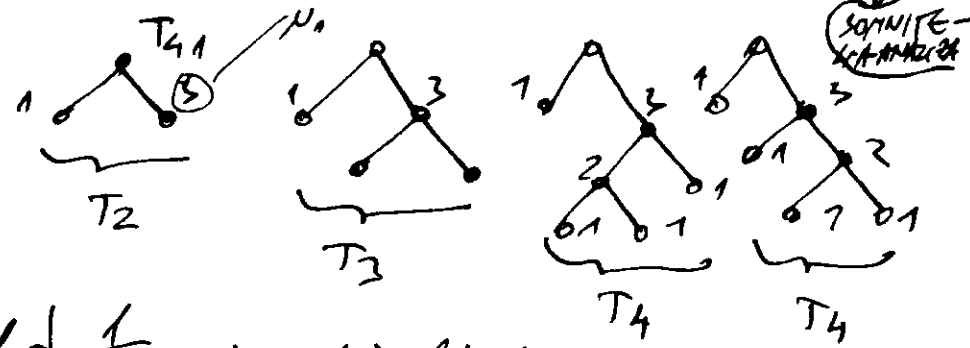
• AVO SE ZEMAT VUMOT DLOV NA DVA KOI SE

179 LADENI:

$N_1=3$

T_{41}, T_{42}

(?)

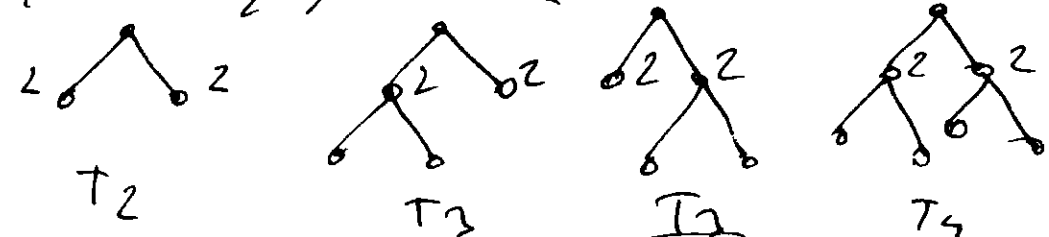


$$H(T_4 | N_1=3) = P(T_2) \log \frac{1}{P(T_2)} + P(T_3) \log \frac{1}{P(T_3)} + P(T_4) \log \frac{1}{P(T_4)}$$

$$= \frac{1}{4} \log 4 + \frac{1}{4} \log 4 + \frac{1}{2} \log 2 = \frac{3}{2}$$

$N_1=2$

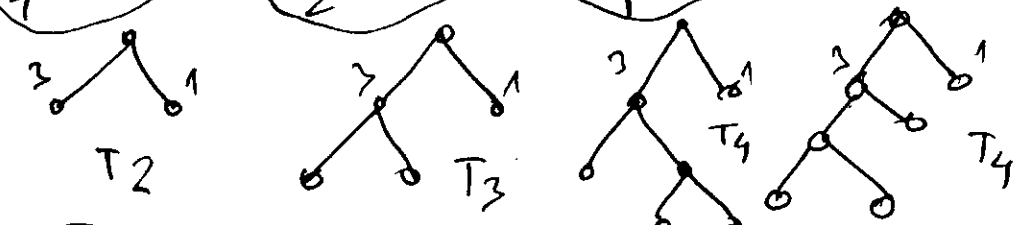
T_{45}



$$H(T_4 | N_1=2) = \left(\frac{1}{4} \log 4\right) + \left(\frac{1}{2} \log 2\right) + \left(\frac{1}{4} \log 4\right) = \frac{3}{2}$$

$N_1=1$

T_{42}, T_{44}



$$H(T_4 | N_1=1) = \left(\frac{2}{4} \log 4\right) + \left(\frac{1}{4} \log 4\right) + \left(\frac{1}{2} \log 2\right) = \frac{3}{2}$$

$$H(T_4 | N_1) = P(N_1=3) H(T_4 | N_1=3) + P(N_1=2) H(T_4 | N_1=2) + P(N_1=1) H(T_4 | N_1=1)$$

$$= \frac{1}{3} \cdot \frac{3}{2} \log 4 + \frac{1}{3} \cdot \frac{1}{4} \log 4 + \frac{1}{3} \cdot \frac{2}{4} \log 4 = \frac{1}{3} + \frac{1}{6} + \frac{1}{3} = \frac{5}{6}$$

$$H(T_3 | N_1) = \frac{1}{3} \cdot \left(\frac{1}{4} \log 4\right) + \frac{2}{4} \log 4 + \frac{1}{4} \log 4 = \frac{2}{3}$$

$$T(T_2 | N_1) = \frac{1}{3} \left(\frac{1}{2} + \frac{1}{2} + \frac{1}{2} \right) = \frac{1}{2}$$

$$H(T_3) = \frac{4}{12} \log 3 = \frac{1}{3} \log 3 \quad T(T_2) = \frac{1}{4} \log 4 = \frac{1}{2}$$

$H(T_4) = \frac{5}{12} \log 12$

$H(T_4) = H(N_1) + H(T_4 | N_1)$

$= \frac{1}{3} \log 3 + \frac{5}{6}$

$H(T_4) = \frac{5}{12} \log(4 \cdot 3) = \frac{5}{12} \log 4 + \frac{5}{12} \log 3 = \frac{5}{6} + \frac{5}{12} \log 3$ (?)

• Aho st. zremt vo rledvid 1 "T1"

$$H(T_4|N_1) = \frac{1}{3} \left(\underbrace{H(T_4|N_1=3)}_{2/5 \text{ lds}} + \underbrace{H(T_4|N_1=2)}_{1/5 \text{ lds}} + \underbrace{H(T_4|N_1=1)}_{2/5 \text{ lds}} \right) =$$

$$= \frac{1}{3} \left(\text{lds} \left(\frac{2}{5} \cdot 2 + \frac{1}{5} \right) - \text{lds} \left(\frac{2}{5} + \frac{2}{5} \right) \right) = \frac{1}{3} \left(\text{lds} - \frac{4}{5} \text{ lds} \right)$$

$$= \frac{1}{3} \left(\text{lds} - \frac{4}{5} \right) \quad (?)$$

$$H(T_4) = \frac{2}{3} \text{ lds} + \frac{1}{3} \text{ lds} - \frac{4}{15}$$

$$H(T_4) = \frac{1}{3} \cdot \text{lds}$$

$$H(N_1) = P(N_1=3) \text{ lds} \frac{1}{P(N_1=3)} + P(N_1=2) \text{ lds} \frac{1}{P(N_1=2)} + P(N_1=1) \text{ lds} \frac{1}{P(N_1=1)} = 3 \frac{1}{3} \text{ lds} = \text{lds}$$

! # (SOMNIVENNA!!!)
KANIZETA!!!

$$H(T_2) = 0 \quad H(T_3) = \left(\frac{1}{2} \text{ lds} \right) \cdot 2 = \text{lds}$$

↓ PROBLÉMA
OD PR. 15

$$H(T_4) = 4 \frac{1}{6} \text{ lds} + \frac{1}{3} \text{ lds} = \frac{2}{3} \text{ lds} + \frac{2}{3} \text{ lds} + \frac{1}{3} \text{ lds}$$

$$H(T_4) = \frac{2}{3} + \text{lds}$$

$$H(T_4|N_1) = \frac{2}{3} [H(T_1) + H(T_2) + H(T_3)] = \frac{2}{3} \text{ lds} = \frac{2}{3} \text{ lds}$$

$$H(T_4|N_1) = \frac{1}{N-1} \sum_{k=1}^{N-1} [H(T_k) + H(T_{N-k})]$$

$$H(T_4|N_1) = P(N_1=4-1) H(T_4|N_1=4-1) + P(N_1=4-2) H(T_4|N_1=4-2) + \dots + P(N_1=1) H(T_4|N_1=1)$$

$$P(N_1) = \frac{1}{4-1}$$

$$H(T_4|N_1) = \frac{1}{4-1} \sum_{i=1}^{4-1} H(T_4|N_1=i)$$

e.g. $H(T_4|N_1) = P(N_1=3) H(T_4|N_1=3) + P(N_1=2) H(T_4|N_1=2) + P(N_1=1) H(T_4|N_1=1)$

$$= \frac{1}{3} \left[\frac{2}{2} \text{ lds} + 1 \cdot \text{lds} + \frac{2}{2} \text{ lds} \right] = \frac{2}{3}$$



$$H(T_4|N_1) = \frac{1}{3} \sum_{i=1}^3 H(T_4|N_1=i) = \frac{1}{3} \sum_{i=1}^3 [H(T_{4-i}) + H(T_i)]$$

$$H(T_4|N_1=3) = P(N_2=2) \cdot H(T_3|N_2=2) + P(N_2=1) \cdot H(T_3|N_2=1) = \frac{1}{2} = H(T_3) = \text{lds}$$

$$H(T_4|N_1=2) = H(T_3)$$

2¹
PODA
ZATOK STO
CERAJ SUMARI
ZA DESKAT I
CERAJ ZA LONAR
STONO

SUMARA

- SUŠTINSKI ZAKLJČOK I DOKAZ:

$$H(T_n | N_n = i) = H(T_{n-i}) + H(T_i)$$

(*) $H(T_4 | N_1 = 3) = H(T_3)$
 90 IMA DVA PATT
 NA LEVA I DESNA STRANA

$$H(T_5 | N_1) = P(N_1=4) \cdot H(T_5 | N_1=4) + P(N_1=3) \cdot H(T_5 | N_1=3) + P(N_1=2) \cdot H(T_5 | N_1=2) + P(N_1=1) \cdot H(T_5 | N_1=1) =$$

$$= \frac{1}{4} [H(T_4) + H(T_3) + H(T_2) + H(T_1)] \quad (2)$$

4 $(n=5)$

$$\sum_{i=1}^4 H(T_i) + H(T_{n-i}) = H(T_1) + H(T_4) + H(T_2) + H(T_3) + H(T_2) + H(T_2) + H(T_4) + H(T_1) =$$

$$= 2 [H(T_1) + H(\bar{2}) + H(\bar{3}) + H(\bar{4})]$$

90 IMA
 ZA LEVA I
 DESNA STRANA

(*) OVA JE POSLEDICA NA TOA ŠTO SE RADITI ZA MARKOV CAIN

ISTOT REZULTAT

ZNAČI: $H(T_n | N_n) = \frac{1}{n-1} \sum_{i=1}^{n-1} H(T_i) + H(T_{n-i}) = \frac{2}{n-1} \sum_{i=1}^{n-1} H(T_i)$

$$H(T_3) = \left(\frac{1}{2} \cdot \log 2\right) \cdot 2 = \log 2$$

$$H(T_4) = H(N_1) + H(T_4 | N_1) = \log 3 + \frac{2}{3}$$

$$H(T_n) = \log(n-1) + \frac{2}{n-1} \sum_{i=1}^{n-1} H(T_i)$$

$$H(T_5) = \log 4 + \frac{2}{4} \sum_{i=1}^4 H(T_i) = \log 4 + \frac{1}{2} (0 + 0 + 1 + \log 3 + \frac{2}{3})$$

$$H(T_5) = 2 + \frac{1}{2} + \frac{1}{3} + \frac{1}{2} \log 3 = \frac{12+3+2}{6} + \frac{1}{2} \log 3 = \frac{17}{6} + \frac{1}{2} \log 3$$

POTREB SE NA
 CLANAKOT OD KAKI-
 PETNA VEDE ŽIBAN
 ZA RANDOM WALK
 ON NUMBER SCHE
 2 4 5 1 1 1 1
 + 1 1 1 1 1 1 1

(f) $H_n = \log(n-1) + \frac{2}{n-1} \sum_{i=1}^{n-1} H_i$

$$(n-1)H_n = (n-1)\log(n-1) + 2 \sum_{k=1}^{n-1} H_k = (n-1)\log(n-1) + 2H_{n-1} + 2 \sum_{k=1}^{n-2} H_k$$

$$H_{n-1} = \log(n-2) + \frac{2}{n-2} \sum_{k=1}^{n-2} H_k \Rightarrow 2 \sum_{k=1}^{n-2} H_k = (n-2)H_{n-1} - (n-2)\log(n-2)$$

$$(n-1)H_n = (n-1)\log(n-1) + 2H_{n-1} + (n-2)H_{n-1} - (n-2)\log(n-2) =$$

$$= (n-1)\log(n-1) - (n-2)\log(n-2) + 4H_{n-1}$$

$$(n-1)H_n = (n-1)\log(n-1) - (n-2)\log(n-2) + 4H_{n-1}$$

$$H_n = \log(n-1) - \frac{n-2}{n-1} \log(n-2) + \frac{1}{n-1} H_{n-1} \Rightarrow$$

$$\frac{H_n}{n} = \frac{H_{n-1}}{n-1} + \frac{\log(n-1) - \frac{n-2}{(n-1) \cdot n} \log(n-2)}{n} \quad \text{①}$$

$$C_n = \frac{(n-1)\log(n-1) - (n-2)\log(n-2)}{n(n-1)}$$

$$\frac{H_n}{n} = \frac{H_{n-1}}{n-1} + C_n$$

$$\lim_{n \rightarrow \infty} C_n = 0$$

$$\lim_{n \rightarrow \infty} \left(\frac{H_n}{n} - \frac{H_{n-1}}{n-1} \right) = 0$$

Edition 1 - SIEĆTAZ FUTURE (JAK IFO VILICITNO ČEM SO UOVRATAK)

$$H(T_n | N_1 = k) = H(T_n, T_{n-k} | N_1 = k) = H(T_n | N_1 = k) + H(T_{n-k} | N_1 = n-k)$$

OVA 6 VOVA ZAVDA ITO LEVO I PRAVO I PESIHO VO PRVO I SE IZVIRAT NEZAVISNO.

SEGA SO OGLEP DEFICIPIATA NA OLCOVA ENTROPITZA:

$$H(T_n | N_1) = \sum_{k=1}^{n-1} P(N_1 = k) \cdot H(T_n | N_1 = k) = \frac{1}{n-1} \sum_{k=1}^{n-1} [H(T_n | N_1 = k) + H(T_{n-k} | N_1 = n-k)]$$

$$= \frac{1}{n-1} \sum_{k=1}^{n-1} [H(T_k) + H(T_{n-k})] = \left[\begin{array}{l} \text{LEVO I} \\ \text{PRAVO I} \\ \text{PRAVO I} \\ \text{SIEĆTAZ} \end{array} \right] = \frac{2}{n-1} \sum_{k=1}^{n-1} H(T_k)$$

VO EDITION 1 VITE PODRKO GO POKAZENAT SO PROMENA NA NOMENKATIVI

$$\frac{1}{n-1} \sum_{k=1}^{n-1} [H(T_k) + H(T_{n-k})] = \frac{1}{n-1} \sum_{k=1}^{n-1} H(T_k) + \frac{1}{n-1} \sum_{k=1}^{n-1} H(T_{n-k})$$

$$\Rightarrow \left[\begin{array}{l} i = n-k \Rightarrow k=1 \quad i=n-1 \\ k=n-1 \quad i=1 \end{array} \right] = \frac{1}{n-1} \sum_{k=1}^{n-1} H(T_k) + \frac{1}{n-1} \sum_{i=1}^{n-1} H(T_i) =$$

$$= \frac{2}{n-1} \sum_{k=1}^{n-1} H(T_k)$$

(f) AS IN Edition 1:

$$C_n = \frac{\log(n-1)}{n} - \frac{\log(n-2)}{n-1} + \frac{2 \log(n-2)}{n(n-1)}$$

$$\frac{H_{n-1}}{n-1} = \frac{H_{n-2}}{n-2} + C_{n-1} \quad C_{n-1} = \frac{\log(n-2)}{n-1} - \frac{\log(n-3)}{n-2} + \frac{2 \log(n-3)}{(n-2)(n-1)}$$

$$\frac{H_n}{n} = \frac{H_{n-2}}{n-2} + C_n + C_{n-1} = \frac{H_{n-2}}{n-2} + \frac{\log(n-1)}{n} + \frac{2 \log(n-2)}{n(n-1)} = \frac{\log(n-3)}{n-2} + \frac{2 \log(n-3)}{(n-2)(n-1)}$$

$$\text{①} = \frac{-(n-1) \log(n-2) + 2 \log(n-3)}{(n-2)(n-1)} = \frac{n-2}{(n-2)(n-1)} \cdot \frac{(n-3) \log(n-3)}{n-2} = \frac{(n-3) \log(n-3)}{(n-2)(n-1)}$$

070 342764

MAKO

- MOZA
- ELOPOLAT
- ANKETA 2008
- TRANSFORMACIJA
- ORUKA 2008

CE SE SVIATI
VO MALCENI
ITERACIJA (TELESCO-
PING THE SUM)

078 270741 / GOL

1) $\frac{H_n}{n} = \frac{H_{n-1}}{n-1} + \frac{\ln(n-1)}{n} - \frac{\ln(n-2)}{n-1} + \frac{2 \ln(n-2)}{n(n-1)}$

2) $\frac{H_n}{n} = \frac{H_{n-2}}{n-2} + \frac{\ln(n-1)}{n} + \frac{2 \ln(n-2)}{n(n-1)} - \frac{\ln(n-3)}{n-2} + \frac{2 \ln(n-3)}{(n-2)(n-1)}$

$C_n = \frac{\ln(n-1)}{n} - \frac{\ln(n-2)}{n-1} + \frac{2 \ln(n-2)}{n(n-1)}$

$\sum_{j=n-1}^n \frac{2 \ln(j-2)}{j(j-1)} = \frac{2 \ln(n-3)}{(n-1)(n-2)} + \frac{2 \ln(n-2)}{n(n-1)}$ OD 20

$\sum_{j=n}^n \frac{2 \ln(j-2)}{j(j-1)} = \frac{2 \ln(n-2)}{n(n-1)}$ OD 10

3) $\frac{H_{n-2}}{n-2} = \frac{H_{n-3}}{n-3} + C_{n-2} = \frac{H_{n-3}}{n-3} + \frac{\ln(n-3)}{n-2} - \frac{\ln(n-4)}{n-3} + \frac{2 \ln(n-4)}{(n-2)(n-3)}$

$\sum_{j=n-2}^n \frac{2 \ln(j-2)}{j(j-1)} = \frac{2 \ln(n-4)}{(n-2)(n-3)} + \frac{2 \ln(n-3)}{(n-1)(n-2)} + \frac{2 \ln(n-2)}{n(n-1)}$

CE SE SVIATI
VO MALCENI
ITERACIJA

30 → 20

$\frac{H_n}{n} = \frac{H_{n-3}}{n-3} + \frac{\ln(n-1)}{n} + \frac{2 \ln(n-2)}{n(n-1)} + \frac{2 \ln(n-3)}{(n-1)(n-2)} - \frac{\ln(n-4)}{n-3} + \frac{2 \ln(n-4)}{(n-2)(n-3)}$

$\frac{H_n}{n} = \frac{\ln(n-1)}{n} + \sum_{j=3}^n \frac{2 \ln(j-2)}{j(j-1)}$

POSLEDNOT
CLEN OD 20

POSLEDNOT
CLEN OD 20

POSLEDNOT
CLEN OD 10

CE SE SVIATI
VO MALCENI

$\lim_{n \rightarrow \infty} \frac{H_n}{n} = \lim_{n \rightarrow \infty} \frac{\ln(n-1)}{n} + \sum_{j=3}^{\infty} \frac{2 \ln(j-2)}{j(j-1)} \leq \sum_{j=3}^{\infty} \frac{2 \ln(j-1)}{(j-1)^2}$

$= \sum_{k=2}^{\infty} \frac{2 \ln k}{k^3} \leq \sum_{j=2}^{\infty} \frac{2 \sqrt{j}}{j^2} = \sum_{j=2}^{\infty} 2 \cdot j^{-3/2} = \sum_{j=2}^{\infty} \frac{2}{\sqrt{j^3}}$

POŠTO $\frac{H_2}{2} = 1$

CONVERGEN!

$\lim_{n \rightarrow \infty} \frac{H_n}{n} = 1.736$
GITS